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The design and implementation of the FIRE engine:
A C++ toolkit for Finite automata
and Regular Expressions

by

B.W. Watson

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The design and implementation of the FIRE engine: A C++ toolkit for FInite automata and Regular Expressions

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August 15, 1994

Abstract
This paper describes the design and implementation of version 1.2 of the FIRE engine. The FIRE engine is a C++ class library implementing finite automata and regular expression algorithms. The algorithms implemented in the toolkit are almost all of those presented in the taxonomies of finite automata algorithms [Wat93a, Wat93b]. The reader is assumed to be familiar with the two taxonomies and with advanced C++ programming techniques.

The toolkit is implemented largely in an object-oriented style, with finite automata and regular expressions being defined as classes. All of the classes and functions in the toolkit are presented in the same format. For each class (or function) the format includes a short description of its behaviour, details of its implementation, and techniques for improving its performance.

*Second printing, with corrections.*
CONTENTS

Contents

1 Introduction 3
   1.1 Coding conventions and performance issues 3
       1.1.1 Performance tuning 4
   1.2 Presentation conventions 4

   I Foundation classes 5
   2 Character ranges 5
   3 Bit vectors 15
   4 States and related classes 24
   5 Transitions and related classes 35
   6 Relations 47

   II Regular expressions and \( \Sigma \)-algebras 67
   7 Regular expressions 67
       7.1 Similarity of \( REs \) 75
       7.2 Derivatives 83
   8 The \( \Sigma \)-algebra 84
       8.1 The \( \Sigma \)-term algebra \( Reg<RE> \) 86

   III Automata and their components 89
   9 Abstract finite automata 89
   10 Deterministic \( FAs \) 91
   11 Constructing a \( DFA \) 98
       11.1 Abstract states and the subset construction 98
       11.2 Derivatives of regular expressions 103
       11.3 Item sets 107
   12 Finite automata 121
       12.1 The \( \Sigma \)-algebra \( Reg<FA> \) 128
   13 Reduced finite automata 133
       13.1 The \( \Sigma \)-algebra \( Reg<RFA> \) 144
   14 Left-biased finite automata 147
   15 Right-biased finite automata 154
   16 Automata constructions 161

   IV Minimizing \( DFAs \) 165
1 Introduction

The **FIRE engine** is a class library implementing finite automata and regular expression algorithms. The algorithms implemented are all of the algorithms presented in the taxonomies of finite automata algorithms [Wat93a, Wat93b]. This paper outlines the design and implementation of the class library, which consists of 9000 lines of C++ code. A number of other toolkits (mostly serving different purposes) are available; they are detailed in [Wat94], which also contains an introduction to the **FIRE engine**.

In order to understand the implementation of the **FIRE engine**, the reader is assumed to be familiar with the taxonomies [Wat93a, Wat93b], and with C++ terminology and programming techniques. An introduction to the C++ language can be found in [Lip91, Str91] (the language reference appears in [Str91, SE90]), while an introduction to C++ programming techniques can be found in [Bud94]. Advanced C++ programming techniques are presented in [Cop92, MeS92, Mur93]. Object-oriented programming, in general, is discussed in [Boo94, Bud91, MeB88].

The most effective way to understand the implementation of the class library is to read the code and the comments embedded within the code. The explanations given in this paper are intended as an outline, to assist in understanding the code.

This paper is structured as follows:

- Part I provides descriptions of all of the foundations classes — those classes used in the implementations of other classes.
- Part II gives descriptions of all classes related to regular expressions and \( \Sigma \)-algebras.
- Part III gives descriptions of all classes implementing finite automata.
- Part IV presents the functions implementing deterministic finite automata minimization algorithms.

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1.1 Coding conventions and performance issues

The toolkit is coded in a style which makes heavy use of the `assert` function. For all except the simplest classes, a member function called `class_invariant` is defined; this member returns 1 if the object is structurally sound, and 0 otherwise. All member functions which are visible to clients should only be invoked on structurally sound objects and must leave the object in a structurally sound condition. The function call `assert(class_invariant())` is explicitly invoked at the beginning and end of all such member functions (except, of course, the function `class_invariant`). The definition of the class invariant is a useful starting point for understanding the structure and implementation of a particular class.

Additional assertions are used in places where the incorrect values of variables (or relations between variables) will cause catastrophic results (i.e., the machine is likely to crash). Assertions are also used in places where they may be of use in understanding the surrounding code. As mentioned below, `assert` can be disabled by compiling with the macro `NDEBUG` defined — meaning that there is no assertion overhead in a production version of the **FIRE engine**.

Almost all of the classes have an `ostream` insertion operator defined (an output operator). For implementation classes the output members are mainly for debugging purposes.

Many classes provide a member function called `reincarnate`. This function clears the object, leaving it in the same condition as it would be in after the argumentless constructor for the object. The only difference is that any memory already allocated by the members of the object is retained for future use.
1.1.1 Performance tuning

There are some general techniques that can be used to improve the performance of the various parts of the FIRE engine toolkit. The easiest method of improving the performance is to compile the toolkit with the macro \texttt{NDEBUG} defined. This will disable the \texttt{assert} function call, and should only be done when the program is functioning correctly and the safety of assertions and class invariants is no longer needed. Before hand-optimizing the code itself, an important first step is to profile the execution of the toolkit. Some preliminary profiling shows that a large amount of time is spent copying and allocating memory. Much of this time can be eliminated by using two techniques:

- Define and use special versions of the \texttt{new} and \texttt{delete} operators for the classes that make extensive use of memory management.
- "Use-count" all classes which copy large amounts of data in their copy constructors or assignment operators. (See any of [Cop92, MeS92, Mur93] for explanations of "use-counting").

1.2 Presentation conventions

Each class is presented in a standard format. The following items appear in the description of a class:

1. The name of the class, and whether it is a user level class (a class intended for use by program components internal and external to the FIRE engine), or an implementation level class (one used only by other classes within the FIRE engine).
2. The Files clause lists the files in which the class declaration and definition are stored. The names of the files associated with a class are usually the class name, followed by .h for interface (or header) files and .cpp for definition files. All file names are short enough for use under MS-DOS.
3. The Uses clause lists the classes of which this class is an immediate client or from which this class inherits directly. (Most classes use themselves; this is not mentioned.)
4. The Description clause gives a brief description of the purpose of the class.
5. The .h file (if there is one) corresponding to the class is listed.
6. The optional Implementation clause outlines the implementation of the class.
7. The optional Performance clause gives some suggestions on possible performance improvements to the implementation of the class.
8. The .cpp file (if there is one) corresponding to the class is listed.
9. The description ends with the $\square$ symbol.
Part I
Foundation classes

2 Character ranges

The header file extras.h contains inline functions for computing minima and maxima of two numbers. The functions are overloaded to provide versions for types char and int.

User class: CharRange

Files: charrang.h, charrang.cpp

Uses:

Description: Sets of characters (denoted by a CharRange) can be used as automata transition labels and as regular expression atoms (unlike most other systems, which use single characters). The sets of characters are limited to a range of characters in the execution character set (no particular character set is assumed, and ASCII or EBCDIC can both be used). The ranges are specified by their upper and lower (inclusive) bounds. The denoted range may not be empty (allowing them to be empty greatly complicates the algorithms). Many member functions are provided, including functions determining if two CharRanges have a character in common, if one is entirely contained in another, or if they denote sets that are adjacent to one another (in the execution character set).
private:
    char lo, hi;
};

// Copy constructor is standard.
inline CharRange( const CharRange& r );
inline CharRange( const char l, const char h );
inline CharRange( const char a );

// Default operator=, destructor are okay

// Some normal member functions:
// Is char a in the range?
inline int contains( const char a ) const;

// Access to the representation:
inline char low() const;
inline char high() const;

// Standard comparison operators.
inline int operator==( const CharRange& r ) const;
inline int operator!=( const CharRange& r ) const;

// Containment operators.
inline int operator<=( const CharRange& r ) const;
inline int operator>=( const CharRange& r ) const;

// Is there something in common between *this and r?
inline int not_disjoint( const CharRange& r ) const;

// Do *this and r overlap, or are they adjacent?
inline int overlap_or_adjacent( const CharRange& r ) const;

// Some operators on CharRanges.
// Merge two CharRanges if they're adjacent:
inline CharRange& merge( const CharRange& r );

// Make *this the intersection of *this and r.
inline CharRange& intersection( const CharRange& r );

// Return the left and right excess of *this with r (respectively):
inline CharRange left_excess( const CharRange& r ) const;
inline CharRange right_excess( const CharRange& r ) const;

// Define an ordering on CharRange's, used mainly in RE::ordering().
int ordering( const CharRange& r ) const;

// The class structural invariant.
inline int class_invariant( ) const;

// Some extra stuff.
friend ostream& operator<<( ostream& os, const CharRange r );
inline CharRange::CharRange( const char I, const char h ) : 
    lo( I ),  
    hi( h ) { 
    assert( class_invariant() ); 
}

inline CharRange::CharRange( const char a ) : 
    lo( a ),  
    hi( a ) { 
    assert( class_invariant() ); 
}

inline int CharRange::contains( const char a ) const { 
    assert( class_invariant() ); 
    return( (lo <= a) && (a <= hi) ); 
}

inline char CharRange::low() const { 
    assert( class_invariant() ); 
    return( lo ); 
}

inline char CharRange::high() const { 
    assert( class_invariant() ); 
    return( hi ); 
}

inline int CharRange::operator==( const CharRange& r ) const { 
    assert( class_invariant() ); 
    assert( r.class_invariant() ); 
    return( (lo == r.lo) && (hi == r.hi) ); 
}

inline int CharRange::operator!=( const CharRange& r ) const { 
    assert( class_invariant() ); 
    assert( r.class_invariant() ); 
    return( !operator==( r ) ); 
}

inline int CharRange::operator<=( const CharRange& r ) const { 
    assert( class_invariant() ); 
    assert( r.class_invariant() ); 
    return( r.lo <= lo && hi <= r.hi ); 
}

inline int CharRange::operator>=( const CharRange& r ) const { 
    assert( class_invariant() ); 
    assert( r.class_invariant() ); 
    return( lo <= r.lo && r.hi <= hi ); 
}

inline int CharRange::noLdisjoint( const CharRange& r ) const { 
    assert( class_invariant() ); 
    assert( r.class_invariant() ); 
    return( (lo <= r.hi) && (hi >= r.lo) ); 
}

inline int CharRange::overlap_or_adjacent( const CharRange& r ) const { 
    assert( class_invariant() ); 
    assert( r.class_invariant() ); 
    return( noLdisjoint( r ) || (lo <= r.hi + 1) || (r.hi <= lo - 1) ); 
}

inline CharRange& CharRange::merge( const CharRange& r ) {
assert( overlap_or_adjacent( r ) );
lo = (char)min( lo, r.lo );
hi = (char)max( hi, r.hi );
return( *this );
}

inline CharRange& CharRange::intersection( const CharRange& r ) {
lo = (char)max( lo, r.lo );
hi = (char)min( hi, r.hi );
return( *this );
}

inline CharRange CharRange::left_excess( const CharRange& r ) const {
auto CharRange ret( min( lo, r.lo ), max( lo, r.lo ) - 1 );
return( ret );
}

inline CharRange CharRange::right_excess( const CharRange& r ) const {
auto CharRange ret( min( hi, r.hi ) + 1, max( hi, r.hi ) );
return( ret );
}

inline int CharRange::class_invariant() const {
// This could possibly be wrong if something goes weird with char being
// signed or unsigned.
return( lo <= hi );
}
#endif

Implementation: CharRanges are represented as a pair of characters: the upper and lower
(inclusive) bounds. There is a potential for problems caused by the fact that the language
standard does not specify whether char is signed or unsigned.

Performance: In most functions that take a CharRange, the parameter passing is by value (since
a pair of characters fit into the parameter passing registers on most processors). On some
machines, performance can be improved by passing a const reference.

/* (c) Copyright 1994 by Bruce W. Watson */
/* Revision: 1.2 $ */
/* $Date: 1994/08/15 13:59:41 $ */
#include "charrang.h"

int CharRange::ordering( const CharRange& r ) const {
if( *this == r ) {
  return( 0 );
} else if( lo == r.lo ) {
  return( hi - r.hi );
} else {
  return( lo - r.lo );
}
}

ostream& operator<<( ostream& os, const CharRange r ) {
if( r.hi == r.lo ) {
  return( os << '\' << r.lo << '\' );
} else {
  return( os << '\n' << '\' << r.lo << '\' << '\' << '\' << r.hi << '\' << '\' );
}
Implementation class: **CRSet**

Files: crset.h, crset.cpp

Uses: **CharRange**

**Description:** A **CRSet** is a set of **CharRanges**, with a restriction: all **CharRanges** in a **CRSet** are pairwise disjoint. When a new **CharRange** is added to a **CRSet**, some **CharRanges** may be split to ensure that the disjointness property is preserved. Members are provided to check if two **CRSets** cover the same characters, and to combine two **CRSets**.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
/** $Revision: 1.2 $ */
/** $Date: 1994/08/15 14:00:25 $ */
#ifndef CRSET_H
#define CRSET_H
#include "charrang.h"
#include <assert.h>
#include <iostream.h>

// A data-structure to hold sets of CharRange's. The CharRange's are kept pairwise disjoint.
// These are generally used in the construction of DFA's, and in DFA minimization.

class CRSet {
public:
    // Some basic constructors.
    // Create the empty set:
    inline CRSet();

    // The singleton set:
    CRSet( const CharRange a );
    CRSet( const CRSet& r );

    // Also need a destructor and an operator=
    inline ~CRSet();

    const CRSet& operator=( const CRSet& r );

    // Some basic operations.
    // Add a CharRange to this set, splitting it as required.
    // It's not possible to add the empty CharRange to a CRSet!
    inline CRSet& add( CharRange a );

    // Append a CharRange onto the end of this set, knowing that a is disjoint with
    // any others in the set.
    inline CRSet& append( const CharRange a );

    // Combine two CRSet's, splitting CharRange's to maintain their pairwise
    // disjointness.
    CRSet& combine( const CRSet& a );

    // Does *this and r cover the same set of characters?
    int equivalent_cover( const CRSet& r ) const;

    // How many CharRanges are in *this?
    inline int size() const;

    // Some iterators:
    // Fetch the i'th CharRange in *this.
    inline const CharRange& iterator( const int i ) const;
};
```
// Is there even an i'th CharRange?
inline int iter_end( const int it ) const;

// Some special members:
inline int class_invariant() const;
friend ostream& operator<<( ostream& os, const CRSet& it );

private:
    // A simple helper.
    void ensure_min_size( int s );

    // Add a CharRange to this set, splitting it as required.
    // Adding begins at i. This is used in CRSet::add().
    void add_at( CharRange a, int i );

    // Is r covered by this set?
    int covered( const CharRange r ) const;

    // A class constant:
    // must be at least 1.
    enum { expansion_size = 5 };

    // Some implementation details:
    int howmany;
    int in_use;
    CharRange *data;
};

inline CRSet::CRSet() :
    howmany( 0 ),
    in_use( 0 ),
    data( 0 ) {
    assert( class_invariant() );
}

inline CRSet::~CRSet() {
    assert( class_invariant() );
    delete [] data;
}

// Insert a without checking to see if it overlaps with any of them already.
inline CRSet& CRSet::append( const CharRange a ) {
    assert( class_invariant() );
    assert( a.class_invariant() );
    ensure_min_size( in_use + 1 );
    data[in_use++] = a;
    assert( class_invariant() );
    return( *this );
}

// Add a CharRange to this set, splitting it as required.
// It's not possible to add the empty CharRange to a CRSet!
inline CRSet& CRSet::add( CharRange a ) {
    assert( class_invariant() );
    add_at( a, 0 );
    return( *this );
}

inline int CRSet::size() const {
    assert( class_invariant() );
    return( in_use );
}

inline const CharRange& CRSet::iterator( const int it ) const {
    assert( class_invariant() );
assert( iter_end( it ) );
return( data[it] );

inline int CRSet::iter_end( const int it ) const {
assert( class_invariant() );
assert( 0 <= it );
return( it >= in_use );
}

inline int CRSet::class_invariant() const {
auto int result;
result = 0 <= howmany
&& 0 <= in_use
&& in_use <= howmany
&& ( howmany == 0 ? data == 0 : data != 0 )
&& 1 <= expansion_size;
// Make sure that all of data[] are pairwise disjoint.
auto int i;
for( i = 0; i < in_use && result; i++ ) {
auto int j;
for( j = i + 1; j < in_use && result; j++ ) {
result = !data[i].not_disjoint( data[j] );
}
return( result );
}

#endif

Implementation: A CRSet contains a pointer to a dynamically allocated array of CharRanges, an integer representing the size of the array, and an integer representing the number of elements in the array that are in use. The member add_at is a good candidate for optimization.

C++:

CRSet::CRSet( const CharRange a ) :
howmany( expansion_size ),
in_use( 1 ),
data( new CharRange [expansion_size] ) {
data[0] = a;
assert( class_invariant() );
}

CRSet::CRSet( const CRSet& r ) :
howmany( r.in_use + expansion_size ),
in_use( r.in_use ),
data( new CharRange [r.in_use + expansion_size] ) {
assert( r.class_invariant() );
auto int i;
for( i = 0; i < in_use; i++ ) {
data[i] = r.data[i];
}
assert( class_invariant() );
}

const CRSet& CRSet::operator=( const CRSet& r ) {
assert( r.class_invariant() );
// May need to allocate some more memory.
if( r.in_use > howmany ) {

auto CharRange *d( new CharRange [howmany = r.in_use + expansion.size] );
delete [] data;
data = d;
}
assert( r.in_use <= howmany );

auto int i;
for( i = 0; i < r.in_use; i++ ) {
data[i] = r.data[i];
}
in_use = r.in_use;
assert( class_invariant() );
return( *this );

void GRS::add_at(CharRange a, int i) {
assert( class_invariant() );
assert( a.class_invariant() );
// Go through data[i] to see if a can be coalesced with data[i].
for( ; i < in_use && !data[i].not_disjoint( a ); i++ );
// Perhaps this is a coalescing candidate.
// There are several possibilities, namely:
// a and data[i] are disjoint
// a <= data[i] (a is entirely contained in data[i]
// data[i] <= a
// a and data[i] are not disjoint (but no containment holds).
// In this case, we have three pieces: the intersection,
// and a slide each from a, data[i].
if( i < in_use ) {
assert( data[i].not_disjoint( a ) );
if( data[i] == a ) {
    // Nothing to do.
} else {
    assert( data[i].not_disjoint( a ) && (data[i] != a) );
    // a needs some splitting. We know that the excess
    // does not fall into data[0..i], we begin adding the excess
    // at data[i+1].
    // For efficiency (not implemented here), we could just
    // append any part of data[i] that doesn't fall within
    // a, since data[i] is pairwise disjoint with all other
    // elements in array data[].
    auto CharRange old( data[i] );
data[i].intersection( a );
    // data[i] now contains old intersection a.
    // Now deal with the left and right excesses of old with a.
    // There are a few possibilities:
    if( a.low() != old.low() ) {
        // Then there is a left excess
        add_at( a.left_excess( old ), i + 1 );
    }
    if( a.high() != old.high() ) {
        // Then there is a right excess
        add_at( a.right_excess( old ), i + 1 );
    }
} else {
    assert( i >= in_use );
    // There may still be something left to put in data[].
    append( a );
}
assert( class_invariantO );
return;
}

CRSet& CRSet::combine( const CRSet& r ) {
assert( r.class_invariantO );
assert( class_invariantO );
auto int i;
for( i = 0; i < r.in_use; i++ ) {
    add( r.data[i] );
}
assert( class_invariantO );
return( *this );
}

int CRSet::equivalent_cover( const CRSet& r ) const {
assert( class_invariantO );
auto int i;
auto int result = 1;
for( i = 0; i < in_use && result; i++ ) {
    result = r.covered( data[i] );
}
return( result );
}

int CRSet::covered( const CharRange r ) const {
assert( class_invariantO );
assert( r.class_invariantO );
auto int i;
for( i = 0; i < in_use && !r.not_disjoint( data[i] ); i++ );
if( i < in_use ) {
    assert( r.not_disjoint( data[i] ) );
    // They have something in common.
    if( r <= data[i] ) {
        // r is completely covered by data[i]
        return( 1 );
    } else {
        // Now, figure out if the excesses are
        // also covered.
        // This could be made more efficient!
        auto int ret1 = 1;
        auto int ret2 = 1;
        if( r.low() != data[i].low() ) {
            // There actually is a left excess:
            ret1 = covered( r.left_excess( data[i] ) );
        }
        if( r.high() != data[i].high() ) {
            // There actually is a right excess:
            ret2 = covered( r.right_excess( data[i] ) );
        }
        return( ret1 && ret2 );
    }
} else {
    assert( i >= in_use );
    // r is not covered by *this.
    return( 0 );
}

// Make sure that data is at least s (adjusting howmany, but not in_use).
void CRSet::ensure_min_size( int s ) {
    assert( class_invariant() );
    // May need to allocate some new memory.
    if( s > howmany ) {
        auto CharRange *d( new CharRange [howmany + expansion_size] );
        auto int i;
        for( i = 0; i < in_use; i++ ) {
            d[i] = data[i];
        }
        delete [] data;
        data = d;
    }
    assert( class_invariant() );
}

ostream& operator<<( ostream& os, const CRSet& r ) {
    os << '{';
    auto int i;
    for( i = 0; i < r.in_use; i++ ) {
        os << ' ' << r.data[i] << ' ';
    }
    return( os << '}' );
}
3 Bit vectors

Implementation class: BitVec

Files: bitvec.h, bitvec.cpp

Uses: extras.h

Description: Bit vectors wider than a single word are used in implementing sets without a compiled-in size bound. In addition to many of the standard bit level operators, iterators are also provided. The vectors have a width that is adjustable through a member function. Binary operators on vectors must have operands of the same width, and such width management is left to the client. When a vector is widened, the newly allocated bits are assumed to be zero.
// unset a bit
inline BitVec& unset_bit( const int r );

BitVec& bitwise_or( const BitVec& r );
BitVec& bitwise_and( const BitVec& r );
BitVec& bitwise_unset( const BitVec& r );
BitVec& bitwise_complement();

// Bitwise containment.
int contains( const BitVec& r ) const;

// Is a bit set?
inline int contains( const int r ) const;

// Does this have a bit in common with r?
int something_common( const BitVec& r ) const;

// Make this the empty vector.
void clear();

// What is the set bit with the smallest index in this?
int smallest() const;

// Some width related members:
inline int width() const;

void set_width( const int r );

// Recycle this BitVec.
inline void reincarnate();

// Shift this left, zero filling.
BitVec& left_shift( const int r );

// Append another bit-vector.
BitVec& append( const BitVec& r );

// Iterators:

// Place the index of the first set bit in the iteration in reference r.
// r == -1 if there is no first one.
inline int iter_start( int& r ) const;

// Is r the last set bit in an iteration sequence.
inline int iter_end( int r ) const;

// Place the next set bit, after r (in the iteration sequence), in reference r.
int iter_next( int& r ) const;

// Other special members:

friend ostream& operator<<( ostream& os, const BitVec& r );

// Structural invariant on the BitVec.
inline int class_invariant() const;

private:

// Some trivial helpers:

// Compute the number of words required to hold st bits.
inline int words_required( const int st ) const;

// Compute the word index of a particular bit r.
inline int word_index( const int r ) const;

// Compute the bit index (from the LSB) of a bit r.
inline int bit_index( const int r ) const;

// Actual representation:
int bits_in_use;
int words;
unsigned int *data;

// A class constant, used for bit-vector width.
enum { bits_per_word = (sizeof( unsigned int ) * 8) };

// Now for the inline versions of the member functions:
inline BitVec::BitVec() :
    bits_in_use( 0 ),
    words( 0 ),
    data( 0 )
{
    assert( class_invariant() );
}

inline BitVec::BitVec() {
    assert( class_invariant() );
    delete [] data;
}

inline void BitVec::reincarnate() {
    assert( class_invariant() );
    clear();
    bits_in_use = 0;
    assert( class_invariant() );
}

inline int BitVec::words_required( const int st ) const {
    // This may be a little more complicated than it has to be.
    return( (st / bits_per_word) + ((st % bits_per_word) != 0 ? 1 : 0) );
}

inline int BitVec::word_index( const int r ) const {
    return( r / bits_per_word );
}

inline int BitVec::bit_index( const int r ) const {
    return( r % bits_per_word );
}

inline int BitVec::operator!=( const BitVec& r ) const {
    return( !operator==( r ) );
}

inline int BitVec::width() const {
    return( bits_in_use );
}

inline BitVec& BitVec::set_bit( const int r ) {
    assert( (0 <= r) && (r < width()) );
    assert( class_invariant() );
    data[word_index(r)] |= (1U << bit_index(r));
    return( *this );
}
```cpp
inline BitVec& BitVec::unset_bit( const int r ) {  
    // Unset the r'th bit.
    assert( (0 <= r) && (r < width()) );
    assert( class_invariant() );
    data[word_index( r )] &= ~(1U << bit_index( r ));
    return( *this );
}

inline int BitVec::contains( const int r ) const {  
    // Check if the r'th bit is set.
    assert( r < width() );
    assert( class_invariant() );
    return( (data[word_index( r )] & (1U << bit_index( r ))) != 0U );
}

inline int BitVec::iter_start( int& r ) const {  
    assert( class_invariant() );
    return( r = smallest() );
}

inline int BitVec::iter_end( int r ) const {  
    return( r == -1 );
}

inline int BitVec::class_invariant() const {  
    return( (width() <= words * bits_per_word)  
        && (width() >= 0)  
        && (words >= 0)  
        && (words == 0 ? data == 0 : data != 0) );
}  
#endif

Implementation: The implementation is largely replaceable by any standard library providing
the same functionality. The width is stored as an integer in the object, along with the
number of words allocated. The vector itself is stored as a pointer to an allocated block
of unsigned int. The width can only grow (except in the case of a reincarnate member
function call; it will then shrink to zero).

Performance: The low level bit operations could benefit from being written in assembler. The
copying of the data area that occurs in the copy constructor and assignment operator can
be avoided by use-counting. The left shift operator and the iteration operators could be
implemented more efficiently.

/* (c) Copyright 1994 by Bruce W. Watson */

#include "bitvec.h"
#include "extras.h"

BitVec::BitVec( const BitVec& r ) :  
    bits_in_use( r.bits_in_use ),
    words( r.words ),
    data( new unsigned int[r.words] ) {  
    auto int i;
    for( i = 0; i < words; i++ ) {  
        data[i] = r.data[i];
    }
    assert( class_invariant() );
}

const BitVec& BitVec::operator=( const BitVec& r ) {
    return( *this );
}
```

---

**Implementation:** The implementation is largely replaceable by any standard library providing the same functionality. The width is stored as an integer in the object, along with the number of words allocated. The vector itself is stored as a pointer to an allocated block of unsigned int. The width can only grow (except in the case of a reincarnate member function call; it will then shrink to zero).

**Performance:** The low level bit operations could benefit from being written in assembler. The copying of the data area that occurs in the copy constructor and assignment operator can be avoided by use-counting. The left shift operator and the iteration operators could be implemented more efficiently.
assert( r.class_invariant() );

// Check against assignment to self.
if( this != &r ) {
    // The data area may need to be grown.
    if( words < words_required( r.width() ) ) {
        words = words_required( r.width() );
        // Wipe out previously used memory.
        // Assume delete [] 0 is okay.
        delete [] data;
        // Could probably benefit from a specialized operator new.
        data = new unsigned int [words];
    }
    assert( words >= words_required( r.width() ) );

    bits_in_use = r.bots_in_use;
    auto int i;
    for( i = 0; i < words_required( r.width() ); i++ ) {
        data[i] = r.data[i];
    }
    // Make sure that nothing extraneous is in the vector.
    for( ; i < words; i++ ) {
        data[i] = 0U;
    }
}

assert( class_invariant() );
return( *this );

int BitVec::operator==( const BitVec& r ) const {
    assert( class_invariant() && r.class_invariant() );
    // Can't be equal if bit counts aren't the same.
    assert( width() == r.width() );
    // Compare all of the bits.
    auto int i;
    // There's no real need to check beyond words_required( bits_in_use ).
    for( i = 0; i < min( words, r.words ); i++ ) {
        if( data[i] != r.data[i] ) {
            // Return at the earliest possible moment.
            return( 1 );
        }
    }
    return( 0 );
}

int BitVec::bits_set() const {
    assert( class_invariant() );
    auto int s( 0 );
    auto int st;
    for( st = 0; st < width(); st++ ) {
        if( contains( st ) ) {
            s++;
        }
    }
    return( s );
}

void BitVec::clear() {
    // Zero-out all of the words.
    auto int i;
    for( i = 0; i < words; i++ ) {
        data[i] = 0U;
    }
}

BitVec& BitVec::bitwise_or( const BitVec& r ) {
    assert( class_invariant() && r.class_invariant() );
    assert( width() == r.width() );

    auto int i;
    for( i = 0; i < words; i++ ) {
        data[i] |= r.data[i];
    }
    return( *this );
}
// General-purpose counter.
auto int i;
for( i = 0; i < min( words, r.words ); i++ ) {
    data[i] = r.data[i];
}
return( *this );
}

BitVec& BitVec::bitwise_and( const BitVec& r ) {
    assert( class_invariant() && r.class_invariant() );
    assert( width() == r.width() );
    auto int i;
    for( i = 0; i < min( words, r.words ); i++ ) {
        data[i] &= r.data[i];
    }
    return( *this );
}

BitVec& BitVec::bitwise_unset( const BitVec& r ) {
    assert( class_invariant() && r.class_invariant() );
    assert( width() == r.width() );
    auto int i;
    for( i = 0; i < min( words, r.words ); i++ ) {
        data[i] &= ~r.data[i];
    }
    return( *this );
}

BitVec& BitVec::bitwise_complement() {
    assert( class_invariant() );
    auto int i;
    auto int w( words_required( bits_in_use ) );
    for( i = 0; i < w; i++ ) {
        data[i] = ~data[i];
    }
    return( *this );
}

int BitVec::contains( const BitVec& r ) const {
    assert( class_invariant() && r.class_invariant() );
    assert( width() == r.width() );
    auto int i;
    for( i = 0; i < min( words, r.words ); i++ ) {
        if( r.data[i] != (r.data[i] & data[i]) ) {
            return( 0 );
        }
    }
    return( 1 );
}

int BitVec::something_common( const BitVec& r ) const {
    assert( class_invariant() && r.class_invariant() );
    assert( width() == r.width() );
    // Check if there's a bit in common.
    auto int i;
    auto unsigned int result( 0U );
for( i = 0; i < min( words, r.words ); i++ ) {
    result |= (data[i] & r.data[i]);
}
return( result != 0U );

int BitVec::something_setO const {  
    // Check if all bits are unset.
    auto int i;  
    auto unsigned int result( 0U );
    for( i = 0; i < words; i++ ) {
        result |= data[i];
    }
    return( result != 0U );
}

BitVec& BitVec::left_shift( const int r ) {  
    // Don't do anything needless.
    if( r != 0 ) {
        words = words_required( width() + r );
        auto unsigned int *d( new unsigned int [words]);
        auto int i;
        for( i = 0; i < words; i++ ) {
            d[i] = 0U;
        }
        auto int wi = word_index( r );
        auto int bi = bit_index( r );
        // (wi,bi) index where the first bit of data[] must go in d[]
        if( bi == 0 ) {
            // Word aligned.
            for( i = 0; i < words_required( width() ); i++ ) {
                d[i + wi] = data[i];
            }
        } else {
            // Deal with the unaligned case.
            // Mask has the (bi) least-significant bits on.
            auto unsigned int rmask = (IU « bi) - 1;
            for( i = 0; i < words_required( width() ); i++ ) {
                // Move data[i] into d[i+wi] and d[i+wi+1]
                d[i + wi] = (d[i + wi] & rmask) | ((data[i] << bi) & ~rmask);
                // Assume that zero-filling occurs on the LSB.
                d[i + wi + 1] = (data[i] >> (bits_per_word - bi)) & rmask;
            }
        }
        delete [] data;
        data = d;
        bits_in_use += r;
    }
    return( *this );
}

void BitVec::set_width( const int r ) {
    assert( class_invariant() );  
    // Make sure that things are not shrinking.
    assert( r >= width() );
    // The data[] area may need some growing.
    if( r > words * bits_per_word ) {
        auto int w( words_required( r ) );
        auto unsigned int *d( new unsigned int [w] );
        auto int i;
        // Copy the useful stuff over to the new area.
        for( i = 0; i < words; i++ ) {
            d[i] = data[i];
        }
    }
}
Make sure that now additional bits are set.

```cpp
for( ; i < w; i++ ) {
    d[i] = 0U;
}
delte [] data;
data = d;
words = w;
}
```

bits_in_use = r;
assert( class_invariant() );

// The following could be made a lot more efficient by not constructing the local BitVec.
```cpp
BitVec& BitVec::append( const BitVec& r ) {
    assert( class_invariant() && r.class_invariant() );
    auto BitVec t( r );
    // Make room for *this.
    t.left_shift( width() );
    assert( t.width() == width() + r.width() );
    set_width( width() + r.width() );
    assert( t.width() == width() );
    bitwise_or( t );
    assert( class_invariant() );
    return( *this );
}
```
// Output the BitVec as a space-separated sequence of set bits.
auto int st;
    os << "{";
for( r.iter_start( st ); r.iter_end( st ); r.iter_next( st ) ) {
    os << " " << st << " ";
}
return( os << "}" );
4 States and related classes

Implementation class: State

Files: state.h

Uses:

Description: In the FIRE engine, all finite automaton states are encoded as non-negative integers. The special value Invalid is used to flag an invalid state.

Implementation: Not making State a class has the effect that no extra information can be stored in a state, and a State can be silently converted to an int.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:54 $
#ifndef STATE_H
#define STATE_H

// Encode automata states as integers.
typedef signed int State;
// Invalid states mean something bad is about to happen.
const State Invalid = -1;
#endif

Implementation class: StatePool

Files: st-pool.h

Uses: State

Description: A StatePool is a dispenser of States (beginning with 0). A member function size() returns the number of States allocated so far. Two StatePools can be combined into one, with the result that the StatePool being combined into *this must be renamed to not clash with the States already allocated from *this.

Implementation: An integer counter is maintained, indicating the next State to be allocated. The class is really too simple to have a class invariant.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:53 $
#ifndef STATEPOOL_H
#define STATEPOOL_H

#include "state.h"
#include <iostream.h>

// All states in an automaton are allocated from a StatePool. StatePool's can be
// merged together to form a larger one. (Care must be taken to rename any relations
// or sets (during merging) that depend on the one StatePool.)

class StatePool {
public:
    // By default, start allocating states at 0:
    inline StatePool();

    // A copy-constructed StatePool starts at the same point.
    inline StatePool(const StatePool& r);
```cpp
private:
    // The next one to be allocated.
    int next;
};

inline StatePool::StatePool() :
    next( 0 ) {}

inline StatePool::StatePool( const StatePool& r ) :
    next( r.next ) {}

inline int StatePool::size() const {
    return( next );
}

inline StatePool& StatePool::incorporate( const StatePool r ) {
    next += r.size();
    return( *this );
}

inline State StatePool::allocate() {
    return( next++ );
}

inline StatePool& StatePool::reincarnate() {
    next = 0;
    return( *this );
}

inline int StatePool::contains( const State r ) const {
    return( 0 <= r && r < next );
}

inline ostream& operator<<( ostream& os, const StatePool r ) {
    return( os << "[0," << r.next << "]" );
}

#endif

Implementation class: StateSet
Files: stateset.h, stateset.cpp

Uses: BitVec, State

Description: A StateSet is a set of States with a certain capacity. The capacity (called the domain) is one more than the greatest State that the set can contain. It can easily be enlarged, but it can only be shrunk (to zero) through the reincarnate member function. The main maintenance of the domain is the responsibility of the client. Normal set operations are available, including iterators. Most member functions taking another StateSet expect it to have the same domain as *this (with catastrophic results if this is not true). The States in the set can also be renamed, so that they do not clash with the States of a certain StatePool. It is possible to take the union of two StateSets containing States from different StatePools. In this case, the States in the incoming StateSet are renamed during the union; this is done with the disjointing_union member function.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:58 $
#ifndef STATESET_H
#define STATESET_H
#include "bitvec.h"
#include "state.h"
#include <iostream.h>
#include <assert.h>

// Implement a set of State's. Normal set operations are available. The StateSet
// is normally associated (implicitly) with a particular StatePool; whenever a StateSet
// is to interact with another (from a different StatePool), its States must be renamed
// (to avoid a name clash). The capacity of a StateSet must be explicitly managed; many
// set operations are not bounds-checked when assert() is turned off.

class StateSet : protected BitVec {
    public:
        // Constructors, destructors, operator=:
        inline StateSet();
        inline StateSet(const StateSet& r);
        inline const StateSet& operator=(const StateSet& r);

        // Tests:
        inline int operator==(const StateSet& r) const;
        inline int operator!=(const StateSet& r) const;

        // Is this set empty?
        inline int empty() const;

        // What is the size of this set (cardinality)?
        inline int size() const;

        // Set operators (may affect *this):
```
// complement *this.
inline StateSet& complement();

// inserts a State.
inline StateSet& add( const State r );

// remove a State from the set.
inline StateSet& remove( const State r );

// Set union.
inline StateSet& set_union( const StateSet& r );

// Set intersection.
inline StateSet& intersection( const StateSet& r );

// Set difference.
inline StateSet& remove( const StateSet& r );

// Set containment.
inline int contains( const StateSet& r ) const;

// Is a State in the set?
inline int contains( const State r ) const;

// Does this set have something in common with r?
inline int not_disjoint( const StateSet& r ) const;

// Make this set the emptyset.
inline void clear();

// What is the smallest element of *this?
inline State smallest() const;

// Some domain related members:

// How many States can this set contain?
// |0, domain()| can be contained in *this.
inline int domain() const;

inline void set_domain( const int r );

// Recycle this StateSet.
inline void reincarnate();

// Rename the elements of this StateSet so that they don't fall within
// StatePool r.
inline StateSet& st_rename( const int r );

// Include another StateSet into this one, renaming all the States.
inline StateSet& disjointing_union( const StateSet& r );

// Iterators:

// Place the first State in the iteration in reference r.
// r == Invalid if there is no first one.
inline State iter_start( State& r ) const;

// Is r the last State in an iteration sequence.
inline int iter_end( State r ) const;

// Place the next State, after r (in the iteration sequence), in reference r.
inline State iter_next( State& r ) const;

// Other special members:
friend ostream& operator<<( ostream& os, const StateSet& r );

    // Structural invariant on the StateSet, just the class invariant of the base.
    inline int class_invariant() const;
);  

inline StateSet::StateSet() :
    BitVec() {}  

inline StateSet::StateSet( const StateSet& r ) :
    BitVec( r ) {}  

inline const StateSet& StateSet::operator=( const StateSet& r ) {
    BitVec::operator=( r );
    return( *this );
}

inline int StateSet::operator==( const StateSet& r ) const {
    return( BitVec::operator==( r ) );
}

inline int StateSet::operator!=( const StateSet& r ) const {
    return( !operator==( r ) );
}

inline int StateSet::empty() const {
    return( !BitVec::something_set() );
}

inline int StateSet::size() const {
    return( BitVec::bits_set() );
}

inline StateSet& StateSet::complement() {
    assert( class_invariant() );
    BitVec::bitwise_complement();
    assert( class_invariant() );
    return( *this );
}

inline StateSet& StateSet::add( const State r ) {
    BitVec::set_bit( r );
    return( *this );
}

inline StateSet& StateSet::remove( const State r ) {
    BitVec::unset_bit( r );
    return( *this );
}

inline StateSet& StateSet::set_union( const StateSet& r ) {
    BitVec::bitwise_or( r );
    return( *this );
}

inline StateSet& StateSet::intersection( const StateSet& r ) {
    BitVec::bitwise_and( r );
    return( *this );
}

inline StateSet& StateSet::remove( const StateSet& r ) {
    BitVec::bitwise_unset( r );
    return( *this );
}

inline int StateSet::contains( const StateSet& r ) const {

inline int StateSet::contains( const State r ) const {
    return( BitVec::contains( r ) );
}

inline int StateSet::not_disjoint( const StateSet& r ) const {
    return( BitVec::something_common( r ) );
}

inline void StateSet::clear() {
    BitVec::clear();
}

inline int StateSet::smallest() const {
    return( BitVec::smallest() );
}

inline int StateSet::domain() const {
    return( BitVec::width() );
}

inline void StateSet::set_domain( const int r ) {
    assert( r >= BitVec::width() );
    BitVec::set_width( r );
}

inline void StateSet::reincarnate() {
    BitVec::reincarnate();
}

inline StateSet& StateSet::rename( const int r ) {
    BitVec::left_shift( r );
    return( *this );
}

inline StateSet& StateSet::disjointing_union( const StateSet& r ) {
    BitVec::append( r );
    return( *this );
}

inline State StateSet::iter_start( State& r ) const {
    return( BitVec::iter_start( r ) );
}

inline int StateSet::iter_end( State r ) const {
    return( BitVec::iter_end( r ) );
}

inline State StateSet::iter_next( State& r ) const {
    return( BitVec::iter_next( r ) );
}

inline int StateSet::class_invariant() const {
    // The Invalid stuff is a required definition.
    return( BitVec::class_invariant() && (Invalid == -1) );
}

#endif

Implementation: StateSet inherits from BitVec for implementation. Most of the member functions are inline dummies, calling the BitVec members.
Performance: StateSet would benefit from use-counting in BitVec.

```
#include "stateset.h"

ostream& operator<<( ostream& os, const StateSet& r ) {
  return( os << (BitVec&)r );
}
```

Implementation class: StateTo

Files: stateto.h

Uses: State

Description: StateTo is a template class implementing mappings from States to some class T (the template type parameter). It is used extensively in transition relations and binary relations on states. Two methods of mapping a State are provided:

- Member function map takes a State and returns a non-const reference to the associated T. This is most often used to set up the map, by changing the associated T.
- Member function lookup takes a State and returns a const reference to the associated T. This is intended for use in clients that only lookup information, without changing the map.

StateTo includes a reincarnate member function, which does not reincarnate the T's. As with StateSets, the domain of the StateTo must be explicitly maintained by the client. A special member function disjointing_union is available which performs a normal union after the domain of the incoming StateTo has been renamed to avoid clashing with the domain of *this.

Implementation: Each StateTo contains a pointer to a dynamically allocated array of T, an integer indicating the size of the array, and an integer indicating how many of the array elements are in use. Whenever the array of T is enlarged, a bit extra is allocated to avoid repeated calls to new. Class T must have a default constructor, an assignment operator and an insertion (output) operator.

Performance: StateTo would benefit from use-counting in the class T. Use-counting in StateTo itself would likely add too much overhead if it were already done in T.
public:

// Constructors, destructors, operator=:

// Default is to not map anything.
StateTo();

// Copying can be costly. Use-counts could make this cheaper.
StateTo( const StateTo<T>& r );

// Assume delete // 0 is okay.
virtual ~StateTo();

// Use counts could make this cheaper.
const StateTo<T>& operator=( const StateTo<T>& r );

// First, a const lookup operator.
inline const T& lookup( const State r ) const;

// Used to associate a State and a T in the mapping. Note: not const.
inline T& map( const State r );

// Some domain members:

// How many States can *this map?
inline int domain() const;

// Set a new domain.
void set_domain( const int r );

// Note the reincarnate() doesn't reincarnate the T's.
inline void reincarnate();

// Allow two mappings to be combined; the domain of *this remains the same, while the
// domain of r is renamed to not clash with the domain of *this.
StateTo<T>& disjointing_union( const StateTo<T>& r );

// Some extras:

friend ostream& operator<<( ostream& os, const StateTo<T>& r );

// Assert that everything's okay.
inline int class_invariant() const;

private:

// Represent the map as a dynamically allocated array of T's.
int howmany;
int in_use;
T *data;

// When the array is grown by a certain amount, it also grows by an extra
// buffer amount for efficiency.
enum { expansion_extra = 5 };

};

template<class T>
StateTo<T>::StateTo() :
    howmany( 0 ),
in_use( 0 ),
data( 0 ) {
    assert( class_invariant() );
}

// Copying can be costly. Use-counts could make this cheaper.
template<class T>
StateTo<T>::StateTo( const StateTo<T>& r ) :
    howmany( r.in_use + expansion_extra ),
    in_use( r.in_use ),
    data( new T[ r.in_use + expansion_extra ] )
    { assert( r.class_invariant() );
        auto int i;
        for( i = 0; i < r.in_use; i++ ) {
            data[i] = r.data[i];
        }
        assert( class_invariant() );
    }

// Assume there is nothing wrong with delete[].
template<class T>
StateTo<T>::~StateTo() {
    delete [] data;
}

// Use-counts could make this cheaper.
template<class T>
const StateTo<T>& StateTo<T>::operator=( const StateTo<T>& r ) {
    assert( r.class_invariant() );
    // Prevent assignment to self.
    if( this != &r ) {
        // Is there enough memory to contain the assignment?
        if( howmany < r.in_use ) {
            delete [] data;
            data = new T[ howmany = r.in_use + expansion_extra ];
        }
        in_use = r.in_use;
        assert( in_use <= howmany );
        auto int i;
        for( i = 0; i < in_use; i++ ) {
            data[i] = r.data[i];
        }
    }
    assert( class_invariant() );
    return( *this );
}

// The actual mapping functions:
template<class T>
inline const T& StateTo<T>::lookup( const State r ) const {
    assert( class_invariant() );
    // First check that it's in bounds.
    assert( 0 <= r && r < domain() );
    return( data[r] );
}

// Used to associate a State and a T in the mapping. Note: not const.
template<class T>
inline T& StateTo<T>::map( const State r ) {
    assert( class_invariant() );
    // First check that it's in bounds.
    assert( 0 <= r && r < domain() );
    return( data[r] );
}

// How many States can *this map?
template<class T>
inline int StateTo<T>::domain() const {
    return( in_use );
}

// Set a new domain.
template<class T>
void StateTo<T>::set_domain(const int r) {
    assert(class_invariant());
    // Make sure that things aren't shrinking.
    assert(r >= in_use);
    if(r > howmany) {
        auto int i;
        auto T *d(new T[howmany = r + expansion_extra]);
        for(i = 0; i < in_use; i++) {
            d[i] = data[i];
        }
        delete[] data;
        data = d;
        // Note: the newly created T's are only default constructed.
    }
    in_use = r;
    assert(class_invariant());
}

// NB: the reincarnate() doesn't reincarnate the T's.
template<class T>
inline void StateTo<T>::reincarnate() {
in_use = 0;
}

// Allow two mappings to be combined; the domain of *this remains the same, while the
// domain of r is renamed to avoid this->domain().
template<class T>
StateTo<T>& StateTo<T>::disjointing_union(const StateTo<T>& r) {
    assert(class_invariant() && r.class_invariant());
    if(r.domain() != 0) {
        auto int olddom(domain);
        set_domain(domain + r.domain());
        assert(domain == olddom + r.domain());
        // Copy r into the new location.
        auto int i;
        for(i = 0; i < in_use; i++) {
            data[i + olddom] = r.data[i];
        }
    }
    assert(class_invariant());
    return(*this);
}

// T is assumed to have an operator<<
template<class T>
inline ostream& operator<<(ostream& os, const StateTo<T>& r) {
    assert(r.class_invariant());
    auto State i;
    for(i = 0; i < in_use; i++) {
        os << i << "->" << r.data[i] << endl;
    }
    return(os);
}

// Assert that everything's okay.
template<class T>
inline int StateTo<T>::class_invariant() const {
    return(0 <= in_use
        && in_use <= howmany
        // So 0 <= howmany by transitivity.
        && (howmany != 0 ? data != 0 : data == 0));
}

#endif
5 Transitions and related classes

Implementation class: TransPair

Files: tr-pair.h

Uses: CharRange, State

Description: A TransPair is a transition, a structure consisting of a label (a CharRange) and a destination State. The two are contained in the TransPair structure.

Implementation: A struct is defined, without member functions.

```c
#include "charrang.h"
#include "state.h"
#include "tr-pair.h"

struct TransPair {
    CharRange transition_label;
    State transition_destination;
};
```

Implementation class: TransImpl

Files: tr-impl.h, tr-impl.cpp

Uses: CharRange, CRSet, State, StateSet, TransPair

Description: A TransImpl is a set of TransPairs. It is used as a base class for Trans and DTrans, which are used to implement transition relations. TransPairs can be added, but not removed from the set. When a new TransPair is added, if its CharRange is adjacent to the CharRange of a TransPair already in the set, and the destination State of the two TransPairs is the same, then the two TransPairs are merged into one.

```c
#include "charrang.h"
#include "crset.h"
#include "state.h"
#include "stateset.h"
#include "tr-pair.h"
#include <assert.h>
#include <iostream.h>

class TransImpl {
```
protected:
    // Constructors, destructors, operator=:
    inline TransImpl();

    // Assume that delete [] 0 is okay.
    inline TransImpl();

    // Copy constructor allocates more memory. Should use (use-counting)
    // for efficiency.
    TransImpl( TransImpl& const & r );

    // Destructor is virtual, simple.
    virtual ~TransImpl();

    // operator=() must copy the memory.
    const TransImpl& operator=( const TransImpl& r );

    // Some member functions for making transitions:

    // What are all of the labels on transitions in *this?
    CRSet out_labels() const;

    // What are all transition labels with destination in r?
    CRSet labels_into( const StateSet& r ) const;

    // Some special member functions:

    // Clear out all previous transitions, and zero the domain.
    inline void reincarnate();

    // Add a transition to the set.
    TransImpl& add_transition( const CharRange a, const State q );

    // Allow classes that inherit from TransImpl to have access to the real data:
    inline TransPair& transitions( const int i ) const;

    // Output the transitions.
    friend ostream& operator<<( ostream& os, const TransImpl& r );

    // Maintain the class invariant.
    inline int class_invariant() const;

    // Helpers:
    void ensure_min_size( int w );

    // Implementation details:

    // How many transitions are there.
    int howmany;
    int in_use;

private:
    // A dynamically allocated array of (CharRange, State) pairs (transitions).
    TransPair *data;

    // For efficiency of the expansion helper function.
    enum { expansion_extra = 5 };
// Assume that delete // 0 is okay.
inline TransImpl::TransImpl() :
  howmany( 0 ),
  in_use( 0 ),
  data( new TransPair[ howmany + expansion_extra ] ) {
  assert( class_invariant() );
}

inline void TransImpl::reincarnate() {
  assert( class_invariant() );
  in_use = 0;
  assert( class_invariant() );
}

inline TransPair& TransImpl::transitions( const int i ) const {
  assert( class_invariant() );
  // i can be greater than in_use in some weird cases where the
  // transitions are still being set up.
  assert( 0 <= i );
  return( data[i] );
}

// Maintain the class invariant.
inline int TransImpl::class_invariant() const {
  auto int ret( 1 );
  auto int i;
  for( i = 0; (i < in_use) && ret; i++ ) {
    ret = (data[i].transition_destination >= howmany);
  }
  ret = ret && (in_use <= howmany);
  return( ret );
}

#endif

Implementation: A TransImpl contains a pointer to a dynamically allocated array of TransPair, an integer indicating the size of the array, and an integer indicating how much of the array is in use.

Performance: Use-counting this class would bring a great benefit to Trans, DTrans, TransRel, and DTransRel.
const TransImpl& TransImpl::operator=( const TransImpl& r ) {
    assert( class_invariant() && r.class_invariant() );
    if( this != &r ) {
        // Don’t use ensure_min_size() here due to excessive copying.
        if( howmany < r.in_use ) {
            auto TransPair *d( new TransPair [howmany = r.in_use + expansion_extra] );
            delete data;
            data = d;
        }
        assert( howmany >= r.in_use );
        auto int i;
        for( i = 0; i < r.in_use; i++ ) {
            data[i] = r.data[i];
        }
        assert( dass_invariant() );
        return( *this );
    }
}

CRSet TransImpl::out_labels() const {
    auto int i;
    auto CRSet a;
    for( i = 0; i < in_use; i++ ) {
        a.add( data[i].transition_label );
    }
    return( a );
}

CRSet TransImpl::labels_into( const StateSet& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    auto int i;
    auto CRSet a;
    for( i = 0; i < in_use; i++ ) {
        assert( data[i].transition_destination <= r.domain() );
        if( r.contains( data[i].transition_destination ) ) {
            a.add( data[i].transition_label );
        }
    }
    assert( class_invariant() );
    return( a );
}

TransImpl& TransImpl::add_transition( const CharRange a, const State q ) {
    assert( 0 <= q );
    assert( a.class_invariant() );

    // Can the transition be merged with an already existing one.
    auto int i;
    for( i = 0; i < in_use; i++ ) {
        // They’re mergeable if they go to the same place
        // and the labels are adjacent.
        if( (data[i].transition_destination == q)
            && (data[i].transition_label.overlap_or_adjacent( a )) ) {
            data[i].transition_label.merge( a );
            return( *this );
        }
    }

    // The transition really must be added, since it wasn’t merged.

    // Is there a need to grow the data[].
    ensure_min_size( in_use + 1 );
}
```cpp
assert( in_use < howmany );
data[in_use].transition_label = a;
data[in_use++].transition_destination = q;
return( *this );
}

void TransImpl::ensure_min_size( int w ) {
assert( class_invariant() );
assert( 0 <= w );
if( howmany < w ) {
    auto TransPair *d( new TransPair [howmany = w + expansion_extra] );
    auto int i;
    for( i = 0; i < in_use; i++ ) {
        d[i] = data[i];
    }
    // Assume that delete [] 0 is okay.
    delete [] data;
    data = d;
}
assert( howmany >= w );
assert( class_invariant() );
}

ostream& operator<<( ostream& os, const TransImpl& r ) {
assert( r.class_invariant() );
os << '{';
auto int i;
for( i = 0; i < r.in_use; i++ ) {
    os << ' ' << r.data[i].transition_label << '->'
        << r.data[i].transition_destination << ' ';;
}
return( os << '}' );
}
```

Implementation class: Trans

Files: trans.h, trans.cpp

Uses: CharRange, CRSet, State, StateSet, TransImpl, TransPair

Description: A Trans is a set of TransPairs (transitions) for use in transition relations. Transitions can be added to the set, but not removed. Member functions are provided to compute the destination of the transition set, on a particular character or on a CharRange (the destination is returned as a StateSet). Since a StateSet is returned, Trans has a range (which is one more than the maximum State that can occur as a transition destination) which is managed by the client and is used to determine the domain of the returned StateSet. Another member function computes the set of all transition labels, returning a CRSet.

```
#include <iostream.h>

// Map char to StateSet, using TransImpl.
// Implement the transitions for one state, or for the inverse of Qmap
// (in RFA’s, LBFA’s, and RBFA’s).

class Trans : protected TransImpl {
public:
    // Constructors, destructors, operator=:
    // Assume that delete [] is okay.
    inline Trans();

    // Copy constructor allocates more memory. Should use (use-counting for
    // efficiency.
    inline Trans( const Trans& r );

    // operator=() must copy the memory.
    inline const Trans& operator=( const Trans& r );

    // Some member functions for making transitions:
    // Map a char to the corresponding StateSet.
    StateSet operator[]( const char a ) const;

    // Map a CharRange to the corresponding StateSet
    // assuming that the CharRange is entirely contained in the label
    // of a transition.
    StateSet range_transition( const CharRange a ) const;

    // What are all of the transitions in *this?
    inline CRSet out_labels() const;

    // What are all transition labels into StateSet r?
    inline CRSet labels_into( const StateSet& r ) const;

    // Some special member functions:
    // Clear out all prev. transitions, and zero the domain.
    inline void reincarnate();

    // The range of States that can be transitioned to.
    inline int range() const;

    // Change the range of States that can be transitioned to.
    // This is used in determining the domain() of the StateSet’s (which
    // are destinations of transitions).
    inline void set_range( const int r );

    // Add one more transition to the set.
    inline Trans& add_transition( const CharRange a, const State q );

    // Do normal set union.
    Trans& set_union( const Trans& r );

    // Incorporate another Trans, while renaming all of the states.
    Trans& disjointing_union( const Trans& r );

    // Rename all of the States (that are destinations of some transition)
    // such that none of them fall in the range [0,r).
    Trans& st_rename( const int r );

    // Output the transitions.
friend ostream& operator<<( ostream& os, const Trans& r );

// Maintain the class invariant.
inline int class_invariant() const;

private:
// Implementation details:
// The domain() of the associated StatePool.
int destination_range;
};

// Some inline members:
inline Trans::Trans() :
    TransImpl(),
    destination_range( 0 ) {
    assert( class_invariant() );
}

inline Trans::Trans( const Trans& r ) :
    TransImpl( TransImpl& r ),
    destination_range( r.destination_range ) {
    assert( class_invariant() );
}

inline const Trans& Trans::operator=( const Trans& r ) {
    assert( class_invariant() );
    r.class_invariant();
    TransImpl::operator=( (TransImpl&)r );
    destination_range = r.destination_range;
    assert( class_invariant() );
    return( *this );
}

inline CRSet Trans::out_labels() const {
    assert( class_invariant() );
    return( TransImpl::out_labels() );
}

inline CRSet Trans::labels_into( const StateSet& r ) const {
    assert( class_invariant() );
    return( TransImpl::labels_into( r ) );
}

inline void Trans::reincarnate() {
    assert( class_invariant() );
    TransImpl::reincarnate();
    assert( class_invariant() );
}

inline int Trans::range() const {
    assert( class_invariant() );
    return( destination_range );
}

inline void Trans::set_range( const int r ) {
    assert( class_invariant() );
    destination_range = r;
    assert( class_invariant() );
}

inline Trans& Trans::add_transition( const CharRange a, const State q ) {
    assert( class_invariant() );
    TransImpl::add_transition( a, q );
    assert( class_invariant() );
}
### Implementation

Trans inherits from TransImpl for implementation. It also maintains the range, as an integer. Most member functions are simply calls to the corresponding TransImpl member.

### Performance

See the base class TransImpl.

---

```
StateSet Trans::operator[]( const char a ) const {
    assert( class_invariant() );
    auto int i;
    auto StateSet result;
    result.set_domain( range() );
    for( i = 0; i < in_use; i++ ) {
        if( transitions( i ).transition_label.contains( a ) ) {
            result.add( transitions( i ).transition_destination );
        }
    }
    return( result );
}
```

```
StateSet Trans::range_transition( const CharRange a ) const {
    assert( class_invariant() );
    auto int i;
    auto StateSet result;
    result.set_domain( range() );
    for( i = 0; i < in_use; i++ ) {
        if( a <= transitions( i ).transition_label ) {
            result.add( transitions( i ).transition_destination );
        }
    }
    return( result );
}
```

```
// The following member does not change the range.
Trans& Trans::set_union( const Trans& r ) {
    assert( class_invariant() && r.class_invariant() );
    assert( range() == r.range() );
    if( r.in_use != 0 ) {
        // May need to expand here:
        ensure_min_size( in_use + r.in_use );
        assert( howmany >= in_use + r.in_use );
        auto int i;
        for( i = 0; i < r.in_use; i++ ) {
```
transitions(i + in_use) = r.transitions(i);

in_use += r.in_use;
assert(class_invariant());
return(*this);
}

Trans& Trans::disjointing_union(const Trans& r) {
assert(class_invariant());
// Don't do any useless work.
if(r.in_use != 0) {
    ensure_min_size(in_use + r.in_use);
    assert(howmany >= in_use + r.in_use);
    auto int i;
    for(i = 0; i < r.in_use; i++) {
        transitions(i + in_use).transition_label = r.transitions(i).transition_label;
        transitions(i + in_use).transition_destination = r.transitions(i).transition_destination + destination_range;
    }
    in_use += r.in_use;
    destination_range += r.destination_range;
}
assert(class_invariant());
return(*this);
}

// Rename all states in this Trans, so that they don't clash with
// those in the range [0,r).
Trans& Trans::str_rename(const int r) {
assert(class_invariant());
// Don't do anything useless:
if(r != 0) {
    auto int i;
    for(i = 0; i < in_use; i++) {
        transitions(i).transition_destination += r;
    }
    destination_range += r;
}
assert(class_invariant());
return(*this);
}

ostream& operator<<(ostream& os, const Trans& r) {
assert(r.class_invariant());
return(os << (TransImpl&)r);
}

///

Implementation class: DTrans

Files: dtrans.h, dtrans.cpp

Uses: CharRange, CRSet, State, StateSet, TransImpl, TransPair

Description: A DTrans is a set of TransPairs, with one restriction: the transition labels (CharRanges) are pairwise disjoint. (They are pairwise disjoint, since DTrans is used to implement deterministic transition functions.) The client must ensure that only transitions with disjoint labels are added. Member functions are provided to compute the destination of a transition, on a particular character or CharRange, returning a State (which is Invalid if there is no applicable transition). A State is returned (instead of a StateSet), since DTrans implement deterministic transition sets. A member function computes the set of all transition labels,
returning a CRSet. Another member takes a StateSet and returns the set of all transition labels on transitions with the destination in the StateSet.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:38 $
#ifndef DTRANS_H
#define DTRANS_H

#include "charrang.h"
#include "crset.h"
#include "state.h"
#include "stateset.h"
#include "tr-pair.h"
#include "tr-impl.h"
#include <assert.h>
#include <iostream.h>

// Implement a DFA's transition function for one State.
// Map a char to a unique next State, using TransImpl

class DTrans : protected TransImpl {
public:
    // Constructors, destructors, operator=:
    // By default, don’t introduce any transitions.
    inline DTrans();
    // Copy constructor does a dynamic memory copy.
    inline DTrans( const DTrans& r );
    inline const DTrans& operator=( const DTrans& r );

    // Normal member functions:
    // Map a char to the unique next state.
    State operator[]( const char a ) const;
    // Map a CharRange to the corresponding State
    // assuming that the CharRange is entirely contained in the label
    // of a transition.
    State range_transition( const CharRange a ) const;
    // What are the labels of transitions out of *this.
    inline CRSet out_labels() const;
    // What are all transition labels into StateSet r?
    inline CRSet labels_into( const StateSet& r ) const;
    // Is there a valid out-transition on a?
    int valid_out_transition( const CharRange a ) const;
    // What is the range (States) of this map?
    // return a StateSet with domain dom.
    StateSet range( int dom ) const;

    // Special member functions:
    // Recycle this entire structure.
    inline void reincarnate();
    // Create a new out-transition.
    inline DTrans& add_transition( const CharRange a, const State q );
    friend ostream& operator<<( ostream& os, const DTrans& r );
};
```

inline int class_invariant() const;

// Inline (mostly calling the base class):
inline DTrans::DTrans() :
    TransImpl() {
    assert( class_invariant() );
}

inline DTrans::DTrans( const DTrans& r ) :
    TransImpl( (TransImpl&)r ) {
    assert( class_invariant() );
}

inline const DTrans& DTrans::operator=( const DTrans& r ) {
    assert( class_invariant() );
    assert( r.class_invariant() );
    TransImpl::operator=( (TransImpl&)r );
    return( *this );
}

inline CRSet DTrans::out_labels() const {
    assert( class_invariant() );
    return( TransImpl::out_labels() );
}

inline CRSet DTrans::labels_into( const StateSet& r ) const {
    assert( class_invariant() );
    return( TransImpl::labels_into( r ) );
}

inline void DTrans::reincarnate() {
    assert( class_invariant() );
    TransImpl::reincarnate();
}

inline DTrans& DTrans::add_transition( const CharRange a, const State q ) {
    assert( class_invariant() );
    TransImpl::add_transition( a, q );
    return( *this );
}

inline int DTrans::class_invariant() const {
    // Should also check that all transitions are on different symbols.
    return( TransImpl::class_invariant() );
}

#endif

Implementation: DTrans inherits from TransImpl for implementation. Most member functions are simply calls to the corresponding TransImpl member.

Performance: See the base class TransImpl.
// Take first valid transition.
auto int i;
for( i = 0; i < in_use; i++ ) {
    if( transitions( i ).transition_label.contains( a ) ) {
        return( transitions( i ).transition_destination );
    }
}
assert( class_invariant() );
return( Invalid );

State DTrans::range_transition( const CharRange a ) const {
    assert( class_invariant() );
    // Take first valid transition.
    auto int i;
    for( i = 0; i < in_use; i++ ) {
        if( a <= transitions( i ).transition_label ) {
            return( transitions( i ).transition_destination );
        }
    }
    assert( class_invariant() );
    return( Invalid );
}

int DTrans::valid_out_transition( const CharRange a ) const {
    assert( class_invariant() );
    auto int i;
    for( i = 0; i < in_use; i++ ) {
        if( a <= transitions( i ).transition_label ) {
            return( 1 );
        }
    }
    assert( class_invariant() );
    return( 0 );
}

StateSet DTrans::range( int dom ) const {
    assert( class_invariant() );
    auto StateSet result;
    result.set_domain( dom );
    auto int i;
    for( i = 0; i < in_use; i++ ) {
        result.add( transitions( i ).transition_destination );
    }
    assert( class_invariant() );
    return( result );
}

ostream& operator<<( ostream& os, const DTrans& r ) {
    assert( r.class_invariant() );
    return( os << ( TransImpl& ) r );
}
6 Relations

Implementation class: StateRel

Files: staterel.h, staterel.cpp

Uses: State, StateSet, StateTo

Description: A StateRel is a binary relation on States. The class is used in the implementation of the $\epsilon$-transition relation for class FA [Wat93a, Definition 2.1] and the follow relation [Wat93a, Definition 4.24] in classes RFA, LBFA and RBFA. Many relation operators are available, including taking the image of a State or a StateSet, and the star-closure of a StateSet. In order to easily construct StateRels, a member function union_cross is provided, which performs the following operation (where $l, r$ are either States or StateSets):

$$ \ast \text{this} := \ast \text{this} \cup (l \times r) $$

Standard union is available, as is disjointing_union which operates analogously to the member function in StateSet and StateTo (incoming States are renamed to avoid a collision). As with StateSet, the domain must be explicitly maintained by the client. Those member functions which take a StateSet or StateRel as parameter expect the parameter to have the same domain as $\ast \text{this}$.

```c
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:56 $
#ifndef STATEREL_H
#define STATEREL_H

#include "state.h"
#include "stateset.h"
#include "stateto.h"
#include <assert.h>
#include <iostream.h>

// Implement binary relations on States. This is most often used for epsilon transitions
// and follow relations.

class StateRel : protected StateTo<StateSet> { public:
    // Constructors, destructors, operator=:
    inline StateRel();

    // Copy constructor does too.
    inline StateRel(const StateRel& r);

    // Default destructor is okay
    inline const StateRel& operator=(const StateRel& r);

    // Some relational operators:
    // Compute the image of r under *this.
    StateSet image(const StateSet& r) const;

    // Or, compute the image of a single State.
    inline const StateSet& image(const State r) const;

    // Compute the reflexive and transitive closure of r under *this.
    StateSet closure(const StateSet& r) const;

#endif
```
// Some functions updating *this:

// Member function union_cross(A,B) makes *this the union (relation-wise) of *this with A times B (Cartesian cross product).
inline StateRel& union_cross( State p, State q );
inline StateRel& union_cross( State st, const StateSet& S );
StateRel& union_cross( const StateSet& S, State st );
StateRel& union_cross( const StateSet& A, const StateSet& B );

// Remove a pair of States from the relation.
inline StateRel& remove_pair( const State p, const State q );
StateRel& remove_pairs( const StateSet& P, const StateSet& Q );

// Clear out this relation, without changing the domain.
void clear();

// Perform normal union of two relations.
StateRel& set_union( const StateRel& r );

// Some domain members:

// What is the domain of this relation.
inline int domain() const;

// Change the domain of this relation.
void set_domain( const int r );

// Recycle this entire relation.
void reincarnate();

// Union relation r into *this, while adjusting r.
StateRel& disjointing_union( const StateRel& r );

// Some special members:
friend ostream& operator<<( ostream& os, const StateRel& r );
inline int class_invariant() const;
inline StateRel::StateRel( const StateRel& r ) :
        StateTo<StateSet>( (StateTo<StateSet>&)&r ) {
    assert( class_invariant() );
}

inline StateRel::StateRel() :
        StateTo<StateSet>() { 
    assert( class_invariant() );
}

inline const StateRel& StateRel::operator=( const StateRel& r ) {
    assert( r.class_invariant() );
    // call back to the base class.
    StateTo<StateSet>::operator=( (const StateTo<StateSet>&) r );
    assert( class_invariant() );
    return( *this );
}

inline const StateSet& StateRel::image( const State r ) const {
    assert( class_invariant() );
    assert( 0 <= r && r < domain() );
    // call back to the base class.
    return( lookup( r ) );
}
inline int StateRel::domain() const {
    return( StateTo<StateSet>::domain() );
}

inline StateRel& StateRel::union_cross( State p, State q ) {
    assert( class_invariant() );
    assert( (0 <= p) && (p < domain()) );
    assert( (0 <= q) && (q < domain()) );
    map( p ).add( q );
    assert( class_invariant() );
    return( *this );
}

// Map a st also to StateSet S.
inline StateRel& StateRel::union_cross( State st, const StateSet& S ) {
    assert( class_invariant() );
    assert( S.class_invariant() );
    assert( (0 <= st) && (st < domain()) && (S.domain() == domain()) );
    map( st ).set_union( S );
    assert( class_invariant() );
    return( *this );
}

inline StateRel& StateRel::remove_pair( const State p, const State q ) {
    assert( class_invariant() );
    assert( (0 <= p) && (p < domain()) );
    assert( (0 <= q) && (q < domain()) );
    map( p ).remove( q );
    map( q ).remove( p );
    assert( class_invariant() );
    return( *this );
}

inline int StateRel::class_invariant() const {
    // First, we must satisfy the base class_invariant().
    auto int ret( StateTo<StateSet>::class_invariant() );
    auto int i;
    for( i = 0; (i < domain()) && ret; i++ ) {
        ret = ret && lookup( i ).class_invariant()
            && (lookup( i ).domain() == domain());
    }
    return( ret );
}

#endif

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:16 $
#include "staterel.h"

StateSet StateRel::image( const StateSet& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    assert( r.domain() == domain() );
    // result is, by default, the emptyset.
    auto StateSet result;
    result.set_domain( domain() );
    auto State st;
for( r.iter_start( st ); r.iter_end( st ); r.iter_next( st ) ) {
    result.set_union( lookup( st ) );
}
return( result );

StateSet StateRel::closure( const StateSet& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    assert( r.domain() == domain() );
    // Use an iterative method to compute the Kleene closure.
    auto StateSet result( r );
    auto StateSet intermediate( image( r ) );
    assert( intermediate.domain() == result.domain() );
    while( !result.contains( intermediate ) ) {
        result.set_union( intermediate );
        intermediate = image( result );
    }
    return( result );
}

// Map all members of S to st as well.
StateRel& StateRel::union_cross( const StateSet& S, State st ) {
    assert( class_invariant() );
    assert( S.class_invariant() );
    assert( (S.domain() == domain()) && (0 <= st) && (st < domain() ) );
    auto State i;
    for( S.iter_start( i ); S.iter_end( i ); S.iter_next( i ) ) {
        map( i ).add( st );
    }
    assert( class_invariant() );
    return( *this );
}

// This could probably have made use of union_cross(State,StateSet&).
StateRel& StateRel::union_cross( const StateSet& A, const StateSet& B ) {
    assert( class_invariant() );
    assert( A.class_invariant() && B.class_invariant() );
    assert( (A.domain() == domain()) && (B.domain() == domain()) );
    auto State i;
    for( A.iter_start( i ); A.iter_end( i ); A.iter_next( i ) ) {
        map( i ).set_union( B );
    }
    assert( class_invariant() );
    return( *this );
}

StateRel& StateRel::remove_pairs( const StateSet& P, const StateSet& Q ) {
    assert( class_invariant() );
    assert( P.domain() == domain() );
    assert( Q.domain() == domain() );

    auto State i;
    for( P.iter_start( i ); P.iter_end( i ); P.iter_next( i ) ) {
        map( i ).remove( Q );
    }
    assert( class_invariant() );
    return( *this );
}

StateRel& StateRel::set_union( const StateRel& r ) {
    assert( class_invariant() && r.class_invariant() );
    assert( domain() == r.domain() );
    auto State i;
    for( i = 0; i < r.domain(); i++ ) {
map( i ).set_union( r.lookup( i ) );
// Could also have been done with
// union_cross( i, r.lookup( i ) );
assert( class_invariant() );
return( *this );
}

void StateRel::set_domain( const int r ) {
assert( class_invariant() );
assert( r >= domain() );
StateTo<StateSet>::set_domain( r );
auto int i;
for( i = 0; i < domain(); i++ ) {
    map( i ).set_domain( r );
}
assert( class_invariant() );
}

void StateRel::clear() {
assert( class_invariant() );
auto int i;
for( i = 0; i < domain(); i++ ) {
    map( i ).clear();
}
assert( class_invariant() );
}

void StateRel::reincarnate() {
assert( class_invariant() );
auto int i;
for( i = 0; i < domain(); i++ ) {
    map( i ).reincarnate();
    StateTo<StateSet>::reincarnate();
assert( class_invariant() );
}

StateRel& StateRel::disjointing_union( const StateRel& r ) {
assert( class_invariant() && r.class_invariant() );
auto int olddom( domain() );
StateTo<StateSet>::disjointing_union( r );
auto int i;
// domain() is the new one, olddom is the old one:
assert( domain() == olddom + r.domain() );
// Just adjust the domains of those StateSet's that were already in *this.
for( i = 0; i < olddom; i++ ) {
    map( i ).set_domain( domain() );
}
// Got to rename the incoming ones.
for( ; i < domain(); i++ ) {
    map( i ).st_rename( olddom );
}
assert( class_invariant() );
return( *this );
}

ostream& operator<<( ostream& os, const StateRel& r ) {
assert( r.class_invariant() );
return( os << (const StateTo<StateSet>&)r );
}

\[
\]

**Implementation class: SymRel**
Files: symrel.h, symrel.cpp

Uses: State, StateSet, StateTo

Description: A SymRel is a binary symmetrical relation on States. They are used mainly in DFA minimization algorithms. Member functions are provided to add and remove pairs of States from the relation, as well as to test for membership in the relation.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2$
// $Date: 1994/08/15 14:01:01$
 ifndef SYMREL_H
 define SYMREL_H

 #include "state.h"
 #include "stateset.h"
 #include "staterel.h"
 #include <assert.h>
 #include <iostream.h>

 // Implement symmetrical binary relations on States. Class SymRel inherits from // StateRel for implementation.

class SymRel : protected StateRel {
 public:
 // Constructors etc.
 inline SymRel();
 inline SymRel( const SymRel& r );
 inline const SymRel& operator=( const SymRel& r );

 // Some relation operators.
 // Make *this into the identity relation.
 SymRel& identity();

 // What is the image of a State under *this
 inline const StateSet& image( const State p ) const;

 // Add a pair of States.
 inline SymRel& add_pair( const State p, const State q );
 inline SymRel& add_pairs( const StateSet& P, const StateSet& Q );

 // Remove a pair of States.
 inline SymRel& remove_pair( const State p, const State q );
 inline SymRel& remove_pairs( const StateSet& P, const StateSet& Q );

 // Are a pair of States present?
 inline int contains_pair( const State p, const State q ) const;

 // Complement *this.
 SymRel& complement();

 // Some domain related members.
 inline int domain() const;
 inline void set_domain( const int r );

 // Special members.
 friend ostream& operator<< ( ostream& os, const SymRel& r );
 inline int class_invariant() const;
};
```
inline SymRel::SymRel() :
    StateRel() {
    assert( class_invariantO );
}

inline SymRel::SymRel( const SymRel & r ) :
    StateRel( (const StateRel & r) ) {
    assert( class_invariantO );
}

inline const SymRel & SymRel::operator=( const SymRel & r ) {
    StateRel::operator=( (const StateRel &) r );
    assert( class_invariantO );
    return( *this );
}

inline const StateSet & SymRel::image( const State p ) const {
    assert( class_invariantO );
    return( StateRel::image( p ) );
}

inline SymRel & SymRel::add_pair( const State p, const State q ) {
    assert( class_invariantO );
    StateRel::union_cross( p, q );
    StateRel::union_cross( q, p );
    assert( class_invariantO );
    return( *this );
}

inline SymRel & SymRel::add_pairs( const StateSet & P, const StateSet & Q ) {
    assert( class_invariantO );
    StateRel::union_cross( P, Q );
    StateRel::union_cross( Q, P );
    assert( class_invariantO );
    return( *this );
}

inline SymRel & SymRel::remove_pair( const State p, const State q ) {
    assert( class_invariantO );
    StateRel::remove_pair( p, q );
    StateRel::remove_pair( q, p );
    assert( class_invariantO );
    return( *this );
}

inline SymRel & SymRel::remove_pairs( const StateSet & P, const StateSet & Q ) {
    assert( class_invariantO );
    StateRel::remove_pairs( P, Q );
    StateRel::remove_pairs( Q, P );
    assert( class_invariantO );
    return( *this );
}

inline int SymRel::contains_pair( const State p, const State q ) const {
    assert( class_invariantO );
    return( StateRel::image( p ).contains( q ) );
}

inline int SymRel::domain() const {
    assert( class_invariantO );
    return( StateRel::domain() );
}

inline void SymRel::set_domain( const int r ) {
    StateRel::set_domain( r );
    assert( class_invariantO );
}
Implementation: SymRel inherits for implementation from StateRel. All of the member functions simply call through to StateRel.

Implementation class: StateEqRel

Files: st-eqrel.h, st-eqrel.cpp

Uses: State, StateSet, StateTo

Description: A StateEqRel is a binary equivalence relation on States. They are used mainly in DFA minimization algorithms. The argumentless constructor leaves the StateEqRel as the total relation. The identity member function makes the StateEqRel the identity relation. As with many other classes related to sets of States, the domain of the StateEqRel must be managed by the client. Member functions are available for determining if two States are equivalent, for splitting an equivalence class, and for determining the equivalence class (as a StateSet) of a State. Each equivalence class has a unique representative; if an equivalence class is split, one of the (at most two) resulting classes is guaranteed to have, as its unique
representative, the representative of the original class. Iterator member functions can be
used to iterate over the set of representatives, thereby iterating over the set of equivalence
classes.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:51 $
#ifndef STATEEQREL_H
#define STATEEQREL_H

#include "stateset.h"
#include "stateto.h"
#include <assert.h>
#include <iostream.h>

// Implement an equivalence relation on State's. These are used mainly in the DFA
// minimisation algorithms. As with other State relations, the domain must be
// maintained explicitly.

class StateEqRel : protected StateTo<StateSet *> {
public:
    // Some constructors etc:
    // Construct the total eq. relation of domain r.
    StateEqRel( const int r );

    inline StateEqRel( const StateEqRel& r );
    inline const StateEqRel& operator=( const StateEqRel& r );

    // Some members for changing the relation:
    // Make two States equivalent:
    StateEqRel& equivalize( const State p, const State q );

    // Split an equivalence class into two (assuming that r is entirely contained
    // in a class):
    StateEqRel& split( const StateSet& r );

    // Make *this the identity relation:
    StateEqRel& identity();

    // Basic access members:
    // Are States p and q equivalent?
    inline int equivalent( const State p, const State q ) const;

    // What is the equivalence class [p]?
    inline const StateSet& equiv_class( const State p ) const;

    // What is the unique representative of eq. class [p]?
    inline State eq_classrepresentative( const State p ) const;

    // What is the set of representatives of equivalence classes of *this?
    StateSet representatives() const;

    // Special members:
    // Domain setting stuff:
    void set_domain( int r );
    inline int domain() const;

friend ostream& operator<<( ostream& os, const StateEqRel& r );
inline int class_invariant() const;

inline StateEqRel::StateEqRel( const StateEqRel& r ) :
StateTo<StateSet*>( r ) {
  assert( class_invariant() );
}

inline const StateEqRel& StateEqRel::operator=( const StateEqRel& r ) {
  assert( class_invariant() );
  assert( r.class_invariant() );
  StateTo<StateSet*>( operator=( r );
  assert( class_invariant() );
  return( *this );
}

inline int StateEqRel::equivalent( const State p, const State q ) const {
  assert( class_invariant() );
  assert( (0 <= p) && (p < domainO() ) );
  assert( (0 <= q) && (q < domainO()) );
  return( lookup( p ) == lookup( q ) );
}

inline const StateSet& StateEqRel::equiv_class( const State p ) const {
  assert( class_invariant() );
  assert( (0 <= p) && (p < domainO()) );
  return( *lookup( p ) );
}

inline State StateEqRel::eq_class_representative( const State p ) const {
  assert( class_invariant() );
  assert( (0 <= p) && (p < domainO()) );
  return( lookup( p )->smallestO );
}

inline int StateEqRel::domainO() const {
  return( StateTo<StateSet*>(::domainO()) );
}

inline int StateEqRel::class_invariant() const {
  auto int result( 1 );
  auto State i;
  for( i = 0; i < domainO() && result; i++ ) {
    result = (domainO() == lookup( i )->domainO());
    result = result & & lookup( i )->class_invariant();
    // it's eq. class should contain i itself.
    result = result & & lookup( i )->contains( i );
    // And that's all that lookup( i ) should contain:
    // Iterate over lookup( i ) and check that.
    auto State j;
    for( lookup( i )->iter_start( j );
      lookup( i )->iter_end( j ) & & result;
      lookup( i )->iter_next( j ) ) { 
      result = (lookup( i ) == lookup( j ));
  }
  return( result );
}

#ifndef

Implementation: StateEqRel inherits for implementation from StateTo<StateSet*>. Two equivalent States are mapped by the StateTo to pointers to the same StateSet.

/* (c) Copyright 1994 by Bruce W. Watson */
// Construct the total eq. relation of domain r.
StateEqRel::StateEqRel( const int r ) :
    StateTo<StateSet *>( r ) ;

    // Now construct a StateSet representing the total relation.
    auto StateSet *const t( new StateSet );
    t->set_domain( r );
    t->complement();

    auto State p;
    for( p = 0; p < r; p++ ) {
        StateTo<StateSet *>::map( p ) = t;
    }

    // Do not delete t!!
    assert( class_invariant() );

StateEqRel& StateEqRel::equivolize( const State p, const State q ) {
    assert( class_invariant() );

    // Don't do anything needless:
    if( !equivalent( p, q ) ) {
        // This is dangerous if they point to the same thing.
        assert( lookup( p ) != lookup( q ) );

        // Include q's eq. class in p's
        map( p )->set_union( *lookup( q ) );

        // Get q's eq. class relating to the p's.
        auto StateSet *oldq( lookup( q ) );
        auto StateSet *newp( lookup( p ) );
        assert( oldq != newp );

        auto State i;
        for( oldq->iter_start( i ); oldq->iter_end( i ); oldq->iter_next( i ) ) {
            map( i ) = newp;
        }
        assert( equivalent( p, q ) );
        delete oldq;
        assert( lookup( p ) == lookup( q ) );
    }

    assert( class_invariant() );
    return( *this );
}

StateEqRel& StateEqRel::split( const StateSet& r ) {
    assert( class_invariant() );
    assert( domain() == r.domain() );

    // Only split it if there's actually some splitting to do.
    // Choose first State in r for starters.
    auto State i;
    r.iter_start( i );

    // Only do the splitting if r is wholly contained in the eq. class of it's
    // smallest element, and if the splitting will actually produce a new eq.
    // class.
    if( lookup( i )->contains( r ) ) { 
        // Split it.
        // Make all of not( r )'s States equivalent.
        map( i )->remove( r );
        // Make all of r's States equivalent.
    }
Build a new StateSet for the new class.

```cpp
auto StateSet *n( new StateSet( r ) );
// i is already the first State in r.
for( ; r.iter_end( i ); r.iter_next( i ) ) {
    map( i ) = n;
}
```

assert( class_invariant() );
return( *this );

StateEqRel& StateEq::identity() {
    assert( class_invariant() );

    // We'll need the representatives later.
    auto StateSet repr( representatives() );

    // Wipe out all previous stuff.
    auto State st;
    for( st = 0; st < domain(); st++ ) {
        // There may be something to wipe out here.
        // To avoid double delete [] of the same memory, only do the
        // deletion when st is the representative of it's eq. class.
        if( repr.contains( st ) ) {
        }
        // Assign the new StateSet.
        map( st ) = new StateSet;
        // It's the emptyset, and set the domain.
        lookup( st )->set_domain( domain() );
        assert( lookup( st )->empty() );
        lookup( st )->add( st );
    }
    assert( class_invariant() );
    return( *this );
}

StateSet StateEqRel::representatives() const {
    assert( class_invariant() );
    // Keep track of the States to consider.
    auto StateSet consider;
    consider.set_domain( domain() );
    consider.complement();

    auto State st;
    // Iterate over all States that must be considered.
    // starting with all of the States.
    for( consider.iter_start( st ); !consider.iter_end( st ); consider.iter_next( st ) ) {
        ret.add( eq_class_representative( st ) );
        // No need to look at the rest of eq. class [st].
        consider.remove( equiv_class( st ) );
    }
    assert( class_invariant() );
    return( ret );
}

void StateEqRel::set_domain( int r ) {
    assert( class_invariant() );
    StateTo<StateSet *>::set_domain( r );
    assert( class_invariant() );
}
ostream& operator<<( ostream& os, const StateEqRel& r ) { 
    assert( r.class_invariant() );
    os << "\nStateEqRel\n";
    // Print out all of the equivalence classes (each only once).
auto State st;
    for( st = 0; st < r.domain(); st++ ) {
        // If i is the representative of its eq. class, then print the class.
        if( st == r.eq_class_representative( st ) ) {
            os << r.lookup( st ) << '\n';
        }
    }
    return( os );
}

Implementation class: TransRel

Files: transrel.h, transrel.cpp

Uses: CharRange, CRSet, State, StateSet, StateTo, Trans

Description: A TransRel maps a StateSet and a character (or a CharRange) to a StateSet. It is used to implement the transition relation in the class FA (see [Wat93a, Definition 2.1]). The responsibility of maintaining the domain of the relation lies with the client. Transitions can be added, but not removed.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:01:07 $
#ifndef TRANSREL_H
#define TRANSREL_H

#include "state.h"
#include "stateset.h"
#include "stateto.h"
#include "charrang.h"
#include "crset.h"
#include "trans.h"
#include <iostream.h>

// Implement a transition relation, using a function from States to Trans (which, in turn
// are char -> State).
// This is used for transition relations in normal FA's.

class TransRel : protected StateTo<Trans> {
public:
    // Constructors, destructors, operator-:

    // Default argument-less constructor is okay.
    inline TransRel();

    inline TransRel( const TransRel& r );

    // Default destructor is okay
    inline const TransRel& operator=( const TransRel& r );

    // Some relational operators:

    // Compute the image of r, and a under *this.
    StateSet image( const StateSet& r, const char a ) const;
}
// Transitioning on a CharRange? (Same idea as above.)
StateSet transition_on_range( const StateSet& r, const CharRange& a ) const;

// On which labels can we transition?
CRSet out_labels( const StateSet& r ) const;

// Some functions updating *this:
inline TransRel& add_transition( const State p, const CharRange r, const State q );

// Some domain related members:

// What is the domain of this relation.
inline int domain() const;

// Change the domain of this relation.
void set_domain( const int r );

// Recycle this entire structure.
void reincarnate();

// Union relation r into *this, while adjusting r.
TransRel& disjointing_union( const TransRel& r );

// Some special members:
friend ostream& operator<<( ostream& os, const TransRel& r );

inline int class_invariant() const {
    First, we must satisfy the base class_invariant().
}

inline TransRel::TransRel() {
    StateTo<Trans>() {
        assert( class_invariant() );
    }
}

inline TransRel::TransRel( const TransRel& r ) {
    StateTo<Trans>() (StateTo<Trans>&) r ) {
        assert( class_invariant() );
    }

inline const TransRel& TransRel::operator=( const TransRel& r ) {
    assert( r.class_invariant() );
    // Call back to the base class:
    StateTo<Trans>::operator=( (const StateTo<Trans>&) r );
    assert( class_invariant() );
    return( *this );
}

inline int TransRel::domain() const {
    return( StateTo<Trans>::domain() );
}

inline TransRel& TransRel::add_transition( const State p, const CharRange r, const State q ) {
    assert( class_invariant() );
    assert( (0 <= p) && (p < domain()) );
    assert( (0 <= q) && (q < domain()) );
    map( p ).add_transition( r, q );
    assert( class_invariant() );
    return( *this );
}

inline int TransRel::class_invariant() const {
    // First, we must satisfy the base class_invariant().
}
auto int res( StateTo<Trans>::class_invariant() );
auto int i;
for( i = 0; (i < domain()) && res; i++ ) {
    res = lookup( i ).class_invariant() && (lookup( i ).range() == domain());
} return( res );
}
#endif

Implementation: TransRel inherits for implementation from StateTo<Trans>. Many of the member functions simply call the corresponding members functions of StateTo or Trans.
map( i ).set_range( r );
} assert( class_invariant() );
}

void TransRel::reincarnate() {
    assert( class_invariant() );
    auto int i;
    // First, reincarnate() all of the components before *this.
    for( i = 0; i < domain(); i++ ) {
        map( i ).reincarnate();
    }
    StateTo<Trans>::reincarnate();
    assert( class_invariant() );
}

TransRel& TransRel::disjointing_union( const TransRel& r ) {
    assert( class_invariant() && r.class_invariant() );
    auto int olddom( domain() );
    StateTo<Trans>::disjointing_union( r );
    assert( StateTo<Trans>::domain() == r.domain() + olddom );
    auto int i:
    for( i = 0; i < olddom; i++ ) {
        map( i ).set_range( domain() );
    }
    for( ; i < domain(); i++ ) {
        map( i ).at_rename( olddom );
    }
    assert( class_invariant() );
    return( *this );
}

ostream& operator<<( ostream& os, const TransRel& r ) {
    assert( r.class_invariant() );
    return( os << (const StateTo<Trans>&)r );
}

Implementation class: DTransRel

Files: dtransre.h, dtransre.cpp

Uses: CharRange, CRSet, DTrans, State, StateSet, StateTo

Description: A DTransRel maps a State and a character (or a CharRange) to a State. It is used to implement the deterministic transition relation in the class DFA (see [Wat93a, Property 2.25]). The domain of the relation must be maintained explicitly by the client. Transitions can be added to the relation, but not removed. The client must ensure that the relation remains deterministic.
// Implement a deterministic transition relation, as a function from States time
// char to State. This is used for transition relations in DFA's.

class DTransRel : protected StateTo<DTrans> {  
public:
        // Constructors, destructors, operator=:
        // Argument-less constructor:
        inline DTransRel();

        inline DTransRel( const DTransRel& r );

        // Default destructor is okay:

        inline const DTransRel& operator=( const DTransRel& r );

        // Some relational operators:

        // Compute the image of r, and a under *this.
        inline State image( const State r, const char a ) const;

        // Compute the image of r, and CharRange a under *this.
        inline State transition_on_range( const State r, const CharRange a ) const;

        // Compute the reverse transition on a range under *this:
        StateSet reverse_transition( const State r, const CharRange a ) const;

        // What are the transition labels between some State r and StateSet s?
        inline CharSet labels_between( const State r, const StateSet& s ) const;

        // Given State r, what are the labels on transitions out of r?
        inline CharSet out_labels( const State r ) const;

        // What are all States reverse reachable from r?
        StateSet reverse_closure( const StateSet& r ) const;

        // Some functions updating *this:

        inline DTransRel& add_transition( const State p, const CharRange a, const State q );

        // Some domain members:

        // What is the domain of this relation.
        inline int domain() const;

        // Change the domain of this relation.
        inline void set_domain( const int r );

        // Recycle this entire structure.
        void reincarnate();

        // Some special members:

        friend ostream& operator<< ( ostream& os, const DTransRel& r );

        inline int class_invariant() const;
    };

inline DTransRel::DTransRel() :
        StateTo<DTrans>() {  
        assert( class_invariant() );
    }
inline DTransRel::DTransRel( const DTransRel& r ) {
    StateTo<DTrans>( (StateTo<DTrans>&)r ) {
        assert( class_invariant() );
    }
}

inline const DTransRel& DTransRel::operator=( const DTransRel& r ) {
    assert( r.class_invariant() );
    // Call back to the base class.
    StateTo<DTrans>::operator=( (const StateTo<DTrans>&) r );
    assert( class_invariant() );
    return( *this );
}

inline State DTransRel::image( const State r, const char a ) const {
    assert( class_invariant() );
    assert( 0 <= r && r < domain() );
    return( lookup( r )[a] );
}

inline State DTransRel::transition_on_range( const State r, const CharRange a ) const {
    assert( class_invariant() );
    assert( 0 <= r && r < domain() );
    return( lookup( r ).range_transition( a ) );
}

inline CRSet DTransRel::labels_between( const State r, const StateSet& s ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    assert( domain() == s.domain() );
    assert( 0 <= r && r <= domain() );
    return( lookup( r ).labels_into( s ) );
}

inline CRSet DTransRel::out_labels( const State r ) const {
    assert( class_invariant() );
    return( lookup( r ).out_labels() );
}

inline DTransRel& DTransRel::add_transition( const State p, const CharRange a, const State q ) {
    assert( class_invariant() );
    map( p ).add_transition( a, q );
    assert( class_invariant() );
    return( *this );
}

inline int DTransRel::domain() const {
    return( StateTo<DTrans>::domain() );
}

inline void DTransRel::set_domain( const int r ) {
    assert( class_invariant() );
    StateTo<DTrans>::set_domain( r );
    assert( class_invariant() );
}

inline int DTransRel::class_invariant() const {
    // First, we must satisfy the base class_invariant().
    auto int i( StateTo<DTrans>::class_invariant() );
    auto int k;
    for( k = 0; (k < domain()) && i; k++ ) {
        i = i && lookup( k ).class_invariant();
    }
    return( i );
inline ostream& operator<<( ostream& os, const DTransRel& r ) {  
    assert( r.class_invariant() );  
    return( operator<<( os, (const StateTo<DTrans>&)r ) );  
}
#endif

Implementation: DTransRel inherits for implementation from StateTo<DTrans>. Many of the member functions are simply calls to the members of StateTo or DTrans.

StateSet DTransRel::reverse_transition( const State r, const CharRange a ) const {  
    assert( class_invariant() );  
    assert( (0 <= r) && (r < domain()) );  
    auto StateSet ret;  
    ret.set_domain( domain() );  
    auto State p;  
    // Iterate over all of the States, and see which have transitions to r.  
    for( p = 0; p < domain(); p++ ) {  
        if( transition_on_range( p, a ) == r ) {  
            ret.add( p );  
        }  
    }  
    return( ret );
}

StateSet DTransRel::reverse_closure( const StateSet& r ) const {  
    assert( class_invariant() );  
    assert( r.class_invariant() );  
    assert( domain() == r.domain() );  
    auto StateSet result( r );  
    auto StateSet intermediate;  
    intermediate.set_domain( domain() );  
    assert( intermediate.empty() );  
    do {  
        result.set_union( intermediate );  
        intermediate.clear();  
        // Go through the State's to see which can reach set result on  
        // some transition.  
        auto State st;  
        for( st = 0; st < domain(); st++ ) {  
            if( result.not_disjoint( lookup( st ).range( domain() ) ) ) {  
                intermediate.add( st );  
            }  
        }  
    } while( !result.contains( intermediate ) );  
    return( result );
}

void DTransRel::reincarnate() {  
    assert( class_invariant() );  
    // Don't reset the domain()  
    auto int i;
for( i = 0; i < domain(); i++ ) {
    map( i ).reincarnate();
}
StateTo<DTrans>::reincarnate();
assert( class_invariant() );
Part II
Regular expressions and $\Sigma$-algebras

7 Regular expressions

Implementation class: \texttt{REops}

Files: \texttt{reops.h}

Uses:

Description: The regular operators (of the $\Sigma$-algebra — see class \texttt{Reg} or [Wat93a, Definition 3.12]) are encoded as elements of an enumeration. These are usually stored in a type-field in regular expressions. They play an important role in computing the homomorphic ($\Sigma$-algebra) image of a regular expression.

Implementation: \texttt{REops} is defined as an \texttt{enum}.

```c
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $  // $Date: 1994/08/15 14:00:47 $
#ifndef REOPS_H
#define REOPS_H

enum REops {
  EPSILON,
  EMPTY,
  SYMBOL,
  OR,
  CONCAT,
  STAR,
  PLUS,
  QUESTION
};
#endif
```

Implementation class: \texttt{RE}

Files: \texttt{re.h, re.cpp}

Uses: \texttt{CharRange, CRSet, REops}

Description: An \texttt{RE} is a regular expression (see [Wat93a, Section 3] for the definition of regular expressions). The argumentless constructor constructs the regular expression $\emptyset$ (the regular expression denoting the empty language). It is not possible to construct more complex regular expressions with the member functions (see the template class \texttt{Reg<RE>} for information on constructing regular expressions). Class \texttt{RE} is the implementation of the carrier set of the $\Sigma$-term algebra (an \texttt{RE} is a term in the $\Sigma$-term algebra).

A member function returns the operator type of the main (or root) operator of the regular expression. (The operator types are enumerated in \texttt{REops}.) Other member functions can be used to determine additional information about the main (root) operator of the regular expression.
expression: the associated \textit{CharRange} (for a \textit{SYMBOL} basis operator), and the subexpressions (for non-basis, binary and unary, operators). A very basic \textit{istream} extraction (input) operator is provided, expecting the input to be a regular expression in prefix notation; the operator does little error checking.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
/* $Revision: 1.2 $ */
/* $Date: 1994/08/15 14:00:46 $ */
ifndef RE_H
#define RE_H
#include "charrang.h"
#include "crset.h"
#include "reaps.h"
#include <assert.h>
#include <iostream.h>

// Regular expressions.
// This class is not too useful for actually constructing them (for that, see
// the Sigma-algebra class in sigma.h). Class RE functions as the terms of the
// Sigma-term algebra. Some members for manipulating them are provided.

class RE {
public:
    // Some constructors, destructors, and operators:
    // By default, create the RE denoting the empty language.
    RE( );
    RE( const RE& r );
    virtual "RE();
    const RE& operator=( const RE& r );

    // Some basic access member functions:
    // How many \textit{SYMBOL} nodes in this regular expression?
    // Mainly for use in constructing RPA's.
    int num_symbols() const;
    // How many operators in this RE?
    // Used in ISImpl (the item set stuff).
    int num_operators() const;
    // What is the main operator of this regular expression?
    inline REops root_operator() const;
    // What is the \textit{CharRange} of this RE, if it is a \textit{SYMBOL} regular expression.
    inline CharRange symbol() const;
    // What is the left RE of this RE operator (if it is a higher operator).
    inline const RE& left_subexpr() const;
    // What is the right RE of this RE operator (if it is a binary operator).
    inline const RE& right_subexpr() const;

    // Some derivatives (Brzozowski's) related member functions:
    // Does this accept epsilon?
    // This is from Definition 3.20 and Property 3.21
    int Null() const;
    // This is from Definition 4.60
```
CRSet First() const;

// Is *this in similarity normal form (SNF)?
int in_snf() const;

// Put *this into similarity normal form.
RE& snf();

// Reduce (optimize) *this by removing useless information.
RE& reduce();

// What is the derivative of *this (w.r.t. r)?
RE derivative( const CharRange& r ) const;

// Some ordering members (largely used in similarity and comparisons)
int ordering( const RE& r ) const;

// Some comparisons:
inline int operator==( const RE& r ) const;
inline int operator!=( const RE& r ) const;
inline int operator<( const RE& r ) const;
inline int operator<=( const RE& r ) const;
inline int operator>( const RE& r ) const;
inline int operator>=( const RE& r ) const;

// Some extras:
friend ostream& operator<<( ostream& os, const RE& r );
friend istream& operator>>( istream& is, RE& r );

protected:

// Some protected helpers, mainly for Reg<RE>:
inline void set_root_operator( REops r );
inline void set_symbol( const CharRange r );

// Make a copy of *this in *r
void shallow_copy( RE* const r ) const;

void reincarnate();

// Put *this into OR normal form (ONF) (used in the snf functions)
RE& onf();

// Destroy node *r, moving its data into *this.
void assimilate_node( RE* const r );

// Some implementation details.
REops op;
RE* left, *right;
CharRange sym;
};

inline RE::RE() : op( EMPTY ),
    left( 0 ),
    right( 0 ) {
    assert( class_invariant() );
}

// Some inline operators:

inline REops RE::root_operator() const {
    return( op );
}
inline CharRange RE::symbol( ) const {
  assert( op == SYMBOL );
  return( sym );
}

// Assume that *this is a higher (non-basis) regular operator.
inline const RE& RE::left_subexpr( ) const {
  assert( ( op != EMPTY ) && ( op != EPSILON ) && ( op != SYMBOL ) );
  return( *left );
}

// Assume that *this is a binary regular operator.
inline const RE& RE::right_subexpr( ) const {
  assert( ( op == OR ) || ( op == CONCAT ) );
  return( *right );
}

inline void RE::set_root_operator( REops r ) {
  // No class assertion here, since *this is still being constructed.
  op = r;
}

inline void RE::set_symbol( const CharRange r ) {
  assert( op == SYMBOL );
  sym = r;
}

inline int RE::operator==( const RE& r ) const {
  return( ordering( r ) == 0 );
}

inline int RE::operator!=( const RE& r ) const {
  return( ordering( r ) != 0 );
}

inline int RE::operator<( const RE& r ) const {
  return( ordering( r ) < 0 );
}

inline int RE::operator>( const RE& r ) const {
  return( ordering( r ) > 0 );
}

inline int RE::operator<=( const RE& r ) const {
  return( ordering( r ) <= 0 );
}

inline int RE::operator>=( const RE& r ) const {
  return( ordering( r ) >= 0 );
}

inline int RE::class_invariant( ) const {
  switch( op ) {
    case EPSILON:
    case EMPTY:
      return( left == 0 && right == 0 );
    case SYMBOL:
      return( left == 0 && right == 0 );
    case OR:
    case CONCAT:
      return( left != 0 && right != 0
               && left_subexpr().class_invariant( )
               && right_subexpr().class_invariant( ) );
    case STAR:
    case PLUS:
Implementation: Regular expressions are implemented as expression trees. The comparison of regular expressions is defined recursively and uses a depth-first traversal.

Performance: Use-counting could avoid the deep-tree copies that are performed. An RE could be stored as a prefix-form string, as in [RW93]; this could give some problems with the heavy use of members such as left_subexpr and right_subexpr.
assert( e.class_invariant() );

op = e.op;
right = 0;

switch( e.root_operator() ) {
    case EPSILON:
        break;
    case EMPTY:
        left = 0;
        break;
    case SYMBOL:
        left = 0;
        // In no other case is it safe to copy the sym.
        sym = e.sym;
        break;
    case OR:
    case CONCAT:
        delete right;
        right = new RE( e.right_subexpr() );
    case STAR:
    case PLUS:
    case QUESTION:
        delete left;
        left = new RE( e.left_subexpr() );
        break;
}

assert( class_invariant() );
return( *this );

// How many symbols are there in *this RE?
int RE::num_symbols() const {
    assert( class_invariant() );
    auto int ret;
    switch( op ) {
        case EPSILON:
        case EMPTY:
        case SYMBOL:
            ret = 1;
            break;
        case OR:
        case CONCAT:
            ret = left->num_symbols() + right->num_symbols();
            break;
        case STAR:
        case PLUS:
        case QUESTION:
            ret = left->num_symbols();
            break;
    }
    return( ret );
}

// How many operators are there in *this RE?
int RE::num_operators() const {
    assert( class_invariant() );
    auto int ret;
    switch( op ) {
        case EPSILON:
        case EMPTY:
        case SYMBOL:
            ret = 1;
            break;
        case OR:
        case CONCAT:
        case STAR:
        case PLUS:
        case QUESTION:
            ret = 1;
            break;
    }
    return( ret );
}
break;
case OR:
case CONCAT:
    ret = left->num_operators() + right->num_operators() + 1;
    break;
case STAR:
case PLUS:
case QUESTION:
    ret = left->num_operators() + 1;
    break;
}
return( ret );

/// Copy *this into *r
void RE::shallow_copy( RE *const r ) const {
    r->op = op;
    r->left = left;
    r->right = right;
    r->sym = sym;
}

void RE::reincarnate() {
    if( left != 0 ) {
        // This is a non-basic operator for reincarnation.
        // Delete what used to be there.
        delete left;
        left = 0;
        if( right != 0 ) {
            // This is a binary operator for reincarnation.
            delete right;
            right = 0;
        }
    }
    op = EMPTY;
    assert( class_invariant() );
}

/// Output in prefix notation.
/// Part of this could be REops, if enums could be used for overload resolution.
ostream& operator<)( ostream& os, const RE& r ) {
    switch( r.root.operator() ) {
    case EMPTY:
        os << '0';
        break;
    case EPSILON:
        os << '1';
        break;
    case SYMBOL:
        os << r.symbol();
        break;
    case OR:
        os << " ! " << r.left_subexpr() << ' ' << r.right_subexpr();
        break;
    case CONCAT:
        os << " . " << r.left_subexpr() << ' ' << r.right_subexpr();
        break;
    case STAR:
        os << '*' << r.left_subexpr();
        break;
    case PLUS:
        os << '+' << r.left_subexpr();
        break;
    case QUESTION:
        os << '?' = << r.left_subexpr();
        break;
    }
return( os );
}

// Input, in prefix notation. Not much error checking is done!
istream& operator>>( istream& is, RE& r ) {
    // Wipe out anything that used to be in r.
    delete r.left;
    delete r.right;
    // If something goes wrong try to leave r as the EMPTY regular expression.
    r.left = r.right = 0;
    r.op = EMPTY;

    auto char c;
    // It could be that there is nothing in the input.
    if( !(is >> c) ) {
        is.clear( ios::badbit | is.rdstate() );
    } else {
        switch( c ) {
        case '1':
            r.op = EPSILON;
            break;
        case '0':
            r.op = EMPTY;
            break;
        case '[':
            // This stuff should really be as the extraction operator
            // in class CharRange.
            is >> c;
            if( c != ']' ) {
                cerr << "First CharRange incorrect\n";
                is.clear( ios::badbit | is.rdstate() );
            } else { // Input first character of the range.
                is >> c;
                auto char c2;
                is >> c2;
                if( c2 != ']' ) {
                    cerr << "First CharRange not terminated properly\n";
                    is.clear( ios::badbit | is.rdstate() );
                } else { // Input first char. of second range.
                    is >> c2;
                    if( c2 != ']' ) {
                        cerr << "Second CharRange incorrect\n";
                        is.clear( ios::badbit | is.rdstate() );
                    } else { // Input first char. of second range.
                        is >> c2;
                        r.sym = CharRange( c, c2 );
                        r.op = SYMBOL;
                        is >> c;
                        if( c != ']' ) {
                            cerr << "Second CharRange not terminated properly\n";
                            is.clear( ios::badbit | is.rdstate() );
                        }
                        is >> c;
                        if( c != ']' ) {
                            cerr << "Close block not found\n";
                            is.clear( ios::badbit | is.rdstate() );
                        }
                    }
                }
            }
        case '\':
            is >> c;
            r.sym = CharRange( c );
        case '
':
            is >> c;
        case '(':
7.2 Derivatives

Two functions on regular expressions, Null and First (see [Wat93a, Example 3.20, Definition 4.60]) — both used in computing derivatives of regular expressions), are defined as recursive member functions in file deriv.cpp. Function Null returns 1 if the empty word (e) is in the language denoted by the regular expression, and 0 otherwise. Function First returns a CRSet denoting the characters which can occur as the first character of a word in the language denoted by the regular expression.

Member function derivative takes a CharRange and returns the derivative of the regular expression with respect to the CharRange. The derivative is computed recursively according to the definition given in [Wat93a, Definition 5.23]. It is used in the implementation of Brzozowski's construction.
• All of its OR operators are left-associative; for example \((E_0 \cup E_1) \cup E_2\) is left-associative, while \(E_0 \cup (E_1 \cup E_2)\) is not.

• The OR regular expression is in one of two forms (where \(\leq\) is the total ordering on regular expressions): either \((E_0 \cup E_1) \cup E_2\) where \(E_1 \leq E_2\) and \((E_0 \cup E_1)\) is in ONF, or \(E_0 \cup E_1\) where \(E_0 \leq E_1\). (The use of \(<\) instead of \(\leq\) implies that if the subexpressions of an OR expression are equal (under the total ordering), then they are collapsed into one expression, implementing the idempotence of OR.)

The ONF is used to provide the equivalent of a multi-ary OR operator, implementing the associativity, commutativity and idempotence of the OR operator, as detailed in [Wat93a, Definition 5.26].

The member function \(\text{in_snf}\) returns 1 if the \(RE\) is in SNF, and 0 otherwise. The member \(\text{snf}\) puts the \(RE\) in SNF, using function \(\text{onf}\) as a helper. Member function \(\text{reduce}\) optimizes an \(RE\), removing any redundant subexpressions. It implements rules 1 to 5 of [Wat93a, Remark 5.32]. The function performs a bottom-up traversal of the regular expression, making such transformations as \(\emptyset \cdot E \rightarrow \emptyset\) and \(\emptyset^* \rightarrow \epsilon\).

All of these members (including the total ordering) are defined in file \(\text{deriv.cpp}\).

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 13:59:46$
#include "re.h"
#include "sigma.h"

// Does *this accept epsilon?
// Implement Definition 3.20 and Property 3.21
int RE::Null() const {
    auto int ret;
    switch( op ) {
    case EPSILON:
        ret = 1;
        break;
    case EMPTY:
    case SYMBOL:
        ret = 0;
        break;
    case OR:
        ret = left->Null() || right->Null();
        break;
    case CONCAT:
        ret = left->Null() && right->Null();
        break;
    case STAR:
        ret = 1;
        break;
    case PLUS:
        ret = left->Null();
        break;
    case QUESTION:
        ret = 1;
        break;
    }
    return( ret );
}

// What is the First of *this?
// Implement Definition 4.60
CRSet RE::First() const {
    assert( class_invariant() );
    auto int ret;
    switch( op ) {
    case EPSILON:
        ret = 0;
        break;
    case EMPTY:
    case SYMBOL:
        ret = 0;
        break;
    case OR:
        ret = left->First() || right->First();
        break;
    case CONCAT:
        ret = left->First() && right->First();
        break;
    case STAR:
        ret = 1;
        break;
    case PLUS:
        ret = left->First();
        break;
    case QUESTION:
        ret = 1;
        break;
    }
    return( ret );
}
```
7.1 Similarity of REs

```c
auto CRSet a;
switch( op ) {
  case EPSILON:
    break;
  case SYMBOL:
    a.append( sym );
    break;
  case OR:
    a = left->First().combine( right->First() );
    break;
  case CONCAT:
    a = left->First();
    if( left->Null() ) {
      a.combine( right->First() );
    }
    break;
  case STAR:
  case PLUS:
  case QUESTION:
    a = left->First();
    break;
}
return( a );
```

// Is *this in similarity normal form (SNF)?
int RE::in_snf() const {
  assert( class_invariant() );
  auto int ret;
  // SNF means that all OR nodes are left-associative, and
  // in the two cases:
  //   A OR B that A < B, or
  //   (A OR B) OR C that B < C and (A OR B) is in SNF.
  switch( op ) {
    case EPSILON:
    case EMPTY:
    case SYMBOL:
      ret = 1;
      break;
    case OR:
      if( (right->op == OR) || !right->in_snf() ) {
        // This can't be in SNF.
        ret = 0;
      } else {
        if( left->op == OR ) {
          // Second case.
          ret = (*left->right) < *right) & & left->in_snf();
        } else {
          // First case.
          ret = (*left < *right) & & left->in_snf();
        }
      }
      break;
    case CONCAT:
      // A CONCAT is in SNF if the subexpressions are in SNF.
      ret = left->in_snf() & & right->in_snf();
      break;
    case STAR:
    case PLUS:
    case QUESTION:
      // A unary subexpr. is in SNF if it's subexpressions are too.
      ret = left->in_snf();
      break;
  }
  return ret;
}
```
return( ret );
}

// Put *this into similarity normal form.
RE& RE::snf() {
    assert( class_invariant() );

    switch( op ) {
    case EPSILON:
        break;
    case EMPTY:
        break;
    case SYMBOL:
        // Already in SNF.
        break;
    case OR:
        // While the right is an OR, rotate them.
        // Note: this is not a true rotation, since we get
        // E \ ( F \ G ) \ ( F \ E \ G ) instead of ( E \ F ) \ G
        // but, by commutativity it's okay.
        while( right->op == OR ) {
            auto RE temp{ left };
            left = right;
            right = left->right;
            left->right = temp;
            assert( class_invariant() );
        }
        assert( right->op != OR );
        left->snf();
        assert( class_invariant() );
        assert( left->in_snf() );
        right->snf();
        assert( class_invariant() );
        assert( right->in_snf() );
        onf();
        assert( class_invariant() );
        break;
    case CONCAT:
        left->snf();
        assert( class_invariant() );
        assert( left->in_snf() );
        right->snf();
        assert( class_invariant() );
        assert( right->in_snf() );
        break;
    case STAR:
    case PLUS:
    case QUESTION:
        left->snf();
        assert( class_invariant() );
        assert( left->in_snf() );
        break;
    }
    assert( in_snf() );
    assert( class_invariant() );
    return( *this );
}

// What is the derivative of *this?
// Implement Definition 5.23
RE RE::derivative( const CharRange& r ) const {
    assert( class_invariant() );

    // Duplicate some of the knowledge appearing in sig-re.cpp
    auto RE e;
    switch( op ) {
    case EPSILON:
    case EMPTY:
        assert( e.op == EMPTY );
7.1 Similarity of REs

break;

case SYMBOL:
    if( r <= sym ) {
        e.op = EPSILON;
        assert( e.class_invariant() );
    }
    else {
        // else, it should be EMPTY.
        assert( e.op == EMPTY );
    }

break;

case OR:
    e.op = OR;
    e.left = new RE( left->derivative( r ) );
    e.right = new RE( right->derivative( r ) );
    assert( e.class_invariant() );

break;

case CONCAT:
    if( left->Null() ) {
        e.op = OR;
        e.right = new RE( right->derivative( r ) );
        e.left = new RE;
        e.left->op = CONCAT;
        e.left->left = new RE( left->derivative( r ) );
        e.left->right = new RE( *right );
        assert( e.class_invariant() );
    } else {
        e.op = CONCAT;
        e.left = new RE( left->derivative( r ) );
        e.right = new RE( *right );
        assert( e.class_invariant() );
    }

break;

case STAR:
    e.op = CONCAT;
    e.left = new RE( left->derivative( r ) );
    e.right = new RE( left->derivative( r ) );
    e.right = new RE( *this );
    assert( e.class_invariant() );

break;

case PLUS:
    e.op = CONCAT;
    e.left = new RE( left->derivative( r ) );
    e.right = new RE;
    e.right->op = STAR;
    e.right->right = 0;
    e.right->left = new RE( *left );
    assert( e.class_invariant() );

break;

case QUESTION:
    e = left->derivative( r );

break;

// By the class_invariant(), the default: cannot occur.

assert( e.class_invariant() );

return( (RE)e );

// Implement rules 1-4 of Remark 5.32 of the Taxonomy.
RE& RE::reduce() {
    assert( class_invariant() );

    switch( op ) {
        case EPSILON:
            case EMPTY:
            case SYMBOL:
                // These basis operators cannot be reduced further.

                e.op = EPSILON;
                e.class_invariant();

                return( (RE)e );
        case OR:
            e.op = OR;
            e.left = new RE( left->derivative( r ) );
            e.right = new RE( right->derivative( r ) );

            e.op = OR;
            e.left = new RE;
            e.left->op = CONCAT;
            e.left->left = new RE( left->derivative( r ) );
            e.left->right = new RE( *right );

            e.op = CONCAT;
            e.left = new RE( left->derivative( r ) );
            e.right = new RE( *right );

            e.op = EPSILON;
            e.class_invariant();

            return( (RE)e );
    }

    // Implement rules 1-4 of Remark 5.32 of the Taxonomy.
    RE& RE::reduce() {
        assert( class_invariant() );

        switch( op ) {
            case EPSILON:
                case EMPTY:
                case SYMBOL:
                    // These basis operators cannot be reduced further.
break;

case OR:
    left->reduce();
    assert( class_invariant() );
    right->reduce();
    assert( class_invariant() );
    if( left->op == EMPTY ) {
        // EMPTY OR E == E
        delete left;
        assimilate_node( right );
        assert( class_invariant() );
    } else if( right->op == EMPTY ) {
        delete right;
        assimilate_node( left );
        assert( class_invariant() );
    }
    // The case where left and right are both EMPTY is handled correctly.
    break;

case CONCAT:
    left->reduce();
    assert( class_invariant() );
    right->reduce();
    assert( class_invariant() );
    if( left->op == EMPTY ) {
        // EMPTY CONCAT E == EMPTY
        delete right;
        assimilate_node( left );
        assert( class_invariant() );
    } else if( right->op == EMPTY ) {
        delete left;
        assimilate_node( right );
        assert( class_invariant() );
    } else if( left->op == EPSILON ) {
        // EPSILON CONCAT E == E
        delete left;
        assimilate_node( right );
        assert( class_invariant() );
    } else if( right->op == EPSILON ) {
        delete right;
        assimilate_node( left );
        assert( class_invariant() );
    }
    break;

case STAR:
    left->reduce();
    assert( class_invariant() );
    if( left->op == EPSILON ) {
        // EPSILON * == EPSILON
        assimilate_node( left );
        assert( class_invariant() );
    } else if( left->op == EMPTY ) {
        // EMPTY * == EMPTY
        delete left;
        op = EPSILON;
        left = right = 0;
        assert( class_invariant() );
    }
    break;

case PLUS:
    left->reduce();
    assert( class_invariant() );
    if( left->op == EPSILON ) {
        // EPSILON + == EPSILON
        assimilate_node( left );
        assert( class_invariant() );
    } else if( left->op == EMPTY ) {
        // EMPTY + == EMPTY
7.1 Similarity of REs

```c
assimilate_node( left );
assert( class_invariant() );
}
break;
case QUESTION:
    left->reduce();
    assert( class_invariant() );
if( left->op == EPSILON ) {
    // EPSILON == EPSILON
    assimilate_node( left );
    assert( class_invariant() );
} else if( left->op == EMPTY ) {
    // EMPTY == EPSILON
    delete left;
    op = EPSILON;
    left = right = 0;
    assert( class_invariant() );
}
break;
}
assert( class_invariant() );
return( *this );

// Put *this into OR normal form (ONF) (used in the snf functions)
RE& RE::onf() {
    assert( class_invariant() );
    assert( op == OR );
    assert( right->op != OR );

    // Two possibilities: (A OR B) OR C, A OR B.
    if( left->op == OR ) {
        // First one.
        // Figure out which order B,C should be in:
        auto int ord( left->right->ordering( *right ) );
        if( ord == 0 ) {
            // B == C
            // destroy C.
            delete right;
            assimilate_node( left );
        } else if( ord > 0 ) {
            // B > C
            // Need to swap B,C
            // Just do it with pointers:
            auto RE *const temp( right );
            right = left->right;
            left->right = temp;

            // Also make sure that (what's now) (A OR C)
            // is also in ONF.
            left->onf();
        } else { // else, it's already in OR NF.
            // Second case.
            // Figure out the order of A,B:
            auto int ord( left->ordering( *right ) );
            if( ord == 0 ) {
                // A == B
                // Wipe out one of the two nodes, and make *this
                // equal to the remaining one.
                delete right;
                assimilate_node( left );
            } else if( ord > 0 ) {
                // A > B
                // Need to swap A,B (left and right):
                auto RE *const temp( right );
```
```cpp
right = left;
left = temp;
}
// else, it's already in OR NF
}
assert( class_invariant() );
assert( in_snf() );
return( *this );
}

// Some ordering members (largely used in similarity)
// (This is a kind of lexicographic ordering.)
int RE::ordering( const RE& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    if( op !~ r.op ) {
        return( op - r.op );
    } else {
        assert( op == r.op );
        auto int ret;
        switch( op ) {
        case EPSILON:
            ret = 0;
            break;
        case SYMBOL:
            ret = sym.ordering( r.sym );
            break;
        case OR:
        case CONCAT:
            ret = left->ordering( r.left_subexpr() );
            if( ret != 0 ) {
                // Do nothing.
            } else { // Compare the right subexprs, since the left
                // ones are equal.
                ret = right->ordering( r.right_subexpr() );
            }
            break;
        case STAR:
        case PLUS:
        case QUESTION:
            ret = left->ordering( r.left_subexpr() );
            break;
        }
        return( ret );
    }
}

// Assimilate *r into *this, and wipe out *r.
void RE::assimilate_node( RE *const r ) {
    // We can't assert( class_invariant() );
    // since we can't be sure that *this is still good.
    // This just delays catching any potential invariant violations.
    assert( r->class_invariant() );
    op = r->op;
    left = r->left;
    right = r->right;
    sym = r->sym;
    assert( class_invariant() );
    r->op = EMPTY;
    r->left = r->right = 0;
    delete r;
}
```
7.2 Derivatives

Two functions on regular expressions, Null and First (see [Wat93a, Example 3.20, Definition 4.60]) — both used in computing derivatives of regular expressions), are defined as recursive member functions in file deriv.cpp. Function Null returns 1 if the empty word (ε) is in the language denoted by the regular expression, and 0 otherwise. Function First returns a CRSet denoting the characters which can occur as the first character of a word in the language denoted by the regular expression.

Member function derivative takes a CharRange and returns the derivative of the regular expression with respect to the CharRange. The derivative is computed recursively according to the definition given in [Wat93a, Definition 5.23]. It is used in the implementation of Brzozowski's construction.
8 The $\Sigma$-algebra

The $\Sigma$-algebra (as defined in [Wat93a, Section 3]) is defined as a template class.

Implementation class: Reg

Files: sigma.h

Uses: CharRange, RE, REops

Description: Reg implements the $\Sigma$-algebra as defined in [Wat93a, Section 3]. The single type (or generic) parameter $T$ is used to specify the carrier set of the $\Sigma$-algebra. The template Reg is an abstract template, in the sense that it is only used to force a common interface between the $\Sigma$-algebras. Most of the member functions are left undefined, with versions specific to the carrier set being defined separately (see files sig-re.cpp, sig-fa.cpp, and sig-rfa.cpp). Regular expressions (class RE) are taken as the $\Sigma$-term algebra. A special member function is defined, taking a regular expression, and constructing the homomorphic image of the RE, in the $\Sigma$-algebra of $T$. The $\Sigma$-algebra is the only way to construct regular expressions, and it is particularly useful for constructing parts of a finite automaton separately. Reg<$T>$ inherits publicly from a $T$, and since a Reg<$T>$ contains nothing extra, it can be cast to a $T$ safely (no slicing occurs, and no information is lost). Presently, only Reg<RE>, Reg<FA>, and Reg<RFA> are available. (Each of RE, FA, and RFA also have constructors taking a const reference to RE; for practical purposes, these constructors are much faster than constructing the homomorphic image using Reg.)

Implementation: The template is used to force a common interface for $\Sigma$-algebras with different carrier sets. Forcing an interface is commonly done using abstract base classes, although this is not possible in this case, as the return types of many of the member functions would not be correct.

Performance: The homomorphic image member function is recursive, and keeps partially constructed results in local variables. Use-counting in the carrier class can improve performance.

```c
/* (c) Copyright 1994 by Bruce Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:50 $
#undef SIGMA_H
#define SIGMA_H

#include "charrang.h"
#include "reops.h"
#include "rre.h"

// The signature of the Sigma-algebra is defined as template Reg. It is only
// an interface template, with most of the member functions being filled in as
// specialized versions. Given class T, the class Reg<$T>$ is a Sigma-algebra with
// $T$ as the carrier set. Reg is the only template required, since the regular
// expression Sigma algebra only has one sort: Reg.
// The Reg of a class T is also publicly derived from T, meaning that once
// a Reg<$T>$ is constructed (perhaps using the Sigma-algebra operators), it can
// be used as a regular T (without slicing).
// A special constructor is provided, which constructs the homomorphic image in
// algebra Reg<$T>$ of some regular expression. (Regular expressions are the Sigma-
// term algebra.)

template<class T>
class Reg : public T {
public:
    // Some constructors. Usually pass control back to the base class.
    Reg() : T() {}
    Reg( const Reg<$T>$& r ) :
```
// Just pass back to the base class.
T( const T &r ) {}

// Construct the homomorphic image of regular expression r, using the Sigma-algebra operators (corresponding to the operators of the regular expression).
Reg( const RE & r ) {
    assert( r.class_invariant() );
    homomorphic_image( r );
    assert( T::class_invariant() );
}

// Default destructor falls through to the "T()"
inline const Reg<T>& operator=( const Reg<T>& r ) {
    // Call through to the base-class's operator=
    T::operator=( r );
    return( *this );
}

// Now the Sigma-algebra signature:

// Sigma-algebra basis
Reg<T>& epsilon();
Reg<T>& empty();
Reg<T>& symbol( const CharRange r );

// Sigma-algebra binary ops
Reg<T>& or( const Reg<T>& r );
Reg<T>& concat( const Reg<T>& r );

// Sigma-algebra unary ops
Reg<T>& star();
Reg<T>& plus();
Reg<T>& question();

protected:

// Helper for constructing the homomorphic image of a regular expression.
// (Can also be used as a type of copy constructor for RE's.)
inline void homomorphic_image( const RE & r );

// Helper for constructing the homomorphic image of a regular expression.
// (Can also be used as a type of copy constructor for RE's.)
template<class T>
void Reg<T>::homomorphic_image( const RE & r ) {
    assert( r.class_invariant() );
    switch( r.root_operator() ) {
    case EMPTY:
        // Default constructor T() should already leave *this accepting the
        // empty language.
        // The following call is redundant:
        empty();
        break;
    case EPSILON:
        // Make *this accept epsilon.
        epsilon();
        break;
    case SYMBOL:
        // Make *this accept the specified symbol.
        symbol( r.symbol() );
        break;
    case OR:
8.1 The $\Sigma$-term algebra $Reg<RE>$

Implementation class: $Reg<RE>$

Files: $sig-re.cpp$

Uses: $CharRange$, $Reg$, $RE$, $REops$

Description: The template instantiation $Reg<RE>$ can be used to construct regular expressions (see the documentation on template class $Reg$).

Implementation: The implementation consists of straight-forward expression tree manipulations. The implementation of $RE$ is declared protected to give $Reg<RE>$ access to it.

Performance: These member functions would benefit from use-counting of class $RE$.

```plaintext
Reg<RE>& Reg<RE>::epsilon() {
    // This may have been something in a previous life.
    reincarnate();
}
```
8.1 The \( \Sigma \)-term algebra Reg<RE>

```
set_root_operator( EPSILON );
assert( class_invariant() );
return( *this );
}
Reg<RE>& Reg<RE>::empty() {
    // See epsilon() case.
    reincarnate();
    set_root_operator( EMPTY );
    assert( class_invariant() );
    return( *this );
}
Reg<RE>& Reg<RE>::symbol( const CharRange r ) {
    // See epsilon case.
    reincarnate();
    set_root_operator( SYMBOL );
    set_symbol( r );
    assert( class_invariant() );
    return( *this );
}
Reg<RE>& Reg<RE>::or( const Reg<RE>& r ) {
    auto RE *const lft( new RE );
    // Make a copy of *this.
    shallow_copy( lft );
    auto RE *const rt( new RE( r ) );
    op = OR;
    left = lft;
    right = rt;
    assert( class_invariant() );
    return( *this );
}
Reg<RE>& Reg<RE>::concat( const Reg<RE>& r ) {
    auto RE *const lft( new RE );
    shallow_copy( lft );
    auto RE *const rt( new RE( r ) );
    op = CONCAT;
    left = lft;
    right = rt;
    assert( class_invariant() );
    return( *this );
}
Reg<RE>& Reg<RE>::star0 {
    auto RE *const d( new RE );
    shallow_copy( d );
    op = STAR;
    left = d;
    right = 0;
    assert( class_invariant() );
    return( *this );
```
Reg<RE>& Reg<RE>::plus() {
    assert( class_invariant() );

    auto RE *const d( new RE );
    shallow_copy( d );
    op = PLUS;
    left = d;
    right = 0;

    assert( class_invariant() );
    return( *this );
}

Reg<RE>& Reg<RE>::question() {
    assert( class_invariant() );

    auto RE *const d( new RE );
    shallow_copy( d );
    op = QUESTION;
    left = d;
    right = 0;

    assert( class_invariant() );
    return( *this );
}
Part III
Automata and their components

9 Abstract finite automata

Implementation class: \textit{FAabs}

Files: \textit{faabs.h}

Use: \textit{DFA} (mentioned)

\textbf{Description:} \textit{FAabs} is an abstract class, used to provide a common interface to the different types of finite automata (which are \textit{FA}, \textit{DFA}, \textit{RFA}, \textit{LBFA}, and \textit{RBFA}). The following operations are provided via member functions:

1. determine how many states there are in the concrete finite automaton;
2. restart the finite automaton in its start states;
3. advance to a new set of states (make a transition) on a character;
4. determine if the automaton is in accepting states;
5. determine if the automaton is stuck (unable to make further transitions);
6. determine if a string is in the language of the automaton, and
7. return a deterministic finite automaton (a \textit{DFA}) which accepts the same language.

By convention, the argumentless constructor of a class inheriting from \textit{FAabs} must construct an object of that class that accepts the empty language (i.e. the object does not accept anything).

\textbf{Implementation:} Since it is an abstract base class (pure virtual), all except one of the member functions are pure virtual; the \textit{acceptable} member function is defined.

```c++
// Copyright 1994 by Bruce Watson +/
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:42 $
#ifndef FAABS_H
#define FAABS_H

class DFA;

class FAabs {
public:
    // Return the number of states (or some reasonably close measure).
    virtual int num_states() const = 0;
    // Reset the current state before beginning to process a string.
    // This is not the default condition for most of the derived classes.
    virtual void restart() = 0;
    // Advance the current state by processing a character.
    virtual void advance( char a ) = 0;
    // Is the current state an accepting (final) one?
    virtual int in_final() const = 0;
    // Is the automaton stuck?
};
#endif
```

virtual int stuck() = 0;

// Is the string w acceptable?
virtual int acceptable( const char *w ) {
    for( restart(); stuck() && *w; advance( *(w++) ) ) {}  
    // It's acceptable if *this is final, and the whole of w was consumed.
    return( in_final() && !*w );
}

// Return a DFA accepting the same language as *this.
virtual DFA determinism() const = 0;

#endif
10 Deterministic FAs

Implementation class: DFA

Files: dfa.h, dfa.cpp

Uses: DFA_components, DSDFAR, DTransRel, FAabs, State, StateSet, StateEqRel, SymRel

Description: A DFA is a deterministic finite automaton (see [Wat93a, Property 2.25] for the definition of DFA). It implements the interface required by the abstract class FAabs. Additionally, it also provides member functions to reverse the automaton, and to minimize it (the minimization functions are described later). A special constructor is provided, which takes a structure containing the essential components of a DFA, and uses those components to construct the automaton; this constructor is used in several DFA constructions involving the subset construction (see [Wat93a, p. 12–13] for the subset construction).
// Some minimization algorithms:
inline DFA& min_Brzoowski();
DFA& min_HopcroftUllman();
DFA& min_dragon();
DFA& min_Hopcroft();
DFA& min_Watson();

// Some member functions relating to useful State's.
// Can all States reach a final State?
int Useful() const;

// Remove any States that cannot reach a final State.
DFA& useless();

// Special member functions:
friend ostream& operator<<( ostream& os, const DFA& r );
inline int class_invariant() const;

protected:
// Given a minimizing equivalence relation, shrink the DFA.
DFA& compress( const StateEqRel& r );
DFA& compress( const SymRel& r );

// Attempt to split the eq. class [p]_P w.r.t. [q]_P
// (Return 1 if it was split, 0 otherwise.)
State split( const State p, const State q, const CharRange a, StateEqRel& P ) const;

// A helper for min_Watson.
int arc_eq( State p, State q, SymRel& S, const StateEqRel& H, const SymRel& Z ) const;

// Implementation details:
StatePool Q;
// S must be a singleton set, or empty.
StateSet S;
StateSet F;
DTransRel T;

// Some simulation details:
State current;

};

// Create a DFA with no State's, accepting nothing.
inline DFA::DFA() :
Q(),
S(),
F(),
T(),
current( Invalid ) {
assert( class_invariant() );
}

// A special constructor used for subset construction etc.
// A DFA constructed this way will always have a start State.
inline DFA::DFA( const DFA_components& r ) :
Q( r.Q ),
S( r.S ),
T( r.T ),
F( r.F ) {
current = Invalid;
}
assert( class_invariant() );
}

inline DFA& DFA::reconstruct( const DFA_components& r ) {
    Q = r.Q;
    S = r.S;
    T = r.T;
    F = r.F;
    current = Invalid;
    assert( class_invariant() );
    return( *this );
}

// See Section 2 (p. 5) of the minimization Taxonomy.
inline DFA& DFA::min_Brzozowski() {
    assert( class_invariant() );
    return( reverse().reverse() );
}

inline int DFA::class_invariant() const {
    return( Q.size() == S.domain()
        && Q.size() == F.domain()
        && Q.size() == T.domain()
        && current < Q.size()
        && S.size() <= 1 );
}

#endif

Implementation: All of the member functions are straight-forward implementations, with the exception of the minimization and reversal member functions. Those are further described later (the minimization member functions are discussed in Part IV). The constructor which takes a DFA_components object is only provided because most C++ compilers do not support template member functions yet. When these are fully supported, template function construct_components will be inlined in the constructor.

Performance: Use-counting the classes used in the automata (such as DTransRel and StateSet) would improve performance on copy construction and assignment.

int DFA::num_states() const {
    assert( class_invariant() );
    return( Q.size() );
}

void DFA::restart() {
    assert( class_invariant() );
    // There really should be only one state in S.
    // this is ensured by the class invariant.
    current = S.smallest();
}

void DFA::advance( char a ) {
assert( class_invariant() );
assert( a != Invalid );
current = T.image( current, a );
}

int DFA::in_final() const {
    assert( class_invariant() );
    return( F.contains( current ) );
}

int DFA::stuck() {
    assert( class_invariant() );
    return( current == Invalid );
}

DFA DFA::determinism() const {
    assert( class_invariant() );
    // Keep it really simple!
    return( *this );
}

DFA& DFA::reverse() {
    // Make sure that *this is structurally sound.
    assert( class_invariant() );
    // Now construct the DFA components from an abstract DFA state.
    reconstruct( construct_components( DSDFARev( F, &T, &S ) ) );
    return( *this );
}

DFA& DFA::compress( const StateEqRel& r ) {
    assert( class_invariant() );
    // All of the components will be constructed into this structure:
    auto DFA_components ret;
    // Give each eq. class of r a State name:
    // newnames maps States to their new names after compression.
    auto StateTo<State> newnames;
    newnames.set_domain( Q.size() );
    auto State st;
    for( st = 0; st < Q.size(); st++ ) {
        // If st is the representative of its class allocate a new name.
        auto State strep( r.eq_class_representative( st ) );
        if( st == strep ) {
            newnames.map( st ) = ret.Q.allocate();
        } else {
            newnames.map( st ) = newnames.lookup( strep );
        }
    }
    // Construct the new transition relation.
    ret.T.set_domain( ret.Q.size() );
    ret.F.set_domain( ret.Q.size() );
    auto CRSet a;
    for( st = 0; st < Q.size(); st++ ) {
        // If st is the representative, construct the transition.
        if( st == r.eq_class_representative( st ) ) {
            auto State stprime( newnames.lookup( st ) );
            // What are st's out-transitions?
            auto CharRange b;
            a = T.out_labels( st );
            // The out-labels of any other element of [st]_r could have
// been used instead. Some other choice may, indeed, lead
// to a smaller DFA. This approach is used for simplicity.
auto int it;
// Iterate over the labels, constructing the transitions.
for( it = 0; it != a.end(); it++ ) {
    b = a.iterator( it );
    ret.T.add_transition( stprime, b,
                         newnames.lookup( T.transition_on_range( st, b ) ) );
}

// st's eq. class may be final.
if( F.contains( st ) ) {
    ret.F.add( stprime );
}

// Set up the new start state.
ret.S.set_domain( ret.Q.size() );
ret.S.add( newnames.lookup( S.smallest() ) );
reconstruct( ret );
assert( class_invariant() );
return( *this );

DFA& DFA::compress( const SymRel& r ) {
    assert( class_invariant() );
    // All of the components will be constructed into this structure:
    auto DFA_components ret;
    // Give each eq. class of r a State name:
    // newnames maps States to their new names after compression.
    auto StateTo<State> newnames;
    newnames.set_domain( Q.size() );
    auto StateSet consider;
    consider.set_domain( Q.size() );
    consider.complement();
    // Build the set of representatives.
    auto StateSet repr;
    repr.set_domain( Q.size() );
    auto State st;
    for( consider.iter_start( st ); consider.iter_end( st ); consider.iter_next( st ) ) {
        // st will always be the representative of its class.
        assert( st == r.image( st ).smallest() );
        repr.add( st );
        // give st the new name
        auto State n( ret.Q.allocate() );
        newnames.map( st ) = n;
        // Go over [st], and give them all the new name.
        auto State z;
        for( r.image( st ).iter_start( z ); r.image( st ).iter_end( z );
             r.image( st ).iter_next( z ) ) {
            newnames.map( z ) = n;
        }
        // Now mark [st] as having been done already.
        consider.remove( r.image( st ) );
        // The outer iterator should still work okay.
    }
}
DETERMINISTIC FAS

// Construct the new transition relation.
ret.T.set_domain( ret.Q.size() );
ret.F.set_domain( ret.Q.size() );

auto CRSet a;
// Go over all of the representatives (eq. classes), constructing the
// transitions.
for( repr.iter.start( st ); !repr.iter.end( st ); repr.iter.next( st ) ) {
    auto State stprime( newnames.lookup( st ) );

    // What are st's out-transitions?
    auto CharRange b;
a = T.out_labels( st );
    // The out-labels of any other element of [st] could have
    // been used instead. Some other choice may, indeed, lead
    // to a smaller DFA. This approach is used for simplicity.
    auto int it;
    // Iterate over the labels, constructing the transitions.
    for( it = 0; !a.iter.end( it ); it++ ) {
        b = a.iterator( it );
        ret.T.add_transition( stprime, b,
                           newnames.lookup( T.transition_on_range( st, b ) ) );
    }

    // st's eq. class may be final.
    if( F.contains( st ) ) {
        ret.F.add( stprime );
    }
}

// Set up the new start state.
ret.S.set_domain( ret.Q.size() );
ret.S.add( newnames.lookup( S.smallest() ) );
reconstruct( ret );
assert( class_invariant() );
return( *this );

// Can all of the State's in *this reach a final State?
// Implement Definition 2.23
int DFA::Useful() const {
    assert( class_invariant() );
    auto StateSet r( T.reverse_closure( F ) );
    return( r.complement().empty() );
}

// Remove any State's from *this that cannot reach a final State. (This is a last
// step in minimization, since some of the min. algorithms may yield a DFA with a
// sink state.)
// Implement Remark 2.39
DFA& DFA::useful() {
    assert( class_invariant() );
    auto StateSet freachable( T.reverse_closure( F ) );

    auto StateTo<State> newnames;
    newnames.set_domain( Q.size() );

    // All components will be constructed into a special structure:
    auto DFA_components ret;

    auto State st;
    for( st = 0; st < Q.size(); st++ ) {
        // If this is a Useful State, carry it over by giving it a name
        // in the new DFA.
        if( freachable.contains( st ) ) {
            ret.T.add( st );
            ret.F.add( st );
            ret.S.add( st );
        }
    }

    return( *this );
}
newnames.map( st ) = ret.Q.allocate();
}

// It is possible that nothing needs to be done (ie. the all States were
// already P useful).
if( Q.size() != ret.Q.size() ) {
    ret.T.set_domain( ret.Q.size() );
    ret.F.set_domain( ret.Q.size() );

    auto CRSet a;
    for( st = 0; st < Q.size(); st++ ) {
        // Only construct the transitions if st is final reachable.
        if( reachable.contains( st ) ) {
            a = T.out_labels( st );
            auto State stprime( newnames.lookup( st ) );

            auto CharRange b;
            auto int it;
            // Construct the transitions.
            for( it = 0; it < a.iter_end(); it++ ) {
                b = a.iterator( it );
                ret.T.add_transition( stprime, b,
                                      newnames.lookup( T.transition_on_range( st, b ) ) );
            }

            // This may be a final State.
            if( F.contains( st ) ) {
                ret.F.add( stprime );
            }
        }
    }
    ret.S.set_domain( ret.Q.size() );

    // Add a start State only if the original one was final reachable.
    if( S.not_disjoint( reachable ) ) {
        ret.S.add( newnames.lookup( S.smallest() ) );
    }
    reconstruct( ret );
}

assert( class_invariant() );
return( *this );
}

ostream& operator<<( ostream& os, const DFA& r ) {
    os << "\nDFA\n";
    os << "Q = " << r.Q << "\n";
    os << "S = " << r.S << "\n";
    os << "F = " << r.F << "\n";
    os << "Transitions = \n" << r.T << "\n";
    os << "current = " << r.current << "\n";
    return( os );
}
11 Constructing a DFA

Implementation class: `DFA_components`

Files: dfacomp.h

Uses: `DTransRel, StatePool, StateSet`

Description: A `DFA_components` is a structure containing the essential parts of a DFA. Class `DFA` has a constructor taking a `DFA_components`. It is used in the subset construction, as the return type of template function `construct_components`.

Implementation: The class is implemented as a struct.

Performance: The template function `construct_components` should really be a constructor template of `DFA`. This would avoid needing `DFA_components`, and the overhead of passing the structure by value.

```c
/* (c) Copyright 1994 by Bruce W. Watson */
//@ Revision: 1.2 $ // $Dak 1994/08/15 14:00:28 $
ifndef DFACOMP_H
define DFACOMP_H

#include "st-pool.h"
#include "stateset.h"
#include "dtransrel.h"

struct DFA_components {
    StatePool Q;
    StateSet S;
    DTransRel T;
    StateSet F;
};

#endif
```

11.1 Abstract states and the subset construction

An abstract state is a certain type of object used in the construction of a DFA. They are essentially states in the subset construction (see [Wat03a, p. 12–13]) with some extra information. The template function `construct_components` uses the extra information to construct a DFA. All abstract states in FIRE engine are implemented using classes with names beginning with `DS`, such as `DSDFARev`. An abstract state must have the following member functions:

- an argumentless constructor, a copy constructor, an assignment operator, and equality and inequality operators;
- `final` — returns 1 if the abstract state is an accepting (final) state, 0 otherwise;
- `out_labels` — returns a `CRSet` representing the set of labels of out-transitions from the abstract state; and
- `out_transition` — takes a `CharRange` (which is an element of the `CRSet` returned by `out_labels`), and returns the abstract state resulting from a transition on the `CharRange`.

Implementation function: `construct_components`

Files: dfaseed.h
11.1 Abstract states and the subset construction

Uses: CharRange, CRSet, DFA_components, State, StateTo

Description: Template function `construct_components` implements the subset construction, with useless state removal [Wat93a, p. 12-13]. It takes an "abstract state" and constructs the components of a DFA (using the abstract state argument as the start state), returning them in a DFA_components structure. It does this by encoding the abstract states as States, using a breadth-first traversal of the start-reachable abstract states. During the traversal, the deterministic transition relation (a DTransRel) is also constructed. As a result of using a traversal starting at the start state, only reachable States and the reachable part of the deterministic transition relation are constructed.

Implementation: This template function should really appear as a template constructor of class DFA. It is implemented separately, as template member functions have only recently been incorporated into the draft of the C++ language standard. The function implements the subset construction (see [Wat93a, pp. 12-13]). Each abstract state is assigned a State, using a StateTo to keep track of the names assigned. The out-transitions from each State are constructed, using the StateTo to lookup the encoding of an abstract state. The abstract state classes are to be written in such a way that there are only finitely many of them, thereby ensuring termination of the function. The useless state removal is hidden in the fact that, by starting with the start state, only reachable States are constructed.

Performance: The performance could be greatly increased by making this template function a template constructor of class DFA, avoiding by-value structure passing.

```c++
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $ // $Date: 1994/08/15 14:00:29 $
#ifndef DFASEED_H
#define DFASEED_H
#include "state.h"
#include "stateto.h"
#include "charrang.h"
#include "crset.h"
#include "dfacomp.h"
#include <cassert>

// A template function, representing an abstract DFA, used in the construction
// of real DFA's.
// It is assumed that class T is a DFFE class (constrained genericity). See
// dsfrf.h or dsfarsev.h for examples of a DFFE class.

// Takes a T (an abstract start State), and constructs a DFA.
template<class T>
DFA_components construct_components( const T& abs_start ) {
  auto DFA_components ret;
  // Map each of the State's to a unique T, and vice-versa.
  auto StateTo<T> names;
  // Allocate a name for the new start State, and insert it into the namer.
  auto State s( ret.Q.allocate() );
  names.set_domain( ret.Q.size() );
  names.map( s ) = abs_start;
  ret.T.set_domain( ret.Q.size() );
  ret.F.set_domain( ret.Q.size() );
  // As invariant:
  // ret.T.domain() == ret.Q.size() & & ret.F.domain() == ret.Q.size()
  // and all States < current are already done.
```
auto State current;
  // Now construct the transitions, and finalness of each of the States.
  if( ret.Q grows as we go.
  for( current = 0; current < ret.Q.size(); current++ ) {
    // If the abstract state associated with current says it's final, then
    // current is too.
    if( names.lookup( current ).final() ) {
      ret.F.add( current );
    }
  }

  // Now go through all of the out-transitions of the abstract
  // state associated with current (ie. names.lookup( current )).
  auto CRSet a( names.lookup( current ).out_labels() );
  auto CharRange b;
  auto int it;
  for( it = 0; it != a.iter.end( it ); it++) {
    b = a.iterator( it );

    // Do something with the destination of the transition.
    auto T dest( names.lookup( current ).out_transition( b ) );

    // See if dest already has a name (by linear search).
    auto State i;
    for( i = 0; i < ret.Q.size() && (names.lookup( i ) != dest); i++) {
      // The abstract state may not have a name yet, so we may need
      // to allocate one.
      if( i == ret.Q.size() ) {
        // Associate i with dest.
        // And maintain the invariant.
        auto State j( ret.Q.allocate() );

        assert( i == j );
        names.set_domain( ret.Q.size() );
        ret.T.set_domain( ret.Q.size() );
        ret.F.set_domain( ret.Q.size() );

        names.map( i ) = dest;
      }

      // Now i is the name of dest.
      assert( dest == names.lookup( i ) );
      // Create the actual transition.
      ret.T.add_transition( current, b, i );
    } // That's all for current.
  }

  // Time to resynchronize the domain() of ret.S with ret.Q.size():
  ret.S.set_domain( ret.Q.size() );
  // Add the start State now.
  ret.S.add( s );

  // Return, constructing the DFA on the way out.
  return( ret );

$endif$

As an example of an abstract state, consider class \texttt{DSDFARev}.

\textbf{Implementation class:} \texttt{DSDFARev}
11.1 Abstract states and the subset construction

Files: dsdfarev.h, dsdfarev.cpp

Uses: CharRange, CRSet, DTransRel, StateSet

Description: Class DSDFARev is used in the reversal of a DFA. It is an abstract state representing a set of states, in the reversal of the deterministic automaton (the reversal of a deterministic finite automaton is not necessarily deterministic). A DSDFARev is constructed in the DFA member function reverse. Member functions support all of the required interface of an abstract state.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:30 $
#ifndef DSDFAREV_H
#define DSDFAREV_H

#include "charrang.h"
#include "crset.h"
#include "stateset.h"
#include "dtransre.h"
#include <assert.h>
#include <iostream.h>

This class is used to represent abstract States in a DFA that is still under construction. It is used in the reversal of a DFA.
// Objects of the class represent States in the subset construction of the reverse of a DFA.

class DSDFARev {
public:
   // Must always have an argument-less constructor.
   inline DSDFARev();

   inline DSDFARev( const DSDFARev& r );

   // A special constructor:
   inline DSDFARev( const StateSet& rq,
                     const DTransRel *T,
                     const StateSet *S );

   inline const DSDFARev& operator=( const DSDFARev& r );

   // The required member functions:
   inline int final() const;
   CRSet out_labels() const;
   DSDFARev out_transition( const CharRange a ) const;
   inline int operator==( const DSDFARev& r ) const;
   inline int operator!=( const DSDFARev& r ) const;

   friend ostream& operator<<( ostream& os, const DSDFARev& r );

   inline int class_invariant() const;

private:
   StateSet which;
   const DTransRel *T;
   const StateSet *S;
};

// Must always have an argument-less constructor.
inline DSDFARev::DSDFARev()
   whichO,
   T( 0 ),
   S( 0 ) {}
inline DSDFARev::DSDFARev( const DSDFARev& r ) : 
    which( r.which ), 
    T( r.T ), 
    S( r.S ) { 
    assert( class_invariant() ); 
}

// A special constructor:
inline DSDFARev::DSDFARev( const StateSet& rq, 
    const DTransRel* rT, 
    const StateSet* rS ) : 
    which( rq ), 
    T( rT ), 
    S( rS ) { 
    assert( class_invariant() ); 
}

inline const DSDFARev& DSDFARev::operator=( const DSDFARev& r ) { 
    assert( r.class_invariant() ); 
    which = r.which; 
    T = r.T; 
    S = r.S; 
    assert( class_invariant() ); 
    return( *this ); 
}

inline int DSDFARev::finalO const { 
    assert( class_invariant() ); 
    return( which.noLdisjoint( *S ) ); 
}

inline int DSDFARev::operator==( const DSDFARev& r ) const { 
    assert( r.class_invariantO ); 
    assert( T == r.T && S == r.S ); 
    return( which == r.which ); 
}

inline int DSDFARev::operator!=( const DSDFARev& r ) const { 
    return( !operator==( r ) ); 
}

inline int DSDFARev::class_invariantO const { 
    return( T != 0 
        && S != 0 
        && which.domainO == T->domainO 
        && which.domainO == S->domainO ); 
}

#endif

#include "dsdfarev.h"

CRSet DSDFARev::out_labels() const { 
    assert( class_invariant() ); 
    auto CRSet a; 
    // Go through all of the States to see which have transitions 
    // into a State represented by this->which.

Implementation: The set of States in the reversed DFA is represented as a StateSet. A DSDFARev maintains pointers to the components of the DFA from which it was created. Two DSDFAReves are equal if their StateSets are equal.

/∗ (c) Copyright 1994 by Bruce W. Watson ∗/
// $Revision: 1.2$
// $Date: 1994/08/15 13:59:49$
#include "dsdfarev.h"
11.2 Derivatives of regular expressions

auto State i;
for( i = 0; i < which.domain(); i++ ) {
    a.combine( T->labels_between( i, which ) );
}
return( a );

DSDFARev DSDFARev::out_transition( const CharRange a ) const {
    auto StateSet result;
    // Set the correct domain.
    result.set_domain( which.domain() );
    // Find the reverse transitions by going through all State's.
    auto State i;
    for( i = 0; i < which.domain(); i++ ) {
        auto State j( T->transition_on_range( i, a ) );
        // We know that there's a transition from i to j on a.
        if( which.contains( j ) ) {
            // Then we want to add the reverse transition.
            result.add( i );
        }
    }
    // Construct the abstract State on the way out of here:
    return( DSDFARev( result, T, S ) );
}

ostream& operator<< ( ostream& os, const DSDFARev& r ) {
    os << DSDFARev::which << *r.T << *r.S << '\n';
    return( os );
}

11.2 Derivatives of regular expressions

Implementation class: DSRE

Files: dsre.h, dsre.cpp

Uses: CharRange, CRSet, RE

Description: A DSRE represents a derivative of a regular expression. A derivative of a regular expression is again a regular expression; a DSRE simply provides an abstract state interface to the derivatives.
class DSRE : protected RE {
public:
    // Must always have an argument-less constructor.
    inline DSRE();

    // A special constructor:
    inline const DSRE& operator=( const DSRE& r );

    inline const DSRE& operator==( const DSRE& r );

    // The required member functions:
    inline int final() const;
    inline CRSet out_labels() const;
    inline DSRE out_transition( const CharRange a ) const;
    inline int operator===( const DSRE& r ) const;
    inline int operator!=( const DSRE& r ) const;

    friend ostream& operator<<( ostream& os, const DSRE& r );

    inline int class_invariant() const;
};

// Must always have an argument-less constructor.
inline DSRE::DSRE() : RE() {
    assert( class_invariant() );
}

inline DSRE::DSRE( const RE& r ) : RE( r ) {
    assert( r.class_invariant() );
    // They must always be in similarity normal form SNF (see deriv.cpp)
    RE::snf();
    assert( class_invariant() );
}

inline const DSRE& DSRE::operator==( const DSRE& r ) {
    assert( r.class_invariant() );
    RE::operator==( (const RE&)r );
    assert( class_invariant() );
    return( *this );
}

inline int DSRE::final() const {
    assert( class_invariant() );
    return( RE::Null() );
}

inline CRSet DSRE::out_labels() const {
    assert( class_invariant() );
    return( RE::First() );
}

inline DSRE DSRE::out_transition( const CharRange a ) const {
    assert( class_invariant() );
    return( RE::derivative( a ) );
}

inline int DSRE::operator===( const DSRE& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    return( RE::operator===( r ) );
}
11.2 Derivatives of regular expressions

inline int DSRE::operator!=( const DSRE& r ) const {
    return( operator==( r ) );
}

inline int DSRE::class_invariant() const {
    return( RE::class_invariant() && RE::is_snf() );
}

#endif

Implementation: A DSRE inherits from RE. The interface required of abstract states (such as member functions final, out_labels, and out_transition) is provided through the use of the RE member functions such as Null, First, and derivative (respectively). When a new DSRE is constructed, it is put into similarity normal form, to ensure that there will only be finitely many DSREs (see [Wat93a, Theorem 5.31]).

Performance: RE member function derivative returns an RE by value, which is then passed to the DSRE constructor. Use-counting RE could eliminate the overhead of copy constructor and destructor calls.

*/ (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 13:59:54 $
#include "dsre.h"

ostream& operator<<( ostream& os, const DSRE& r ) {
    assert( r.class_invariant() );
    return( os << "\nDSRE\n" << (const RE&)r );
}

Implementation class: DSREopt

Files: dsre-opt.h, dsre-opt.cpp

Uses: CharRange, CRSet, RE

Description: This class is identical to DSRE, with one addition: the underlying regular expression (the derivative) is reduced — redundant subexpressions are eliminated. This makes it more likely that two derivatives (denoting the same language) will be found to be equal (using the DSREopt equality operator) in template function construct_components, thus giving a DFA with fewer States.

*/ (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:35 $
 ifndef DSREOPT_H
 #define DSREOPT_H

 #include "charrang.h"
 #include "crset.h"
 #include "re.h"
 #include <assert.h>
 #include <iostream.h>

 // This class is used to represent abstract States in a DFA that is still under
 // construction. It is used in the construction of a DFA from an RE (using
 // Brzozowski's optimized construction). It differs from DSRE in that the derivatives
 // are optimized (redundant information is removed — see deriv.cpp)
// Most member functions are calls through to the corresponding RE members.  
// DSREopt inherits from RE for implementation.

class DSREopt : protected RE {
public:
    // Must always have an argument-less constructor.
    inline DSREopt();

    // A special constructor:
    inline DSREopt( const RE& r );

    inline const DSREopt& operator=( const DSREopt& r );

    // The required member functions:
    inline int final() const;
    inline CRSet out_labels() const;
    inline DSREopt out_transition( const CharRange a ) const;
    inline int operator==( const DSREopt& r ) const;
    inline int operator!=( const DSREopt& r ) const;

friend ostream& operator<<( ostream& os, const DSREopt& r );

    inline int class_invariant() const;
};

// Must always have an argument-less constructor.
inline DSREopt::DSREopt() : RE() {
    assert( class_invariant() );
}

inline DSREopt::DSREopt( const RE& r ) : RE( r ) {
    assert( r.class_invariant() );
    // They must always be in similarity normal form SNF (see deriv.cpp)
    RE::snf();
    // They must also be reduced.
    RE::reduce();
    assert( class_invariant() );
}

inline const DSREopt& DSREopt::operator=( const DSREopt& r ) {
    assert( r.class_invariant() );
    RE::operator=( (const RE&)r );
    assert( class_invariant() );
    return( *this );
}

inline int DSREopt::final() const {
    assert( class_invariant() );
    return( RE::Null() );
}

inline CRSet DSREopt::out_labels() const {
    assert( class_invariant() );
    return( RE::First() );
}

inline DSREopt::DSREopt::out_transition( const CharRange a ) const {
    assert( class_invariant() );
    return( RE::derivative( a ) );
}

inline int DSREopt::operator==( const DSREopt& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
}
11.3 Item sets

Implementation: Redundant information is eliminated by calling RE member function reduce in the constructor from an RE.

Performance: See DSRE.

Description: ISImpl is a base class implementing an item set [Wat93a, Definition 5.58]. It is used as a base class in abstract state classes DSIS, DSDeRemer, and DSIS_opt. It includes a constructor from a pointer to RE (the regular expression associated with the item set). The pointed-to RE is assumed to remain in existence while the DFA components are constructed by template function construct_components. Additionally, the abstract state members final and out_labels are defined. The member function move_dots takes a CharRange and provides the functionality of "derivatives of item sets" (as defined in [Wat93a, Definition 5.61]), while member function D implements the dot closure relation D defined in [Wat93a, Definition 5.63].
This class is used to represent abstract States in a DFA that is still under construction. It is used as the base class in representing item sets, for the item set construction, DeRemer's construction, and the optimized item set construction. This implements Section 5.5 of the Taxonomy.

class ISImpl {
    protected:
    // Must always have an argument-less constructor.
    inline ISImpl();

    // A special constructor:
    ISImpl(const RE *r);

    const ISImpl &operator=( const ISImpl & r );

    // The required member functions:
    inline int final() const;
    CRSet out_labels() const;

    inline int operator==( const ISImpl & r ) const;
    inline int operator!=( const ISImpl & r ) const;

    // Move the dots across sym nodes.
    inline void move_dots( const CharRange & r );

    friend ostream & operator<( ostream & os, const ISImpl & r );

    inline int class_invariant() const;

    // Some implementation details:
    // Which RE is this with-respect-to:
    const RE * e;

    // Where are the dots:
    // The indices in the BitVec's are the pre-order traversal orders of the nodes in *e.
    // before indicates the nodes with a dot before them;
    // after indicates the nodes with a dot after them.
    BitVec before;
    BitVec after;

    private:
    // Some helpers:

    // Move the dots as far as possible.
    // This is named according to Definition 5.63 of the Taxonomy.
    void D( const RE & r, int & node_num );

    // Fetch the labels throughout the tree.
    // Helper to out_labels().
    CRSet traverse_labels( const RE & r, int & node_num ) const;

    // Move the dots throughout the RE tree:
    // Helper to move_dots().
    void ISImpl::traverse_move_dots( const RE & r, const CharRange & a, int & node_num );
}

// Must always have an argument-less constructor.
// The resulting ISImpl won't satisfy the class_invariant()
inline ISImpl::ISImpl() :
    e( 0 )
{

}
11.3 Item sets

assert( class_invariant() );
// It will be final if there is a dot after the root (which is node 0
// in the traversal order).
// See construction 5.69 of the Taxonomy.
return( after.contains( 0 ) );

inline CRSet ISImpl::out_labels() const {
    assert( class_invariant() );
    // Start at the root, node 0.
    auto int node( 0 );
    return( traverse_labels( *e, node ) );
}

inline int ISImpl::operator==( const ISImpl& r ) const {
    assert( r.class_invariant() );
    // All dots must be in the same places.
    return( (before == r.before) && (after == r.after) );
}

inline int ISImpl::operator!=( const ISImpl& r ) const {
    return( !operator==( r ) );
}

// Move the dots across sym nodes.
inline void ISImpl::move_dots( const CharRange& r ) {
    assert( class_invariant() );
    // Traverse the RE tree.
    // i maintains the node number.
    auto int i( 0 );
    // Start without dots after the nodes.
    after.clear();
    traverse_move_dots( *e, i );
    assert( i == e->num_operators() );

    // Now go back to the root, and do the closure.
    i = 0;
    D( *e, i );
    assert( i == e->num_operators() );
}

inline int ISImpl::class_invariant() const {
    return( e->class_invariant()
            && before.width() == e->num_operators()
            && after.width() == before.width() );
}

#ifndef

Implementation: As explained in [Wat93a, Section 5.5], a dot (an item in an item set) can occur
either before or after a regular expression, or any of its subexpressions. The subexpressions
of the associated regular expression (including the expression itself) are encoded by their
node number (starting at 0) in a prefix traversal of the RE expression tree. The dots before
subexpressions are stored as bits in a BitVec called before, while the dots after subexpressions
are stored as bits in a BitVec, called after (in both cases, the bits indicate if the corresponding
dot is present). Two ISImpls are equal if these two BitVecs are equal, i.e. their dots before
and after are equal. The member functions move_dots and D both make prefix traversals of
the RE expression tree, keeping track of the prefix number of each subexpression.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
CONSTRUCTING A DFA

// $Date: 1994/08/15 11:00:00 $
#include "is-impl.h"

ISImpl::ISImpl( const RE *r ) {
  e( r ) {
    assert( r->class_invariant() );
    auto int i{ r->num_operators() };
    before.set_width( i );
    after.set_width( i );
    assert( !(before.something_set() || after.something_set()) );

    // Now put a dot before the root of *e.
    before.set_bit( 0 );

    // Now propagate the dots.
    auto int nodes( 0 );
    D( *r, nodes );
    assert( nodes == r->num_operators() );
    assert( class_invariant() );
  }
}

const ISImpl& ISImpl::operator=( const ISImpl& r ) {
  assert( r.class_invariant() );
  e = r.e;
  before = r.before;
  after = r.after;
  assert( class_invariant() );
  return( *this );
}

// Move the dots throughout the RE tree:
// Only dots in front a SYMBOL node, matching CharRange a, can advance.
void ISImpl::traverse_move_dots( const RE& r, const CharRange& a, int& node_num ) {
  assert( r.class_invariant() );

  switch( r.root_operator() ) {
  case EPSILON:
    case EMPTY:
      // Remove any dot that may be before.
      before.unset_bit( node_num );
      node_num++; break;
  case SYMBOL:
    // If there's a match, advance the bit.
    if( (a <= r.symbol()) && before.contains( node_num ) ) {
      after.set_bit( node_num );
    }
    before.unset_bit( node_num );
    node_num++; break;
  case OR:
  case CONCAT:
    before.unset_bit( node_num );
    node_num++; traverse_move_dots( r.left_subexpr(), a, node_num ); traverse_move_dots( r.right_subexpr(), a, node_num ); break;
  case STAR:
  case PLUS:
  case QUESTION:
    before.unset_bit( node_num );
    node_num++; traverse_move_dots( r.left_subexpr(), a, node_num ); break;
  }
11.3 Item sets

void ISImpl::D( const RE& r, int& node_num ) {
    // Need to remember if there was a dot in front of this node, and what the preorder number of this node is.
    auto int b;
    auto int nn;
    auto int node_right;
    auto int sub_parity;

    switch( r.root_operator() ) {
    case EPSILON:
        // The dot always moves through an EPSILON.
        if( before.contains( node_num ) ) {
            after.set_bit( node_num );
        }
        node_num++;
        break;
    case EMPTY:
    case SYMBOL:
        // The dot never moves through EMPTY or SYMBOL.
        // Just advance to the next preorder number.
        node_num++;
        break;
    case OR:
        // Is there one before this node?
        b = before.contains( node_num );
        nn = node_num++;
        // If b, the propagate the dot inside the (left of the) OR too.
        if( b ) {
            // node_num is now the node num of the left subexpr.
            before.set_bit( node_num );
        }
        // Propagate through the left subexpression.
        D( r.left_subexpr(), node_num );
        // If the dot is after the left subexpr, then propagate out of the OR.
        if( after.contains( nn + 1 ) ) {
            after.set_bit( nn );
        }
    case CONCAT:
        nn = node_num++;
        // If there's a dot before the CONCAT, propagate it into the left subexpr.
        if( before.contains( nn ) ) {
            before.set_bit( node_num );
        }
        // Propagate through the left subexpr.
        D( r.left_subexpr(), node_num );
        // If the dot is after the left subexpr, move it before the right one.
        if( after.contains( nn + 1 ) ) {
            before.set_bit( node_num );
        }
    case:
        break;
    }
}

// Move the dots as far as possible.
// This is named according to Definition 5.63 of the Taxonomy.
// node_num is the preorder number of r (in the associated RE).

// Need to remember if there was a dot in front of this node, and what the preorder number of this node is.
auto int b;
auto int nn;
auto int node_right;
auto int sub_parity;

switch( r.root_operator() ) {
    case EPSILON:
        // The dot always moves through an EPSILON.
        if( before.contains( node_num ) ) {
            after.set_bit( node_num );
        }
        node_num++;
        break;
    case EMPTY:
    case SYMBOL:
        // The dot never moves through EMPTY or SYMBOL.
        // Just advance to the next preorder number.
        node_num++;
        break;
    case OR:
        // Is there one before this node?
        b = before.contains( node_num );
        nn = node_num++;
        // If b, the propagate the dot inside the (left of the) OR too.
        if( b ) {
            // node_num is now the node num of the left subexpr.
            before.set_bit( node_num );
        }
        // Propagate through the left subexpression.
        D( r.left_subexpr(), node_num );
        // If the dot is after the left subexpr, then propagate out of the OR.
        if( after.contains( nn + 1 ) ) {
            after.set_bit( nn );
        }
    case CONCAT:
        nn = node_num++;
        // If there's a dot before the CONCAT, propagate it into the left subexpr.
        if( before.contains( nn ) ) {
            before.set_bit( node_num );
        }
        // Propagate through the left subexpr.
        D( r.left_subexpr(), node_num );
        // If the dot is after the left subexpr, move it before the right one.
        if( after.contains( nn + 1 ) ) {
            before.set_bit( node_num );
        }
    case:
        break;
}
// Save the node number of the right subexpr.
node_right = node_num;

// Propagate through the right.
D( r.right_subexpr(), node_num );
// If the dot is after the right subexpr, move it behind the CONCAT.
if( after.contains( node_right ) ) {
    after.set_bit( nn );
}

break;

case STAR:

    // A dot before the expression means that it moves after too.
    if( before.contains( node_num ) ) {
        after.set_bit( node_num );
    }

    // Fall through here.

case PLUS:

    // The fall through from the STAR is correct!
nn = node_num++;  // If there's a dot before the STAR/PLUS and into the subexpr.
    if( before.contains( nn ) ) {
        before.set_bit( node_num );
    }

    // Remember if there's already a dot before the subexpr.
    sub_parity = before.contains( node_num );

    // Propagate through the subexpr.
    D( r.left_subexpr(), node_num );

    // Check if a dot is after the subexpr.
    if( after.contains( nn + 1 ) ) {
        // Move the dot out.
        after.set_bit( nn );

        // and move the bit around before the subexpr.
        // if it wasn't already there before.
        if( !sub_parity ) {
            // Re-propagate the dots.

            // nn + 1 is the node num of the subexpr.
            before.set_bit( ++nn );
            D( r.left_subexpr(), nn );

        }

    }

    break;

case QUESTION:

    nn = node_num++;

    // If there's a dot before the QUESTION, move it after, and into the subexpr.
    if( before.contains( nn ) ) {
        after.set_bit( nn );
        before.set_bit( node_num );
    }

    // Propagate through the subexpr.
    D( r.left_subexpr(), node_num );

    // If there's a dot after the subexpr, propagate it out.
    if( after.contains( nn + 1 ) ) {
        after.set_bit( nn );
    }

    break;

}
11.3 Item sets

assert( r.class_invariant() );

auto CRSet b;
switch( r.root_operator() ) {
    case EPSILON:
        node_num++;
        break;
    case EMPTY:
        // Check if this node has a dot before it.
        if( before.contains( node_num ) ) {
            b.append( r.symbol() );
        }
        node_num++;
        break;
    case SYMBOL:
        if( before.contains( node_num ) ) {
            b.append( r.symbol() );
        }
        node_num++;
        break;
    case OR:
    case CONCAT:
        node_num++;
        b = traverse_labels( r.left_subexpr(), node_num );
        b = combine( traverse_labels( r.right_subexpr(), node_num ) );
        break;
    case STAR:
    case PLUS:
    case QUESTION:
        node_num++;
        b = traverse_labels( r.left_subexpr(), node_num );
        break;
}

return( b );

ostream& operator<<( ostream& as, const ISImpl& r ) {
    assert( r.class_invariant() );
    return( as << "\nISImpl\n" << *(r.e) << "\nbefore: " << r.before
      << "\nafter: " << r.after << '\n' );
}

Implementation class: DSIS

Files: dsis.h, dsis.cpp

Uses: CharSet, CRSet, ISImpl, RE

Description: Class DSIS provides the full abstract state interface of an item set (see class ISImpl). DSIS inherits from ISImpl for implementation. It only adds the out_transition member function. It is used in the item set construction [Wat93a, Construction 5.69].
class DSIS : protected ISImpl {}
public:
   // Must always have an argument-less constructor.
   inline DSIS();
   // A special constructor:
   inline DSIS( const RE *r );
   inline const DSIS& operator=( const DSIS& r );
   // The required member functions:
   inline int final() const;
   inline CRSet out_labels() const;
   DSIS out_transition( const CharRange a ) const;
   inline int operator==( const DSIS& r ) const;
   inline int operator!=( const DSIS& r ) const;

   friend ostream& operator<<( ostream& os, const DSIS& r );
   inline int class_invariant() const;
};

// The default constructor won't leave *this in a condition satisfying the
// class invariant.
inline DSIS::DSIS() :
   ISImpl() {
}
inline DSIS::DSIS( const RE *r ) :
   ISImpl( r ) {
   assert( r->class_invariant() );
   assert( class_invariant() );
}
inline const DSIS& DSIS::operator=( const DSIS& r ) {
   assert( r.class_invariant() );
   ISImpl::operator=( (const ISImpl&)r );
   // *this may not satisfy the invariant until after the assignment.
   assert( class_invariant() );
   return( *this );
}
inline int DSIS::final() const {
   assert( class_invariant() );
   return( ISImpl::final() );
}
inline CRSet DSIS::out_labels() const {
   assert( class_invariant() );
   return( ISImpl::out_labels() );
}
inline int DSIS::operator==( const DSIS& r ) const {
   assert( class_invariant() );
   assert( r.class_invariant() );
   return( ISImpl::operator==( r ) );
}
inline int DSIS::operator!=( const DSIS& r ) const {
   return( !operator==( r ) );
}
inline int DSIS::class_invariant() const {
   // Should really check that before and after are fully closed.
   return( ISImpl::class_invariant() );
Implementation: Most of the implementation is provided by class ISImpl.

```
#include "dsis.h"

DSIS::out_transition( const CharRange a ) const {
    assert( class_invariant() );
    auto DSIS ret( *this );
    ret.move_dots( a );
    return( ret );
}

ostream& operator<<( ostream& os, const DSIS& r ) {
    assert( r.class_invariant() );
    return( os << "\nDSIS\n" << (const ISImpl&)r );
}
```

Implementation class: DSDeRemer

Files: dsderem.h, dsderem.cpp

Uses: CharRange, CRSet, ISImpl, RE

Description: Class DSDeRemer provides the abstract state interface of an item set (see class ISImpl). DSDeRemer inherits from ISImpl for implementation. It is used in DeRemer's construction [Wat93a, Construction 5.75]. DeRemer's construction can yield smaller DFAs than the item set construction.

```
#include "dsderem.h"
#include "crset.h"
#include "re.h"
#include "is-impl.h"
#include <assert.h>
#include <iostream.h>

II

This class is used to represent abstract States in a DFA that is still under
construction. It represents the item sets in the DeRemer construction.
See Construction 5.75
DSDeRemer inherits from ISImpl for implementation.

class DSDeRemer : protected ISImpl { 20
public:
    // Must always have an argument-less constructor.
    inline DSDeRemer();
    // A special constructor:
    inline DSDeRemer( const RE& r );
```
inline const DSDeRemer& operator=( const DSDeRemer& r );

// The required member functions:
inline int final() const;
inline CRSet out_labels() const;
DSDeRemer out_transition( const CharRange a ) const;
inline int operator==( const DSDeRemer& r ) const;
inline int operator!=( const DSDeRemer& r ) const;

friend ostream& operator<<( ostream& os, const DSDeRemer& r );

inline int class_invariant() const;

private:
// Helpers: the optimization function of the DeRemer construction.
inline void opt();
void traverse_opt( const RE& r, int& i );

inline DSDeRemer::DSDeRemer() :
    ISImpl() {
}

inline DSDeRemer::DSDeRemer( const RE *r ) :
    ISImpl( r ) {
    assert( r->class_invariant() );
    // The ISImpl constructor will have closed before/after, but not optimized yet.
    opt();
    assert( class_invariant() );
}

inline const DSDeRemer& DSDeRemer::operator==( const DSDeRemer& r ) {
    assert( r.class_invariant() );
    ISImpl::operator==( (const ISImpl&)r );
    assert( class_invariant() );
    return( *this );
}

inline int DSDeRemer::final() const {
    assert( class_invariant() );
    return( ISImpl::final() );
}

inline CRSet DSDeRemer::out_labels() const {
    assert( class_invariant() );
    return( ISImpl::out_labels() );
}

inline int DSDeRemer::operator==( const DSDeRemer& r ) const {
    assert( class_invariant() );
    return( ISImpl::operator==( r ) );
}

inline int DSDeRemer::operator!=( const DSDeRemer& r ) const {
    return( !operator==( r ) );
}

inline void DSDeRemer::opt() {
    auto int i( 0 );
    traverse_opt( *e, i );
    assert( i == (e->num_operators() ) );
}

inline int DSDeRemer::class_invariant() const {
    // Should really check that before and after are fully closed,
11.3 Item sets

// and that they are optimized.
return( ISImpl::class_invariant() );
}
#endif

**Implementation:** The implementation is largely provided by the base class ISImpl. The member function `opt` (with helper `traverse_opt`) implements DeRemer's dot filter function $X$ as defined in [Wat93a, Definition 5.72]. The filter function makes a single prefix traversal of the RE expression tree. By applying the filter, it becomes more likely that template function `construct_components` will find two item sets (denoting the same language) to be equal (this will result in fewer states in the constructed DFA).

```cpp
/* (c) Copyright 1993 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1993/08/15 19:59:48 $
#include "dsderem.h"

DSDeRemer DSDeRemer::out_transition( const CharSet& a ) const {
    assert( class_invariant() );
    auto DSDeRemer ret( *this );
    ret.move_dots( a );
    ret.opt();
    return( ret );
}

void DSDeRemer::traverse_opt( const RE& r, int& node_num ) {
    assert( r.class_invariant() );
    // Recursively traverse r and apply the optimizations of Defn. 5.72
    // (function X) of the Taxonomy.
    switch( r.root_operator() ) {  
        case EPSILON:
        case EMPTY:
        case SYMBOL:
            node_num++;
            break;
        case OR:
            before.unset_bit( node_num );
            // Fall through is good.
        case CONCAT:
            node_num++;
            traverse_opt( r.left_subexpr(), node_num );
            traverse_opt( r.right_subexpr(), node_num );
            break;
        case STAR:
            before.unset_bit( node_num );
            node_num++;
            after.unset_bit( node_num );
            traverse_opt( r.left_subexpr(), node_num );
            break;
        case PLUS:
        case QUESTION:
            node_num++;
            traverse_opt( r.left_subexpr(), node_num );
            break;
    }
    ostream& operator<< ( ostream& os, const DSDeRemer& r ) {
        assert( r.class_invariant() );
        return( os << "\n" << (const ISImpl&): );
    }
```
Implementation class: **DSISopt**

Files: dsis-opt.h, dsis-opt.cpp

Uses: **CharRange, CRSet, ISImpl, RE**

Description: Class **DSIS.opt** provides the abstract state interface of an item set (see class **ISImpl**). **DSIS.opt** inherits from **ISImpl** for implementation. It is used in the optimized item set construction [Wat93a, Construction 5.82]. The optimized item set construction (called Oconstr in [Wat93a]) can yield much smaller DFAs than either the normal item set construction or DeRemer's construction.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:32 $
#ifndef DSIS_OPT_H
#define DSIS_OPT_H

#include "charrang.h"
#include "crset.h"
#include "re.h"
#include "is-impl.h"
#include <assert.h>
#include <iostream.h>

// This class is used to represent abstract States in a DFA that is still under
// construction. It represents the item sets in the DeRemer construction.
// See Construction 5.75
// DSIS_opt inherits from ISImpl for implementation.

class DSIS_opt : protected ISImpl {
  public:
    // Must always have an argument-less constructor.
    inline DSIS_opt();

    // A special constructor:
    inline DSIS_opt( const RE *r );

    inline const DSIS_opt& operator=( const DSIS_opt& r );

    // The required member functions:
    inline int final() const;
    inline CRSet out_labels() const;
    DSIS_opt out_transition( const CharRange a ) const;
    inline int operator==( const DSIS_opt& r ) const;
    inline int operator!=( const DSIS_opt& r ) const;

    friend ostream& operator<<( ostream& os, const DSIS_opt& r );

    inline int class_invariant() const;

  private:
    // Helpers: the optimization function of the DeRemer construction.
    inline void opt();
    void traverse_opt( const RE& r, int& i );
  }

inline DSIS_opt::DSIS_opt() : ISImpl() {
}

inline DSIS_opt::DSIS_opt( const RE *r ) :
```
Implementation: The implementation is largely provided by the base class ISImpl. The member function opt (with helper traverse_opt) implements the optimized dot filter function \( Y \) as defined in [Wat93a, Definition 5.79]. The filter function makes a single prefix traversal of the RE expression tree. By applying the filter, it becomes more likely that template function construct_components will find two item sets (denoting the same language) to be equal. The reason why the optimized item set construction yields smaller DFAs than DeRemer's construction is that the \( Y \) filter is more effective at removing redundant information than the \( X \) filter used in DeRemer's construction (see [Wat93a, Remark 5.81]).
DSIS_opt DSIS_opt::out_transition( const CharRange a ) const {
    assert( class_invariantO );
    auto DSIS_opt ret(*this);
    ret.move_dots( a );
    ret.opt();
    return( ret );
}

void DSIS_opt::traverse_opt( const RE& r, int& node_num ) {
    assert( class_invariantO );
    assert( r.class_invariant() );
    switch( r.root_operation() ) {
    case EPSILON:
        before.unset_bit( node_num );
        node_num++;
        break;
    case EMPTY:
        node_num++;
        break;
    case SYMBOL:
        // Don't remove a bit before — see Definition 5.79
        before.unset_bit( node_num );
        node_num++;
        break;
    case OR:
    case CONCAT:
        before.unset_bit( node_num );
        node_num++;
        traverse_opt( r.left_subexpr(), node_num );
        traverse_opt( r.right_subexpr(), node_num );
        break;
    case STAR:
    case PLUS:
    case QUESTION:
        before.unset_bit( node_num );
        node_num++;
        traverse_opt( r.left_subexpr(), node_num );
        break;
    }
    return( os << "nDSIS_opt\n" << (const ISImpl& ) r );
}
12 Finite automata

Implementation class: FA

Files: fa.h, fa.cpp

Uses: FAabs, RE, State, StatePool, StateRel, StateSet, TransRel

Description: Class FA implements finite automata as defined in [Wat93a, Definition 2.1]. It inherits from FAabs, and implements the interface defined there. A constructor taking an RE is provided, implementing Thompson's top-down construction [Wat93a, Construction 4.5]. This constructor can be more efficient than using the $\Sigma$-algebra $\text{Reg}<\text{FA}>$.

```c
/* (c) Copyright 1994 by Bruce Watson */
// $Revision: 1.2 $
// $Date: 1994/06/15 14:00:41 $
#ifndef FA_H
#define FA_H
#include "faabs.h"
#include "st-pool.h"
#include "stateset.h"
#include "staterel.h"
#include "transrel.h"
#include "re.h"
#include <assert.h>
#include <iostream.h>

// Implement general finite automata (Definition 2.1 of the Taxonomy).
// These can be constructed by the Sigma-algebra (see sig-fa.cpp) or using
// the special constructor from RE's.

class FA : virtual public FAabs {
public:
    // Constructors, destructors, operator=:
    FA();

    // Provide a copy constructor:
    FA( const FA& r );

    // A special constructor, implementing the Thompson's top-down construction
    // (see Construction 4.5).
    FA( const RE& e );

    // Default operator= is okay.

    // Basic FAabs member functions (without overriding acceptable( )):

    virtual int num_states() const;
    virtual void restart();
    virtual void advance( char a );
    virtual int in_final() const;
    virtual int stuck();
    virtual DFA determinism() const;

    // Special member functions:

    friend ostream& operator<<( ostream& os, const FA& r );

    inline int class_invariant() const;

protected:
    // Functions states_reqd, td (see Construction 4.5) for use in the constructor from RE.
    int states_reqd( const RE& e );
    void td( const State s, const RE& e, const State f );
};
```
// recycle this FA:
void reincarnate();

// Implementation details:

StatePool Q;
StateSet S, F;

// Transitions maps each State to its out-transitions.
TransRel Transitions;
// E is the epsilon transition relation.
StateRel E;

// Simulation stuff:
StateSet current;

inline int FA::class_invariant() const {
// Should also check that current is E-closed.
return (Q.size() == S.domain() && Q.size() == F.domain() &&
Q.size() == Transitions.domain() && Q.size() == E.domain() &&
Q.size() == current.domain() &&
S.class_invariant() &&
F.class_invariant() &&
Transitions.class_invariant() &&
E.class_invariant() &&
current.class_invariant());
}

#endif

Implementation: Most of the member functions are implemented in a straight-forward manner.
The constructor from an RE makes a traversal of the RE to determine the number of States
to allocate before-hand.

Performance: Since Thompson's construction yields finite automata with certain structural
properties (see [Wat93a, end of Definition 4.1]), a special "Thompson's construction" automata class could be efficiently implemented, instead of the general finite automaton structure defined in FA.
current.set_domain( Q.size() );
assert( class_invariant() );
}

FA::FA( const FA& r ) :
Q( r.Q ),
S( r.S ),
F( r.F ),
Transitions( r.Transitions ),
E( r.E ),
current( r.current ) {
assert( class_invariant() );
}

FA::FA( const RE& e ) {
    // First, allocate enough states by setting up the domains.
    // See Section 4.1.1 of the Taxonomy.
    auto int states_required( states_reqd( e ) );
    S.set_domain( states_required );
    F.set_domain( states_required );
    Transitions.set_domain( states_required );
    E.set_domain( states_required );
    current.set_domain( states_required );

    auto State s( Q.allocate() );
    auto State f( Q.allocate() );
    S.add( s );
    F.add( f );
    Id( s, e, f );
    assert( class_invariant() );
}

    // The following follows directly from inspecting Thompson's construction
    // (Definition 4.1).
    int FA::states_reqd( const RE& e ) {
        assert( e.class_invariant() );

        auto int ret;
        switch( e.root_operator() ) {
            case EPSILON:
                ret = 2;
                break;
            case EMPTY:
                ret = 2 + states_reqd( e.left_subexpr() ) + states_reqd( e.right_subexpr() );
                break;
            case OR:
                ret = 2 + states_reqd( e.left_subexpr() ) + states_reqd( e.right_subexpr() );
                break;
            case CONCAT:
                ret = states_reqd( e.left_subexpr() ) + states_reqd( e.right_subexpr() );
                break;
            case STAR:
            case PLUS:
            case QUESTION:
                ret = 2 + states_reqd( e.left_subexpr() );
                break;
        }
        return( ret );
    }

    int FA::num_states( ) const {
        return( Q.size() );
    }

    // The simulation details follow directly from Section 2 of the Taxonomy.
    void FA::restart() {
        assert( class_invariant() );
        current = E.closure( S );
    }
void FA::advance( char a ) {
    assert(class_invariant());
    // Compute the epsilon-closure.
    current = E.closure( Transitions.image( current, a ) );
    assert( current == E.closure( current ) );
    assert(class_invariant());
}

int FA::in_final() const {
    assert(class_invariant());
    return( current.not_disjoint( F ) );
}

int FA::stuck() {
    assert(class_invariant());
    return( current.empty() );
}

DFA FA::determinism() const {
    // Make sure that *this is structurally sound.
    assert(class_invariant());
    // Now construct the DFA components.
    return( DFA( construct_components( DSFA( S, &Transitions, &E, &F ) ) ) );
}

ostream& operator<< ( ostream& os, const FA& r ) {
    assert( r.class_invariant() );
    os << "\nFA\n";
    os << "Q = " << r.Q << "\n";
    os << "S = " << r.S << "\n";
    os << "F = " << r.F << "\n";
    os << "Transitions = \n" << r.Transitions << "\n";
    os << "current = " << r.current << "\n";
    return( os );
}

void FA::td( const State S, const RE& e, const State F ) {
    assert(e.class_invariant());
    // Implement function td (of Construction 4.5).
    // Construct an FA to accept the language of e, between S and F.
    switch( e.root_operator() ) {
    case EPSILON:
        E.union_cross( S, F );
        break;
    case EMPTY:
        break;
    case SYMBOL:
        Transitions.add_transition( S, e.symbol(), F );
        break;
    case OR:
    {
        auto State p( Q.allocate() );
        auto State q( Q.allocate() );
        td( p, e.left_subexpr(), q );
        auto State r( Q.allocate() );
        auto State t( Q.allocate() );
        td( r, e.right_subexpr(), t );
        E.union_cross( S, p );
        E.union_cross( S, r );
        E.union_cross( q, F );
    }
E.union_cross( t, f );
} break;
case CONCAT:
{
    auto State p( Q.allocate() );
    auto State q( Q.allocate() );
    td( s, e.left_subexpr(), p );
    td( q, e.right_subexpr(), f );
    E.union_cross( p, q );
} break;
case STAR:
{
    auto State p( Q.allocate() );
    auto State q( Q.allocate() );
    td( p, e.left_subexpr(), q );
    E.union_cross( s, p );
    E.union_cross( q, p );
    E.union_cross( q, f );
    E.union_cross( s, f );
} break;
case PLUS:
{
    auto State p( Q.allocate() );
    auto State q( Q.allocate() );
    td( p, e.left_subexpr(), q );
    E.union_cross( s, p );
    E.union_cross( q, p );
    E.union_cross( q, f );
} break;
case QUESTION:
{
    auto State p( Q.allocate() );
    auto State q( Q.allocate() );
    td( p, e.left_subexpr(), q );
    E.union_cross( s, p );
    E.union_cross( q, f );
    E.union_cross( s, f );
} break;
}

void FA::reincarnate() {
    Q.reincarnate();
    S.reincarnate();
    F.reincarnate();
    Transitions.reincarnate();
    E.reincarnate();
    current.reincarnate();
}

Implementation class: DSFA

Files: dsfa.h, dsfa.cpp

Uses: CharRange, CRSet, StateRel, StateSet, TransRel
Description: Class DSFA implements the abstract class interface required to construct a DFA from an FA. A DSFA is constructed in the FA member function determinism, and then passed to template function construct_components, which constructs the DFA components.

```cpp
/* (c) Copyright 1994 by Bruce W. Watson */
/* $Revision: 1.2 $ */
/* $Date: 1994/08/15 14:00:31 $ */
#ifndef DSFA_H
#define DSFA_H
#include "charrang.h"
#include "crset.h"
#include "stateset.h"
#include "staterel.h"
#include "transrel.h"
#include <cassert.h>
#include <iostream.h>

This class is used to represent abstract States in a DFA that is still under construction. It is used in the construction of a DFA from an FA. The behaviour of a DSFA follows the simulation behaviour of an FA as in the subset construction.

class DSFA {
    public:
        // Must always have an argument-less constructor.
        inline DSFA();

        inline DSFA( const DSFA& r );

        // A special constructor:
        DSFA( const StateSet& rq,
              const TransRel* T,
              const StateRel* E,
              const StateSet* F );

        inline const DSFA& operator=( const DSFA& r );

        // The required member functions:
        inline int final( const CRSet out_labels() const;
        DSFA out_transition( const CharRange a ) const;
        inline int operator==( const DSFA& r ) const;
        inline int operator!=( const DSFA& r ) const;

        friend ostream& operator<<( ostream& os, const DSFA& r );
        inline int class_invariant() const;

    private:
        StateSet which;
        const TransRel* T;
        const StateRel* E;
        const StateSet* F;
    };

    // Must always have an argument-less constructor.
    inline DSFA::DSFA() :
        which(),
        T( 0 ),
        E( 0 ),
        F( 0 ) {
        // No assertion of the class invariant, since it won't qualify.
    }

    inline DSFA::DSFA( const DSFA& r ) :
```

Implementation: A DSFA contains a StateSet, representing the set of FA States that make up a DFA state in the subset construction (see [Wat93a, p. 12-13] for an explanation of the subset construction). The implementation follows directly from the way in which an FA processes an input string.
128

assert( class_invariant( ) );
}

CRSet DSFA::out_labels() const {
    assert( class_invariant( ) );
    return( T->out_labels( which ) );
}

DSFA DSFA::out_transition( const CharRange a ) const {
    assert( class_invariant( ) );
    /\ Construct the abstract State on the way out of here:
    /\ Assume that which is already E closed, then transition on CharRange
    /\ a; the StateSet gets E closed again in the constructor.
    return( DSFA( T->transition_on_range( which, a ), T, E, F ) );
}

ostream& operator< <( ostream& os, const DSFA& r ) {
    os << "\nDSFA\n" << r.which << "\nT << "\nE << "\nF << '\n';
    return( os );
}

12.1 The \( \Sigma \)-algebra \( \text{Reg}<\text{FA}> \)

Implementation class: \( \text{Reg}<\text{FA}> \)

Files: sig-fa.cpp

Uses: \text{CharRange}, \text{FA}, \text{Reg}, \text{State}

Description: The template instantiation \( \text{Reg}<\text{FA}> \) implements Thompson’s \( \Sigma \)-algebra of finite automata, as defined in [Wat93a, Definition 4.1]. The operators can be used directly to construct complex finite automata, or the \( \text{Reg}<\text{FA}> \) constructor from \( \mathcal{R} \mathcal{E} \) (the constructor defined in class \( \text{Reg} \)) can be used to construct the homomorphic image of the regular expression (the constructor uses these operators indirectly).

Implementation: The implementation follows directly from [Wat93a, Definition 4.1]. In some places it is complicated by the management of domains of \text{StateSets}, \text{StateRels}, and \text{TransReIs}. 

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:19 $
#include "charrang.h"
#include "sigma.h"
#include "fa.h"

// Implement the Sigma-algebra operators (Definition 4.29 of the Taxonomy).
Reg<FA>& Reg<FA>::epsilonO {
    // *this may have been something in a previous life.
    // Wipe out all previous components;
    reincarnate();
    auto State s( Q.allocate( ) );
    auto State f( Q.allocate( ) );
    S.set_domain( Q.size() );
    S.add( s );
    F.set_domain( Q.size() );
    F.add( f );
    Transitions.set_domain( Q.size() );
12.1 The $\Sigma$-algebra Reg<FA>

```cpp
Reg<FA>& Reg<FA>::empty() {
    // See epsilon case.
    reincarnate();
    auto State s = Q.allocate();
    auto State f = Q.allocate();
    S.set_domain( Q.size() );
    S.add( s );
    F.set_domain( Q.size() );
    F.add( f );
    Transitions.set_domain( Q.size() );
    E.set_domain( Q.size() );
    current.set_domain( Q.size() );
    assert( class_invariant() );
    return( *this );
}

Reg<FA>&& Reg<FA>::symbol( const CharRange r ) {
    // See epsilon case.
    reincarnate();
    auto State s = Q.allocate();
    auto State f = Q.allocate();
    S.set_domain( Q.size() );
    S.add( s );
    F.set_domain( Q.size() );
    F.add( f );
    Transitions.set_domain( Q.size() );
    Transitions.add_transition( s, r, f );
    E.set_domain( Q.size() );
    current.set_domain( Q.size() );
    assert( class_invariant() );
    return( *this );
}

Reg<FA>&& Reg<FA>::or( const Reg<FA>& r ) {
    assert( class_invariant() );
    Q.incorporate( r.Q );
    auto State s = Q.allocate();
    auto State f = Q.allocate();
    S.disjointing_union( r.S );
    assert( class_invariant( r ) );
    return( *this );
}
```

S.set_domain(Q.size());
F.disjointing_union(r.F);
F.set_domain(Q.size());
Transitions.disjointing_union(r.Transitions);
Transitions.set_domain(Q.size());
E.disjointing_union(r.E);
E.set_domain(Q.size());
E.union_cross(s, S);
E.union_cross(F, f);
S.clear();
S.add(s);
F.clear();
F.add(f);
current.set_domain(Q.size());
assert(class_invariant());
return(*this);

Reg<FA>& Reg<FA>::concat(const Reg<FA>& r) {
assert(class_invariant());
assert(r.class_invariant());
// See the or operator.
// All state-related stuff in r must be adjusted.

// Save the old domain().
auto int olddom(Q.size());

// Incorporate the other StatePool.
Q.incorporate(r.Q);

// Adjust the domain of the Start states.
S.set_domain(Q.size());

// Transitions are just unioned.
Transitions.disjointing_union(r.Transitions);

// The epsilon trans. are unioned and F times r.S are added.
E.disjointing_union(r.E);
F.set_domain(Q.size());

// r.S will be needed for epsilon transitions.
auto StateSet S1(r.S);
// rename it to the new StatePool size.
S1.rename(olddom);
E.union_cross(F, S1);

// F remains the Final states.
F = r.F;
F.set_rename(olddom);
current.set_domain(Q.size());
assert(class_invariant());
return(*this);
}

Reg<FA>& Reg<FA>::star() {
assert(class_invariant());

// Create some new States, and adjust the domains.
auto State s( Q.allocate() );
auto State f( Q.allocate() );

S.set_domain( Q.size() );
F.set_domain( Q.size() );
Transitions.set_domain( Q.size() );
E.set_domain( Q.size() );
current.set_domain( Q.size() );

E.union_cross( s, S );
E.union_cross( s, f );
E.union_cross( F, S );
E.union_cross( F, f );

S.clear();
S.add( s );
F.clear();
F.add( f );

assert( class_invariant() );
return( *this );

Reg<FA>& Reg<FA>::plus() {
    assert( class_invariant() );

    // Create some new States, and adjust the domains.
    auto State s( Q.allocate() );
    auto State f( Q.allocate() );

    S.set_domain( Q.size() );
    F.set_domain( Q.size() );
    Transitions.set_domain( Q.size() );
    E.set_domain( Q.size() );
    current.set_domain( Q.size() );

    E.union_cross( s, S );
    E.union_cross( s, f );
    E.union_cross( F, S );
    E.union_cross( F, f );

    S.clear();
    S.add( s );
    F.clear();
    F.add( f );

    assert( class_invariant() );
    return( *this );
}

Reg<FA>& Reg<FA>::question() {
    assert( class_invariant() );

    // Create some new States, and adjust the domains.
    auto State s( Q.allocate() );
    auto State f( Q.allocate() );

    S.set_domain( Q.size() );
    F.set_domain( Q.size() );
    Transitions.set_domain( Q.size() );
    E.set_domain( Q.size() );
    current.set_domain( Q.size() );

    E.union_cross( s, S );
    E.union_cross( s, f );
    E.union_cross( F, f );
S.clear();
S.add( s );
F.clear();
F.add( f );

assert( class_invariant() );
return( *this );
13 Reduced finite automata

Implementation class: \textit{RFA}

Files: \texttt{rfa.h}, \texttt{rfa.cpp}

Uses: \texttt{DFA}, \texttt{FAabs}, \texttt{LBFA}, \texttt{RBFA}, \texttt{RE}, \texttt{StatePool}, \texttt{StateRel}, \texttt{StateSet}, \texttt{Trans}

Description: Class \textit{RFA} implements reduced finite automata, as defined in [Wat93a, Definition 4.24]. It inherits from \textit{FAabs}, and implements the interface defined by abstract base \textit{FAabs}. A constructor taking an \texttt{RE} provides a very efficient top-down implementation of homomorphism \texttt{rfa} defined in [Wat93a, Definition 4.30].

```c
/* (c) Copyright 1994 by Bruce Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:48 $
#ifndef RFA_H
#define RFA_H

#include "faabs.h"
#include "st-pool.h"
#include "stateset.h"
#include "staterel.h"
#include "trans.h"
#include "re.h"
#include "dfa.h"
#include <iostream.h>

// The existence of LBFA and RBFA are needed, for friendship.
// Can't include rфа.h or rфа.h due to circularity.
class LBFA;
class RBFA;

// Implement RFA's (reduced finite automata -- Definition 4.24 of the Taxonomy)
// as a type of FAabs.
class RFA : virtual public FAabs {
public:
  // Some constructors:
  RFA();
  // Need a copy constructor:
  RFA( const RFA& r );
  // The Sigma-homomorphism constructor (see Definition 4.30):
  RFA( const RE& e );
  // Default destr, operator= are okay.
  // Normal FAabs member functions (don't override acceptable()):
  virtual int num_states() const;
  virtual void restart();
  virtual void advance( char a );
  virtual int in_final() const;
  virtual int stuck();
  virtual DFA determinism() const;

  // An alternate implementation of determinism() using the LBFA interpretation.
  virtual DFA determinism2() const;

  // Special members:
```
friend ostream& operator<<(ostream& os, const RFA& r);

// These two must be friends so that they can construct themselves properly.
friend class LBFA;
friend class RBFA;

inline int class_invariant() const;

protected:

// A helper: compute all of the components into place.
// for use in the constructor.
void rfa_into(const RE& e, StatePool& Qp, StateSet& f, StateSet& l, Trans& Qm, StateRel& foll, int& N);

// Recycle this:
void reincarnate();

// Implementation details, protected for Reg<RFA> use.
StatePool Q;
StateSet first, last;
Trans Qmap_inverse;
StateRel follow;
int Nullable;

// Some simulation stuff.
int final;
StateSet current;

inline int RFA::class_invariant() const {
    return (Q.size() == first.domain() &&
            Q.size() == last.domain() &&
            Q.size() == Qmap_inverse.range() &&
            Q.size() == follow.domain() &&
            first.class_invariant() &&
            last.class_invariant() &&
            Qmap_inverse.class_invariant() &&
            follow.class_invariant() &&
            current.class_invariant() );
}

#endif

Implementation: Most of the member functions are implemented in a straight-forward manner.

In implementing RFA, a choice must be made between the two interpretations of States: the left-biased or the right-biased interpretation (see the bottom of p. 73 in [Wat93a] for a brief explanation of the different interpretations). The interpretation chosen here is the right-biased one, since it appears to be the most efficient one in practice.

/* (c) Copyright 1994 by Bruce Watson */
#include "rfa.h"
#include "dsrfa.h"
#include "dsrfa2.h"
#include "dfaseed.h"
RFA::RFA() :
    Nullable( 0 ) { 
        assert( class_invariant() );
    }

RFA::RFA( const RFA& r ) :
    Q( r.Q ),
    first( r.first ),
    last( r.last ),
    Qmap_inverse( r.Qmap_inverse ),
    follow( r.follow ),
    Nullable( r.Nullable ),
    final( r.final ),
    current( r.current ) {
        assert( class_invariant() );
    }

    // The Sigma-homomorphism constructor (see Definition 4.30):
    RFA::RFA( const RE& e ) {
        assert( e.class_invariant() );
        auto int size( e.num_symbols() );
        first.set_domain( size );
        last.set_domain( size );
        Qmap_inverse.set_range( size );
        follow.set_domain( size );

        // Now compute all of the components into place in *this
        rfa_into( e, Q, first, last, Qmap_inverse, follow, Nullable );
        current.set_domain( size );
        assert( class_invariant() );
    }

    // Provide implementation for all of the FAs required members.
    int RFA::num_states() const {
        return( Q.size() );
    }

    // The following functions make use of the RBFA interpretation of states
    // (Section 4.3 of the Taxonomy).
    void RFA::restart() {
        assert( class_invariant() );
        // Ready to see any char corresponding to a first.
        current = first;
        // Only acceptable if *this accepts epsilon (ie. Nullable is true).
        final = Nullable;
        assert( class_invariant() );
    }

    void RFA::advance( char a ) {
        assert( class_invariant() );

        // Figure out which are valid out-transitions.
        current.intersection( Qmap_inverse[a] );

        // current is the set of States we can really be in to see a.
        // We're accepting if at least one of the current States is also a last.
        final = current.not_disjoint( last );
        current = follow.image( current );
        assert( class_invariant() );
    }

    int RFA::in_final() const {
        assert( class_invariant() );
        return( final );
    }
13 REDUCED FINITE AUTOMATA

int RFA::stuck() {
    assert( class_invariant() );
    return( current.empty() );
}

DFA RFA::determinism() const {
    assert( class_invariant() );

    // Now construct the DFA components.
    return( construct_components( DSRFA( first,
                                          &Qmap_inverse,
                                          &follow,
                                          &first,
                                          &last,
                                          Nullable ) ) );
}

DFA RFA::determinism2() const {
    assert( class_invariant() );

    // Now construct, using the right-biased interpretation.
    return( construct_components( DSRFA2( &Qmap_inverse,
                                           &follow,
                                           &first,
                                           &last,
                                           Nullable ) ) );
}

// Implement a top down version of Sigma homomorphism rfa (Definition 4.36)
void RFA::rfa_into( const RE& e,
                    StatePool& Qp,
                    StateSet& f,
                    StateSet& l,
                    Trans& Qm,
                    StateRel& foll,
                    int& N ) {

    assert( e.class_invariant() );

    switch( e.root_operator() ) {
        case EPSILON:
            // Nothing to do, but make it nullable.
            N = 1;
            return;
        case EMPTY:
            N = 0;
            return;
        case SYMBOL:
            // There is a symbol here, so give it a State.
            {
                auto State q( Qp.allocate() );
                f.add( q );
                l.add( q );
                Qm.add_transition( e.symbol(), q );
                // Nothing to do to follow.
                N = 0;
            }
            return;
        case OR:
        case CONCAT:
            // Lots of common code, so do the two together.
            {
                // Compute the rfa of e.left_subexpr() into *this.
                rfa_into( e.left_subexpr(), Qp, f, l, Qm, foll, N );

                // Use some locals to store the pieces of the RFA of
// the right subexpression.
auto StateSet fr;
fr.set_domain( f.domain() );
auto StateSet lr;
lr.set_domain( l.domain() );
auto Trans Qmr;
Qmr.set_range( Qm.range() );
auto StateRel follr;
follr.set_domain( foll.domain() );
auto int Nr;

// Now compute the RFA components corresponding to the right subexpr.
// Note: states are still taken from Qp.
rfa_into( e.right_subexpr(), Qp, fr, lr, Qmr, follr, Nr );

Qm.set_union( Qmr );
if( e.root_operator() == OR ) {
    f.set_union( fr );
    l.set_union( lr );
    foll.set_union( follr );
    N = N || Nr;
} else {
    // This is a CONCAT.
    assert( e.root_operator() == CONCAT );
    foll.set_union( follr );
    foll.union_cross( l, fr );

    if( N ) {
        f.set_union( fr );
    }
    if( Nr ) {
        l.set_union( lr );
    } else {
        l = lr;
    }
    N = N && Nr;
}
}
return;
case STAR:
case PLUS:
    // First, do the subexpression.
    rfa_into( e.left_subexpr(), Qp, f, l, Qm, foll, N );
    foll.union_cross( l, f );
    if( e.root_operator() == STAR ) {
        N = 1;
    } else {
        assert( e.root_operator() == PLUS );
    }
}
return;
case QUESTION:
    rfa_into( e.left_subexpr(), Qp, f, l, Qm, foll, N );
    N = 1;
return;
}

void RFA::reincarnate() {
    Q.reincarnate();
    first.reincarnate();
    last.reincarnate();
    Nullable = final = 0;
    current.reincarnate();
}
13 REDUCED FINITE AUTOMATA

```cpp
ostream& operator<<( ostream& os, const RFA& r ) {
    assert( r.class_invariant() );
    os << "\nRFA\n";
    os << "Q = " << r.Q << \n;  
    os << "first = " << r.first << \n;  
    os << "last = " << r.last << \n;  
    os << "Qmap_inverse = " << r.Qmap_inverse << \n;  
    os << "follow = \n" << r.follow << \n;  
    os << "Nullable = " << (r.Nullable ? "true" : "false");
    os << "unfinal = " << (r.final ? "true" : "false");
    os << "incorrect = " << r.current << \n;  
    return( os );
}
```

Implementation class: DSRFA

Files: dsrf a.h, dsrf a.cpp

Uses: CharRange, CRSet, StateRel, StateSet, Trans

Description: Class DSRFA implements the abstract class interface required to construct a DFA from an RFA. A DSRFA is constructed in the RFA member function determinism, and then passed to template function construct_components, which constructs the components of the DFA.

```c
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $  
// $Date: 1994/08/15 14:00:37 $
#define DSRFA_H

#include "charrang.h"
#include "crset.h"
#include "stateset.h"
#include "staterel.h"
#include "transrel.h"
#include <assert.h>
#include <iostream.h>

class DSRFA {  
public:
    // Must always have an argument-less constructor.
    inline DSRFA();  
    inline DSRFA( const DSRFA& r );
    // A special constructor:
    DSRFA( const StateSet& rq,  
           const Trans +Qmap_inverse,  
           const StateRel +follow,  
           const StateSet +first,  
           const StateSet +last,  
           const int rfinalness );

    inline const DSRFA& operator=( const DSRFA& r );
    // The required member functions:
    inline int final() const;  
    CRSet out_labels() const;
};
```
DSRFA out_transition( const CharSet a ) const;
inline int operator==( const DSRFA& r ) const;
inline int operator!=( const DSRFA& r ) const;

friend ostream& operator<<( ostream& os, const DSRFA& r);
inline int class_invariantO const;

private:
    // This stuff may vary from object to object.
    StateSet which;
    int finalness;
    // This stuff should be the same for all objects corresponding
    // to a particular RFA.
    const Trans *Qmap_inverse;
    const StateRel *follow;
    const StateSet *first;
    const StateSet *last;

    // Must always have an argument-less constructor.
    inline DSRFA::DSRFA():
        which(0),
        Qmap_inverse(0),
        follow(0),
        first(0),
        last(0),
        finalness(0) {
        // No assert since it won't satisfy the class invariant.
    }

    inline DSRFA::DSRFA( const DSRFA& r):
        which( r.which ),
        Qmap_inverse( r.Qmap_inverse ),
        follow( r.follow ),
        first( r.first ),
        last( r.last ),
        finalness( r.finalness ) {
        assert( class_invariantO );
    }

    inline const DSRFA& DSRFA::operator=( const DSRFA& r ) {
        assert( r.class_invariantO );
        // *this may not satisfy the class invariant yet.
        which = r.which;
        Qmap_inverse = r.Qmap_inverse;
        follow = r.follow;
        first = r.first;
        last = r.last;
        finalness = r.finalness;
        assert( class_invariantO );
        return( *this );
    }

    inline int DSRFA::final() const {
        assert( class_invariantO );
        return( finalness );
    }

    inline int DSRFA::operator==( const DSRFA& r ) const {
        assert( class_invariantO );
        assert( r.class_invariantO );
        assert( Qmap_inverse == r.Qmap_inverse );
        && follow == r.follow
        && first == r.first
Implementation: A DSRFA contains a StateSet, representing the set of RFA States that make up a DFA state in the subset construction (under the right-biased interpretation of states). It also contains an integer indicating if the abstract state is a final one.
Implementation class: **DSRFA2**

Files: dsrfa2.h, dsrfa2.cpp

Uses: CharRange, CRSet, StateRel, StateSet, Trans

Description: Class **DSRFA2** implements the abstract class interface required to construct a DFA from an RFA. **DSRFA2** is an alternative to **DSRFA**, using the left-biased interpretation of RFA States, instead of the right-biased interpretation used in **DSRFA**. A **DSRFA2** is constructed in the RFA member function **determinism2**, and then passed to template function **construct components**, which constructs the components of the DFA.
inline int class_invariant() const:

private:
  // Another special constructor, for use in out.transitions():
  DSRFA2( const StateSet& rq,
          const Trans *Qmap_inverse,
          const StateRel *follow,
          const StateSet *first,
          const StateSet *last );

  // This stuff may vary from object to object.
  StateSet which;
  int finalness;

  // This stuff should be the same for all objects corresponding
  // to a particular RFA.
  const Trans *Qmap_inverse;
  const StateRel *follow;
  const StateSet *first;
  const StateSet *last;
};

// Must always have an argument-less constructor.
inline DSRFA2::DSRFA2() :
  which(),
  Qmap_inverse( 0 ),
  follow( 0 ),
  first( 0 ),
  last( 0 ),
  finalness( 0 ) {
  // No assert since it won't satisfy the class invariant.
}

inline DSRFA2::DSRFA2( const DSRFA2& r ) :
  which( r.which ),
  Qmap_inverse( r.Qmap_inverse ),
  follow( r.follow ),
  first( r.first ),
  last( r.last ),
  finalness( r.finalness ) {
  assert( class_invariant() );
}

inline const DSRFA2& DSRFA2::operator=( const DSRFA2& r ) {
  assert( r.class_invariant() );
  // *this may not satisfy the class invariant yet.
  which = r.which;
  Qmap_inverse = r.Qmap_inverse;
  follow = r.follow;
  first = r.first;
  last = r.last;
  finalness = r.finalness;
  assert( class_invariant() );
  return( *this );
}

inline int DSRFA2::final() const {
  assert( class_invariant() );
  return( finalness );
}

inline int DSRFA2::operator==( const DSRFA2& r ) const {
  assert( class_invariant() );
  assert( r.class_invariant() );
  assert( Qmap_inverse == r.Qmap_inverse && follow == r.follow

&& first == r.first
&& last == r.last);
return( (which == r.which) && (finalness == r.finalness) );
}

inline int DSRFA2::operator!=( const DSRFA2& r ) const {
return( !operator==( r ) );
}

inline int DSRFA2::class_invariantO const {
return( Qmap_inverse != 0
&& follow != 0
&& first != 0
&& last != 0
&& which.domainO == Qmap_inverse->range()
&& which.domainO == follow->domain()
&& which.domainO == first->domain()
&& which.domainO == last->domain() );
}

#endif

Implementation: A DSRFA2 contains a StateSet, representing the set of RFA States that make up a DFA state in the subset construction (under the left-biased interpretation of states). It also contains an integer indicating if the abstract state is a final one.

Performance: For higher performance in constructing a DFA, use class DSRFA instead.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 13:59:56 $
#include "dsrfa2.h"

// Use the LBFA interpretation of State's (Section 4.2 of the Taxonomy).

// A special constructor (for use in constructing the start state):
DSRFA2::DSRFA2( const Trans *Qmap_inverse,
const StateRel *rfollow,
const StateSet *rfirst,
const StateSet *rlast,
const int rfinalness ) :
Qmap_inverse( rQmap_inverse ),
follow( rfollow ),
first( rfirst ),
last( rlast ),
finalness( rfinalness ) {
// We assume that we receive the finalness magically as a parameter.
which.set_domain( follow->domain() );
// The emptiness of which is the indicator of the start state.
assert( which.empty() );
assert( class_invariant() );
}

// Another special constructor, for use in out_transitions():
DSRFA2::DSRFA2( const StateSet& rq,
const Trans *Qmap_inverse,
const StateRel *rfollow,
const StateSet *rfirst,
const StateSet *rlast ) :
which( rq ),
Qmap_inverse( rQmap_inverse ),
follow( rfollow ),
first( rfirst ),
lst( rlast ) {
assert( !rq.empty() );
Reduced Finite Automata

13.1 The \( \Sigma \)-algebra \( \text{Reg}<\text{RFA}> \)

**Implementation class:** \( \text{Reg}<\text{RFA}> \)

**Files:** sig-rfa.cpp

**Uses:** CharRange, Reg, RFA

**Description:** The template instantiation \( \text{Reg}<\text{RFA}> \) implements the \( \Sigma \)-algebra of reduced finite automata, as defined in [Wat93a, Definition 4.29]. The operators can be used directly to construct complex RFAs, or the \( \text{Reg}<\text{RFA}> \) constructor from \( RE \) (which is defined in template class \( \text{Reg} \)) can be used to construct the homomorphic image of the regular expression.

**Implementation:** The implementation follows directly from [Wat93a, Definition 4.29].
13.1 The $\Sigma$-algebra Reg<RFA>

Nullable = 1;
assert( class_invariant() );
return( *this );
}

Reg<RFA>& Reg<RFA>::empty() {
    // See epsilon case.
    reincarnate();
    assert( Nullable == 0 );
    return( *this );
}

Reg<RFA>& Reg<RFA>::symbol( const CharRange r ) {
    // See epsilon case.
    reincarnate();
    auto State q( Q.allocate() );
    first.set_domain( Q.size() );
    first.add( q );
    last.set_domain( Q.size() );
    last.add( q );
    Qmap_inverse.set_range( Q.size() );
    Qmap_inverse.add_transition( r, q );
    follow.set_domain( Q.size() );
    // Nothing to add to follow.
    Nullable = 0;
    current.set_domain( Q.size() );
    assert( class_invariant() );
    return( *this );
}

Reg<RFA>& Reg<RFA>::or( const Reg<RFA>& r ) {
    assert( class_invariant() );
    assert( r.class_invariant() );
    // All state-related stuff in r must be adjusted.
    Q.incorporate( r.Q );
    first.disjointing_union( r.first );
    last.disjointing_union( r.last );
    Qmap_inverse.disjointing_union( r.Qmap_inverse );
    follow.disjointing_union( r.follow );
    Nullable = Nullable || r Nullable;
    current.set_domain( Q.size() );
    assert( class_invariant() );
    return( *this );
}

Reg<RFA>& Reg<RFA>::concat( const Reg<RFA>& r ) {
    assert( class_invariant() );
    assert( r.class_invariant() );
// See the or operator.
// All state-related stuff in r must be adjusted.
// First, incorporate the follow sets.

follow.disjointing_union( r.follow );
// Rename the incoming StateSet's first and last StateSets.
auto StateSet fil( r.first );
fil.st.rename( Q.size() );
auto StateSet lal( r.last );
lal.st.rename( Q.size() );

Q.incorporate( r.Q );
// Adjust last as well.
last.set_domain( Q.size() );
follow.union.cross( last, fil );

first.set_domain( Q.size() );
if( Nullable ) {
    first.set.union( fil );
}
if( r.Nullable ) {
    last.set.union( lal );
} else {
    last = lal;
}

Qmap_inverse.disjointing_union( r.Qmap_inverse );
Nullable = Nullable && r Nullable;

current.set_domain( Q.size() );
assert( class_invariant() );
return( *this );
}

Reg<RFA>& Reg<RFA>::star() {
assert( class_invariant() );

// Nothing to do to Q, first, last, Qmap_inverse.
follow.union.cross( last, first );
Nullable = 1;
assert( class_invariant() );
return( *this );
}

Reg<RFA>& Reg<RFA>::plus() {
assert( class_invariant() );

// Don't change Q, first, last, Qmap_inverse, Nullable.
follow.union.cross( last, first );
assert( class_invariant() );
return( *this );
}

Reg<RFA>& Reg<RFA>::question() {
assert( class_invariant() );

// Don't change Q, first, last, Qmap_inverse, follow.
Nullable = 1;
assert( class_invariant() );
return( *this );
}
14 Left-biased finite automata

Implementation class: LBFA

Files: lba.h, lba.cpp

Uses: DFA, FAabs, RFA, State, StatePool, StateRel, StateSet, Trans

Description: Class LBFA implements left-biased finite automata, as defined in [Wat93a, Definition 4.20]. It inherits from FAabs, and implements the interface defined by abstract base FAabs. A constructor taking an RFA provides a method of (indirectly) constructing an LBFA from regular expressions. The constructor implements isomorphism decode as defined in [Wat93a, Definition 4.28]. A special member function takes an an RFA and implements the non-isomorphic mapping convert defined in [Wat93a, Definition 4.35]. This mapping provides an alternative way of constructing an LBFA from an RFA (see [Wat93a, p. 36-37]).
friend ostream& operator<<( ostream& os, const LBFA& r );

inline int class_invariant() const;

protected:
    // Implementation stuff, protected for Reg<LBFA>:
    StatePool Q;
    // Single start state.
    State s;
    StateSet F;
    // Qmap (see Definition 4.24) stored as its inverse.
    Trans Qmap_inverse;
    StateRel follow;

    // Simulation stuff:
    StateSet current;
};

inline int LBFA::class_invariant() const {
    return ( Q.contains( s )
        && Q.size() == F.domain()
        && Q.size() == Qmap_inverse.range()
        && Q.size() == follow.domain()
        && Q.size() == current.domain()
        && F.class_invariant()
        && Qmap_inverse.class_invariant()
        && follow.class_invariant()
        && current.class_invariant() );
}

#endif

/* (c) Copyright 1994 by Bruce Watson */
// $Revision: 1.2 $  14/04/02
#include "lhfa.h"
#include "dfaseed.h"
#include "dslbfa.h"

LBFA::LBFA() {
    // Ensure that *this has the LBFA properties.
    // * a single start State, with no in-transitions,
    s = Q.allocate();
    F.set_domain( Q.size() );
    Qmap_inverse.set_range( Q.size() );
    follow.set_domain( Q.size() );
    current.set_domain( Q.size() );
    assert( class_invariant() );
}

// Implement homomorphism decode (Definition 4.28).
LBFA::LBFA( const RFA& r ) :
    Q( r.Q ),
    F( r.last ),
    Qmap_inverse( r.Qmap_inverse ),
    follow( r.follow ) {
    assert( r.class_invariant() );

    // The new start state:
    s = Q.allocate();
    F.set_domain( Q.size() );
Qmap_inverse.set_range( Q.size() );
follow.set_domain( Q.size() );
current.set_domain( Q.size() );

// ...which is final if the RFA is nullable.
if( r Nullable ) {
    F.add( s );
}
auto StateSet fst( r.first );
fst.set_domain( Q.size() );
follow.union_cross( s, fst );
assert( r.class_invariant() );
}

int LBFA::num_states() const {
    assert( class_invariant() );
    return( Q.size() );
}

void LBFA::restart() {
    assert( class_invariant() );
current.clear();
current.add( s );
    assert( class_invariant() );
}

void LBFA::advance( char a ) {
    assert( class_invariant() );

    // There are two possible ways to implement this. It's not clear which is
    // more efficient.
    // current = Qmap_inverse[a].intersection( follow.image( current ) );
    current = follow.image( current ).intersection( Qmap_inverse[a] );
    assert( class_invariant() );
}

int LBFA::in_final() const {
    return( current.not_disjoint( F ) );
}

int LBFA::stuck() {
    return( current.empty() );
}

DFA LBFA::determinism() const {
    // Make sure that this is structurally sound.
    assert( class_invariant() );

    // Need s as a singleton StateSet.
    auto StateSet S;
    S.set_domain( Q.size() );
    S.add( s );
    // Now construct the DFA components.
    return( construct_components( DSLBFA( S, &Qmap_inverse, &follow, &F ) ) );
}

// Implement homomorphism decode (Definition 4.28).
LBFA& LBFA::decode( const RFA& r ) {
    // Implement Definition 4.28 of the Taxonomy.
    assert( r.class_invariant() );
    Q = r.Q;
    F = r.last;
    Qmap_inverse = r.Qmap_inverse;
    follow = r.follow;

    // The new start state:
s = Q.allocate();
F.set_domain( Q.size() );
Qmap_inverse.set_range( Q.size() );
follow.set_domain( Q.size() );
current.set_domain( Q.size() );

// ...which is final if the RFA is nullable.
if( r.nullable ) {
    F.add( s );
}
auto StateSet fst( r.first );
fst.set_domain( Q.size() );
follow.union_cross( s, fst );
assert( class_invariant() );
return( *this );

// Implement non-homomorphic RFA->LBFA mapping convert (Defn. 4.35).
LBFA& LBFA::convert( const RFA& r ) {
    // Implement Definition 4.35 of the Taxonomy.
    assert( r.class_invariant() );

    Q = r.Q;
    // r.first must be a singleton set for this to work properly!
    // (See Property 4.37):
    assert( r.first.size() == 1 );
    s = r.first.smallest();

    F = r.last;
    Qmap_inverse = r.Qmap_inverse;
    follow = r.follow;

    current = r.current;

    assert( class_invariant() );
    return( *this );
}

ostream& operator<<( ostream& os, const LBFA& r ) {
    // r.class_invariant();
    os << "LBFAs
    os << "Q = " << r.Q << \n; os << "s = " << r.s << \n; os << "F = " << r.F << \n; os << "Qmap_inverse = " << r.Qmap_inverse << \n; os << "follow = " << r.follow << \n; os << "current = " << r.current << \n; return( os );

Implementation class: DSLBFA

Files: dslbfa.h, dslbfa.cpp

Uses: CharRange, CRSet, StateRel, StateSet, Trans

Description: Class DSLBFA implements the abstract class interface required to construct a DFA from an LBFA. A DSLBFA is constructed in the LBFA member function determinism, and then passed to template function construct_components, which constructs the components of the DFA.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
/\ $Date: 1994/08/15 14:00:33 $ \\
ifndef DLBLA_H \\
define DLBLA_H \\
#include "charrang.h" \\
#include "crset.h" \\
#include "state.h" \\
#include "stateset.h" \\
#include "staterel.h" \\
#include "transrel.h" \\
#include <assert.h> \\
#include <iostream.h> \\
/
This class is used to represent abstract States in a DFA that is still under 
construction. It is used in the construction of a DFA from an LBFA. The behaviour 
of a DSLBFA follows the simulation behaviour of an LBFA. 

class DLBLA { 
public: 
   // Must always have an argument-less constructor. 
inline DLBLA(); 
   
inline DLBLA( const DLBLA& r ); 
   // A special constructor: 
   DLBLA( const StateSet& rq, 
      const Trans *Qmap_inverse, 
      const StateRel *follow, 
      const StateSet *F ); 
   
inline const DLBLA& operator=( const DLBLA& r ); 
   // The required member functions: 
inline int final() const; 
CRSet out_labels() const; 
DLBLA out_transition( const CharRange a ) const; 
inline int operator==( const DLBLA& r ) const; 
inline int operator!=( const DLBLA& r ) const; 
   
friend ostream& operator<<( ostream& os, const DLBLA& r ); 
   
inline int class_invariant() const; 
private: 
   StateSet which; 
   const Trans *Qmap_inverse; 
   const StateRel *follow; 
   const StateSet *F; 
}; 
   // Must always have an argument-less constructor. 
inline DLBLA::DLBLA() : 
   which(), 
   Qmap_inverse( 0 ), 
   follow( 0 ), 
   F( 0 ) 
   // No assertion of the class invariant since it isn't satisfied yet. 
} 
inline DLBLA::DLBLA( const DLBLA& r ) : 
   which( r.which ), 
   Qmap_inverse( r.Qmap_inverse ), 
   follow( r.follow ), 
   F( r.F ) 
   assert( class_invariant() );
inline const DSLBFA& DSLBFA::operator=( const DSLBFA& r ) {
    assert( r.class_invariant() );
    // *this does not satisfy the class invariant yet.
    which = r.which;
    Qmap_inverse = r.Qmap_inverse;
    follow = r.follow;
    F = r.F;
    assert( class_invariant() );
    return( *this );
}

inline int DSLBFA::final() const {  
    assert( class_invariant() );
    return( which.not_disjoint( *F ) );
}

inline int DSLBFA::operator==( const DSLBFA& r ) const {
    assert( class_invariant() );
    assert( r.class_invariant() );
    assert( Qmap_inverse == r.Qmap_inverse && follow == r.follow && F == r.F );
    return( which == r.which );
}

inline int DSLBFA::operator!=( const DSLBFA& r ) const {
    return( operator==( r ) );
}

inline int DSLBFA::class_invariant() const {
    return( Qmap_inverse != 0 && follow != 0 && F != 0 &&
            which.domain() == Qmap_inverse->range() &&
            which.domain() == follow->domain() &&
            which.domain() == F->domain() );
}

#endif

Implementation: A DSLBFA contains a StateSet, representing the set of LBFA States that make up a DFA state in the subset construction. It also includes pointers to some components of the LBFA that constructed it.

#include "dslbfa.h"

// A special constructor:
DSLBFA::DSLBA( const StateSet& rq,
    const Trans *rQmap_inverse,
    const StateRel *rfollow,
    const StateSet *rF ) :
    which( rq ),
    Qmap_inverse( rQmap_inverse ),
    follow( rfollow ),
    F( rF ) {
    assert( class_invariant() );
}

CRSet DSLBFA::out_labels() const {
    assert( class_invariant() );
    // *Qmap_inverse contains the transitions into State's.
    // So, we take the image of which under follow to get which's out-transitions.
    return( Qmap_inverse->labels_into( follow->image( which ) ) );
}
// Pretty much follow the stuff in lbfa.cpp
DSLBFADSLBFADSLBFADSLBFA::out_transition(const CharRange a) const {
    assert(class_invariant());
    return(DSLBFADSLBFADSLBFA(follow->image(which).intersection(
         Qmap_inverse->range_transition(a)),
         Qmap_inverse,
         follow,
         F));
}

ostream& operator<<(ostream& os, const DSLBFADSLBFA& r) {
    os << "\nDSLBF\n" << r.which << r.Qmap_inverse << r.follow << r.F << "\n";
    return(os);
}
15 Right-biased finite automata

Implementation class: RBFA

Files: rbfa.h, rbfa.cpp

Uses: DFA, FAabs, RFA, State, StatePool, StateRel, Trans

Description: Class RBFA implements right-biased finite automata, as defined in [Wat93a, Construction 4.43]. It inherits from FAabs, and implements the interface defined by abstract base class FAabs. A constructor taking an RFA provides a method of (indirectly) constructing an RBFA from regular expressions. The constructor implements isomorphism \( R \circ \text{decode} \circ R \) as defined in [Wat93a, p. 41]. A special member function takes a reference to an RFA and implements the non-isomorphic mapping \( R \circ \text{convert} \circ R \) defined in [Wat93a, p. 43]. This mapping provides an alternative way of constructing an RBFA from an RFA (see [Wat93a, p. 42-43]).

```cpp
/* (c) Copyright 1994 by Bruce Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:45 $
#ifndef RBFA_H
#define RBFA_H

#include "faabs.h"
#include "st-pool.h"
#include "stateset.h"
#include "staterel.h"
#include "trans.h"
#include "rfa.h"
#include <assert.h>
#include <iostream.h>

class RBFA : virtual public FAabs {

public:
    // Constructors, destructors, operator=:
    // Default constructor builds the empty language acceptor.
    RBFA();
    // Default constr, copy constr, destr, operator= are okay.
    // Special constructor (from RFA) implementing \( R \circ \text{decode} \circ R \) (see
    // Construction 4.43 of the Taxonomy):
    RBFA( const RFA& r );
    // Standard FAabs member functions. Don't override acceptable();
    virtual int num_states() const;
    virtual void restart();
    virtual void advance( char a );
    virtual int in_final() const;
    virtual int stuck();
    virtual DFA determinism() const;
    // Some Sigma-algebra conversions:
    // as the duals of those given in Definitions 4.28 and 4.35
    virtual RBFA& decode( const RFA& r );
    virtual RBFA& convert( const RFA& r );

    // Special member functions:
```
friend ostream& operator<<( ostream& os, const RBFA& r );

inline int class_invariant() const;

protected:
  // Implementation stuff, protected for Reg<RBFA>:
  StatePool Q;
  StateSet S;
  // Single final state as required (Construction 4.43):
  State f;
  // See rfa.h and dfafa.h for explanation of Qmap_inverse:
  Trans Qmap_inverse;
  StateRel follow;

  // Simulation stuff:
  StateSet current;
};

inline int RBFA::class_invariant() const {
  return( Q.contains( f )
      && Q.size() == S.domain()
      && Q.size() == Qmap_inverse.range()
      && Q.size() == follow.domain()
      && Q.size() == current.domain()
      && S.class_invariant()
      && Qmap_inverse.class_invariant()
      && follow.class_invariant()
      && current.class_invariant() );
}

#endif

Implementation: Class RBFA is a friend of class RFA. Most of the member functions are implemented in a straight-forward manner.

RBFA::RBFA() {
  // Give *this the properties required of Construction 4.43
  f = Q.allocate();
  S.set_domain( Q.size() );
  Qmap_inverse.set_range( Q.size() );
  follow.set_domain( Q.size() );
  current.set_domain( Q.size() );

  assert( class_invariant() );
}

// Implement R, decode: R (Construction 4.45 of the Taxonomy):
RBFA::RBFA( const RFA& r ) :
  Q( r.Q ),
  S( r.first ),
  Qmap_inverse( r.Qmap_inverse ),
  follow( r.follow ) {
  assert( r.class_invariant() );

  // Create the new final state.
  f = Q.allocate();
RIGHT-BIADED FINITE AUTOMATA

```cpp
S.set_domain( Q.size() );
Qmap_inverse.set_range( Q.size() );
follow.set_domain( Q.size() );
current.set_domain( Q.size() );

if( r.Nullable ) {
    S.add( f );
}
auto StateSet la( r.last );
lq.set_domain( Q.size() );
follow.union_cross( la, f );
assert( class_invariant() );
}

int RBFA::num_states() const {
    assert( class_invariant() );
    return( Q.size() );
}

void RBFA::restart() {
    assert( class_invariant() );
current = S;
    assert( class_invariant() );
}

void RBFA::advance( char a ) {
    // Compute which states we can even transition out of, then
    // transition out of them.
    assert( class_invariant() );
current = follow.image( current.intersection( Qmap_inverse[a] ) );
    assert( class_invariant() );
}

int RBFA::in_final() const {
    assert( class_invariant() );
    return( current.contains( f ) );
}

int RBFA::stuck() {
    assert( class_invariant() );
    return( current.empty() );
}

DFA RBFA::determinism() const {
    // Make sure that *this is structurally sound.
    assert( class_invariant() );
    // Now construct the DFA components.
    return( construct_components( DSRBFA( S, &Qmap_inverse, &follow, f ) ) );
}

// implement homomorphism R; decode : R (Construction 4.45):
RBFA& RBFA::decode( const RFA& r ) {
    // Implement R; decode : R of the Taxonomy (p. 41).
    assert( r.class_invariant() );
    Q = r.Q;
    S = r.first;
    Qmap_inverse = r.Qmap_inverse;
    follow = r.follow;
    // Allocate the new final state.
    f = Q.allocate();
    S.set_domain( Q.size() );
    Qmap_inverse.set_range( Q.size() );
    follow.set_domain( Q.size() );
    current.set_domain( Q.size() );
```
current.clear();

if( r.Nullable ) {
    S.add( f );
}

auto StateSet la( r.last );
la.set_domain( Q.size() );
follow.union_cross( la, f );

assert( class_invariant() );
return( *this );

// Implement non-homomorphic RFA>RBFA mapping R; conver; R (see
// page 43 of the Taxonomy):
RBFA& RBFA::convert( const RFA& r ) {
    // Implement R; convert ; R of the Taxonomy (p. 43).
    Q = r.Q;
    // r.last must be a singleton for this to work.
    // (see Property 4.37 of the Taxonomy):
    f = r.last.smallest();
    assert( r.last.size() == 1 );

    S = r.first;
    Qmap_inverse = r.Qmap_inverse;
    follow = r.follow;

    current = r.current;

    assert( class_invariant() );
    return( *this );
}

ostream& operator< <( ostream& os, const RBFA& r ) {
    assert( r.class_invariant() );
    os << "\nRBFA\n";
    os << "Q = " << r.Q << "\n";
    os << "S = " << r.S << "\n";
    os << "f = " << r.f << "\n";
    os << "Qmap_inverse = " << r.Qmap_inverse << "\n";
    os << "follow = " << r.follow << "\n";
    os << "current = " << r.current << "\n";
    return( os );
}


Implementation class: DSRBFA

Files: dsrbfa.h, dsrbfa.cpp

Uses: CharRange, CRSet, StateRel, StateSet, Trans

Description: Class DSRBFA implements the abstract class interface required to construct a DFA
from an RBFA. A DSRBFA is constructed in the RBFA member function determinism, and
then passed to template function construct_components, which constructs the components of
the DFA.

/* (c) Copyright 1994 by Bruce W. Watson */

$Revision: 1.2 $
$Date: 1994/08/15 14:00:34 $
ifndef DSRBFA_H
#define DSRBFA_H
#define DSRBFA_H
#include "charrang.h"
#include "crset.h"
#include "state.h"
#include "stateset.h"
#include "staterel.h"
#include "transrel.h"
#include <assert.h>
#include <iostream.h>

// This class is used to represent abstract States in a DFA that is still under
// construction. It is used in the construction of a DFA from an RBFA. The behaviour
// of a DSRBFA follows the simulation behaviour of an RBFA.

class DSRBFA {
    public:
        // Must always have an argument-less constructor.
        inline DSRBFA();

        // A special constructor:
        inline DSRBFA( const DSRBFA& r );

        // The required member functions:
        inline int final() const;
        CRSet out_labels() const;
        DSRBFA out_transition( const CharRange a ) const;
        inline int operator==( const DSRBFA& r ) const;
        inline int operator!=( const DSRBFA& r ) const;

        friend ostream& operator<<( ostream& os, const DSRBFA& r );

        inline int class_invariant() const;

    private:
        StateSet which;
        const Trans *Qmap_inverse;
        const StateRel *follow;
        State f;
};

// Must always have an argument-less constructor.
inline DSRBFA::DSRBFA():
    which(),
    Qmap_inverse( 0 ),
    follow( 0 ),
    { ( Invalid ) } {
        // this will not yet satisfy the class invariant.
    }

inline DSRBFA::DSRBFA( const DSRBFA& r ) :
    which( r.which ),
    Qmap_inverse( r.Qmap_inverse ),
    follow( r.follow ),
    { ( f.f ) } {
        assert( class_invariant() );
    }

inline const DSRBFA& DSRBFA::operator=( const DSRBFA& r ) { 
    assert( r.class_invariant() );
}
// *this may not satisfy the invariant yet.
which = r.which;
Qmap_inverse = r.Qmap_inverse;
follow = r.follow;
f = r.f;
assert( class_invariant() );
return( *this );
}

inline int DSRBFA::final() const {
assert( class_invariant() );
return( which.contains( f ) );
}

inline int DSRBFA::operator==( const DSRBFA& r ) const {
assert( class_invariant() );
assert( r.class_invariant() );
assert( Qmap_inverse == r.Qmap_inverse && follow == r.follow && f == r.f );
return( which == r.which );
}

inline int DSRBFA::operator!=( const DSRBFA& r ) const {
return( !operator==( r ) );
}

inline int DSRBFA::class_invariant() const {
return( Qmap_inverse != 0 && follow != 0 && 0 <= I && f < follow->domain() && which.domain() == Qmap_inverse->range() && which.domain() == follow->domain() );
}

#endif

Implementation: A DSRBFA contains a StateSet, representing the set of RBFA States that make up a DFA state in the subset construction. It also includes pointers to some components of the RBFA which constructed it.

/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 13:59:53 $
#include "dsrbfa.h"

// A special constructor:
DSRBFA::DSRBFA( const StateSet& rq,
    const Trans *Qmap_inverse,
    const StateRel *follow,
    const State rf ) :
    which( rq ),
    Qmap_inverse( rQmap_inverse ),
    follow( rfollow ),
    f( rf ) {
assert( class_invariant() );
}

CRSet DSRBFA::out_labels() const {
assert( class_invariant() );
return( Qmap_inverse->labels_into( which ) );
}

// Pretty much follow the stuff in rbfa.cpp
DSRBFA DSRBFA::out_transition( const CharRange a ) const {
assert( class_invariant() );
return( DSRBFA( follow->image( Qmap_inverse->range_transition( a ) )


intersection( which ),
Qmap_inverse,
follow,
f );

ostream& operator<<( ostream& os, const DSRBFA& r ) {
    os << "DSRBFA\n" << r.which << *r.Qmap_inverse << *r.follow << r.f << '\n';
    return( os );
}
16 Automata constructions

The finite automata construction functions are functions which take an RE (or a pointer to an RE) and return an automaton of some type (one of FA, RFA, LBFA, RBFA, or DFA). Functions implementing Brzozowski's construction [Wat93a, Section 5.3] are given in dconstrs.cpp and dconstrs.h. Functions implementing the item set constructions [Wat93a, Section 5.5] are given in iconstrs.cpp and iconstrs.h. Files constrs.h and constrs.cpp contain all of the remaining functions (from [Wat93a, Section 4]).

All of the function declarations (or prototypes) are accompanied by their corresponding reference in [Wat93a].

```c
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $  
// $Date: 1994/08/15 14:00:24 $ 
#ifndef CONSTRS_H  
#define CONSTRS_H

#include "re.h"
#include "fa.h"
#include "rfa.h"
#include "lbfa.h"
#include "rbfa.h"
#include "sigma.h"
#include "dfa.h"
#include <assert.h>

 inline FA Thompson( const RE& r ) {  
  return( FA( r ) );
}

 inline RFA rfa( const RE& r ) {  
  return( RFA( r ) );
}

 inline DFA MYG_shortcut( const RE& r ) {  
  return( RFA( r ).determinism2() );
}
```
// Construction 4.50 (variation of Aho-Sethi-Ullman)
inline DFA ASU_shortcut( const RE& r ) {
    return( RFA( r ).determinism() );
}

// LBFA constructions:
// using RFA's.

// Construction 4.32
inline LBFA BerrySethi( const RE& r ) {
    return( RFA( r ) );
}

// Construction 4.38
LBFA BS_variation( const RE& r );

// DFA from LBFA construction:
// Construction 4.39
inline DFA MYG( const RE& r ) {
    return( BerrySethi( r ).determinism() );
}

// RBFA constructions:
// using RFA's.

// Construction 4.45
inline RBFA BerrySethi_dual( const RE& r ) {
    return( RFA( r ) );
}

// Construction 4.48
RBFA BS_variation_dual( const RE& r );

// Construction 4.50
inline DFA ASU( const RE& r ) {
    return( BS_variation_dual( r ).determinism() );
}
#endif

/
* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/18 13:59:42 $
#include "constrs.h"

LBFA BS_variation( const RE& r ) {
    assert( r.class_invariant() );

    // Construct the homomorphic image of r.
    auto RFA l( r );

    // Construct the begin-marker.
    auto Reg<RFA> m;
    m.symbol( '$' );

    // Attach the begin-marker.
    m.concat( (const Reg<RFA>&)l );

    // Construct the dummy LBFA, for use with the non-homomorphic mapping
    // convert (Defn. 4.35 of the Taxonomy)
    auto LBFA LB;

    return( LB.convert( m ) );
}
RBFA BS_variation_dual( const RE& r ) {
    // See above.
    assert( r.class_invariant() );

    auto RFA ( r );
    // Construct the end-marker.
    auto Reg<RFA> m;
    m.symbol( '$' );

    // Attach the end-marker.
    ((Reg<RFA>&)I).concat( m );

    auto RBFA RB;
    return( RB.convert( 1 ) );
}

DFA Brz( const RE& r ) {
    assert( r.class_invariant() );
    return( construct_components( DSRE( r ) ) );
}

DFA Brz_optimized( const RE& r ) {
    assert( r.class_invariant() );
    return( construct_components( DSREopt( r ) ) );
}

#ifndef DCONSTRS_H
#define DCONSTRS_H

#include "re.h"
#include "dfa.h"
#include <assert.h>

// The numbers in parentheses refer to the Construction number in the Taxonomy.

// DFA constructions:
// Brzowski's constructions:
// Normal (Construction 5.34)
DFA Brz( const RE& r );

// With strong similarity (Construction 5.34 + Remark 5.32)
DFA Brz_optimized( const RE& r );
#endif
#ifndef ICONSTRS_H
#define ICONSTRS_H

#include "re.h"
#include "dfa.h"
#include <assert.h>

// The numbers in parentheses refer to the Construction number in the Taxonomy.

// Item set constructions:
// Iconstr (Construction 5.69)
DFA Iconstr( const RE *r );

// DeRemer's (Construction 5.75)
DFA DeRemer( const RE *r );

// Oconstr (Construction 5.82)
DFA Oconstr( const RE *r );

#endif

/* (c) Copyright 1994 by Bruce W. Watson */
/* $Revision: 1.2 $ */
/* $Date: 1994/08/15 14:00:43 $ */
#include "dconstrs.h"
#include "dsis.h"
#include "dsderem.h"
#include "dsis-opt.h"
#include "dfaseed.h"

DFA Iconstr( const RE *r ) {
    assert( r->class_invariant() );
    return( construct_components( DSIS( r ) ) );
}

DFA DeRemer( const RE *r ) {
    assert( r->class_invariant() );
    return( construct_components( DSDeRemer( r ) ) );
}

DFA Oconstr( const RE *r ) {
    assert( r->class_invariant() );
    return( construct_components( DSIS_opt( r ) ) );
}
Part IV
Minimizing DFAs

An interesting DFA minimization function is the one due to Brzozowski (appearing in [Wat93b, Section 2]).

User function: DFA::min_Brzozowski

Files: dfa.h

Uses: see the entry for class DFA

Description: Brzozowski's minimization algorithm differs from all of the other minimization algorithms, and appears as inline DFA member function min_Brzozowski. The member function simply minimizes the DFA.

Implementation: The implementation appears in file dfa.h. The implementation follows directly from [Wat93b, Section 2].

The remainder of Part IV deals with the other minimization functions, based upon their derivations in [Wat93b].
17 Helper functions

Two member functions of class \textit{DFA} are used as helpers in those minimization functions which are based upon an equivalence relation (see [Wat93b, Section 3]).

\section*{Implementation function: \textit{DFA::compress}}

\textbf{Files:} dfa.cpp  
\textbf{Uses:} see the entry for class \textit{DFA}

\textbf{Description:} Member function \textit{compress} is overloaded to take either a \textit{StateEqRel} (an equivalence relation on \textit{States}), or a \textit{SymRel} (a symmetrical relation on \textit{States}). The relation is assumed to be a refinement of relation \textit{E} [Wat93b, Definition 3.3]. The member function implements function \textit{merge} defined in [Wat93b, Transformation 3.1].

\textbf{Implementation:} The implementation appears in file dfa.cpp. The implementation follows directly from [Wat93b, Transformation 3.1].

\section*{Implementation function: \textit{DFA::split}}

\textbf{Files:} min.cpp  
\textbf{Uses:} \textit{CharRange}, \textit{DFA}, \textit{State}, \textit{StateEqRel}, \textit{StateSet}

\textbf{Description:} Member function \textit{split} takes two \textit{States} (\textit{p} and \textit{q}), a \textit{CharRange} \textit{a}, and a \textit{StateEqRel}. The two \textit{States} are assumed to be representatives in the \textit{StateEqRel}. The equivalence class of \textit{p} is then split into those \textit{States} (equivalent to \textit{p}) which transition on \textit{a} to a \textit{State} equivalent to \textit{q}, and those that do not. If in fact there was a split (the equivalence class of \textit{p} is split into two equivalence classes), \textit{p} will still be the unique representative of one of the two new equivalence classes; in this case, the (unique) representative of the other equivalence class is returned. If no such split occurred, function \textit{split} returns \textit{Invalid}.

\textbf{Implementation:} The implementation is a simple, iterative one. It implements the equivalent of the assignment to \textit{Q'} in Algorithm 4.4 of [Wat93b].

```cpp
State DFA::split( const State p, const State q, const CharRange a, StateEqRel& P ) const {
    assert( class_invariant() );
    assert( P.class_invariant() );
    if ( p and q must be representatives of their eq. classes. 
    assert( p == P.eq.class_representative( p ) );
    assert( q == P.eq.class_representative( q ) );
    // Split [p] with respect to [q] and CharRange a.
    auto StateSet part;
    part.set_domain( Q.size() );
    // Iterate over [p], and see whether each member transitions into [q]
    // on CharRange a.
    auto State st;
    for( P.eqv.class( p ), iter_start( st );
```
P.equiv_class( p ).iter_end( st );
P.equiv_class( p ).iter_next( st ) {  
    auto State dest( T.transition_on_range( st, a ) );
    // It could be that dest == Invalid.
    // if not, check if dest is in [q].
    if ( dest != Invalid && P.equivalent( dest, q ) ) {
        part.add( st );
    }
}

// The following containment must hold after the splitting.
assert( P.equiv_class( p ).contains( part ) );

// Return non-zero if something was split.
if( (part != P.equiv_class( p )) && !part.empty() ) {
    auto StateSet otherpiece( P.equiv_class( p ) );
    otherpiece.remove( part ); assert( !otherpiece.empty() );

    P.split( part );
    assert( class_invariant() );

    // Now, we must return the representative of the newly created
    // equivalence class.
    auto State x( P.eq_class_representative( otherpiece.smallest() ) );
    assert( x != Invalid );
    // It could be that p is not in part.
    auto State y( P.eq_class_representative( part.smallest() ) );
    assert( y != Invalid );
    assert( x != y );

    if( p == x ) {
        // If p is the representative of 'otherpiece', then
        // return the representative of 'part'.
        assert( otherpiece.contains( p ) );
        return( y );
    } else {
        // If p is the representative of 'part', then return the
        // representative of 'otherpiece'.
        assert( part.contains( p ) );
        return( x );
    }
} else {
    assert( (part == P.equiv_class( p )) || (part.empty()) );
    // No splitting to be done.
    return( Invalid );
}
18 Aho, Sethi, and Ullman's algorithm

User function: `DFA::min_dragon`

Files: `min-asu.cpp`

Uses: `CRSet`, `DFA`, `State`, `StateEqRel`, `StateSet`

Description: Member function `min_dragon` is an implementation of Algorithm 4.6 in [Wat93b]. It is named the "dragon" function after Aho, Sethi, and Ullman’s "dragon book" [ASU86].

Implementation: The algorithm begins with the approximation $E_0$ to the State equivalence relation $E$. It then repeatedly partitions the equivalence classes until the greatest fixed point ($E$) is reached.

```cpp
#include "crset.h"
#include "state.h"
#include "state_eqrel.h"
#include "dfa.h"

// Implement Algorithm 4.6 of the Taxonomy.
DFA& DFA::min_dragon() {
    assert( class_invariant() );
    assert( Useful() );

    // We'll need the combination of all of the out-transitions of all of the
    // States, when splitting equivalence classes.
    auto CRSet C;
    auto State st;
    for( st = 0; st < Q.size(); st++ ) {
        Ccombine( T.out_labels( st ) );
    }

    // $P$ is the equivalence relation approximation to $E$. It is initialized
    // to the total relation with domain $Q.size()$.
    auto StateEqRel P( Q.size() );
    // We now initialize it to $E_0$.
    P.split( F );

    // reps is the set of representatives of $P$.
    auto StateSet reps( P.representatives() );

    // [st] is going to be split w.r.t. something.
    // The following is slightly convoluted.
    reps.iter_start( st );
    while( !reps.iter_end( st ) ) {
        // Try to split [st] w.r.t. every class [q].
        auto State q;

        // Keep track of whether something was indeed split.
        // Having to use this variable could be avoided with a goto
        auto int something_split( 0 );

        // Iterate over all q, and try to split [st] w.r.t. all other
        // equivalence classes [q].
        // Stop as early as possible.
        for( reps.iter_start( q ); reps.iter_end( q )
            && !something_split; reps.iter_next( q ) ) {
            // Iterate over the possible transition labels, and
            // do a split if possible.
            auto int i;
```
for( i = 0; !C.iter_end( i ) && !something_split; i++ ) {
    something_split = (split( st, q, C.iterator( i ), P ) != Invalid);
}

// If something was split, restart the outer repetition.
if( something_split ) {
    // The set of representatives will have changed due to
    // the split.
    reps = P.representatives();
    reps.iter_start( st );
    // Now continue the outer repetition with the restarted
    // representatives.
} else {
    // Just go on as usual.
    reps.iter_next( st );
}

compress( P );
assert( class_invariant() );

return( *this );
19 Hopcroft and Ullman’s algorithm

User function: DFA::min_HopcroftUllman

Files: min-hu.cpp

Uses: CRSet, DFA, State, StateSet, SymRel

Description: This member function provides a very convoluted implementation of Algorithm 4.7 of [Wat93b]. It computes the distinguishability relation (relation $D$ in [Wat93b]).

Implementation: The algorithm is not quite the same as that presented in [Wat93b, Algorithm 4.7]. The member function computes the equivalence relation (on States) by first computing distinguishability relation $D$. Initially, the pairs of distinguishable States are those pairs $p, q$ where one is a final State and the other is not. Pairs of States that reach $p, q$ are then also distinguishable (according to the efficiency improvement property given in [Wat93b, Section 4.7, p. 16]). Iteration terminates when all distinguishable States have been considered.

Performance: This algorithm can be expected to run quite slowly, since it makes use of reverse transitions in the deterministic transition relation (DTransRel). The transition relations (TransRel and DTransRel) are optimized for forward transitions.

```cpp
// Implement (a version of) Algorithm 4.7 of the minimization Taxonomy.
DFA& DFA::min_HopcroftUllman() {
    assert( class_invariant() );
    assert( UsefulU() );

    // We need the combination of all transition labels, to iterate over transitions.
    auto State st;
    auto CRSet C;
    for( st = 0; st < Q.size(); st++ ) {
        C.combine( T.out_labels( st ) );
    }

    // First, figure out which are the non-final States.
    auto StateSet nonfinal( F );
    nonfinal.complement();

    // Use Z to keep track of the pairs that still need to be considered for distinguishedness.
    auto SymRel Z;
    Z.set_domain( Q.size() );

    // We begin with those pairs that are definitely distinguished.
    // this includes pairs with different parity.
    Z.add_pairs( F, nonfinal );

    // It also includes pairs with differing out-transitions.
    // iterate over C (the CRSet) and check the have-haves and have-nots
    auto StateSet have;
    have.set_domain( Q.size() );
    auto StateSet havenot;
    havenot.set_domain( Q.size() );

    auto int it;
```
for( it = 0; !C.iter_end( it ); it++ ) {
  // have and havenot must be empty for the update to be correct.
  assert( have.empty() && havenot.empty() );
  auto State p;
  // Iterate over the States and check which have a transition.
  for( p = 0; p < Q.size(); p++ ) {
    // Does p have the transition?
    if( T.transition_on_range( T.iterator( it ), it ) != Invalid ) {
      have.add( p );
    } else {
      havenot.add( p );
    }
  }
  // have and havenot are distinguished from one another.
  // (under the assumption that Useful() holds)
  Z.add_pairs( have, havenot );
  have.clear();
  havenot.clear();
}

// G will be use to accumulate the (distinguishability) relation D.
auto SymRel G;
G.set_domain( Q.size() );

// Now consider all of the pairs until nothing changes.
while( 1 ) {
  auto State p;
  // Go looking for the next pair to do.
  for( p = 0; p < Q.size() && Z.image( p ).empty(); p++ );
  // There may be nothing left to do.
  if( p == Q.size() ) {
    break;
  } else {
    assert( !Z.image( p ).empty() );
    // Choose q such that {p,q} is in Z.
    // We know that such a q exists.
    auto State q( Z.image( p ).smallest() );
    assert( q != Invalid );
    // Move {p,q} from Z across to G.
    Z.remove_pair( p, q );
    G.add_pair( p, q );
    // Now, check out the reverse transitions from p and q.
    auto int i;
    // Iterate over all of the labels.
    for( i = 0; !C.iter_end( i ); i++ ) {
      auto StateSet pprime( T.reverse_transition( p, C.iterator( i ) ) );
      auto StateSet qprime( T.reverse_transition( q, C.iterator( i ) ) );
      // pprime and qprime are all distinguished.
      // Iterate over both sets and flag them as distinguished.
      auto State pp;
      for( pprime.iter_start( pp ); !ppprime.iter_end( pp );
          pprime.iter_next( pp ) ) {
        auto State qp;
        for( qprime.iter_start( qp ); !qprime.iter_end( qp );
            qprime.iter_next( qp ) ) {
          // Mark pp, qp for processing if they
          // haven't already been done.
          if( !G.contains_pair( pp, qp ) ) {
            Z.add_pair( pp, qp );
          }
        }
      }
    }
  }
}
Now, compute $E$ from $D$

```cpp
G.complement();
```

And use it to compress the DFA.

```cpp
compress( G );
```

assert( class_invariant() );

return( *this );
20 Hopcroft's algorithm

User function: DFA::min_Hopcroft

Files: min-hop.cpp

Uses: CRSet, DFA, State, StateEqRel, StateSet

Description: Member function min_Hopcroft implements Hopcroft's \(n \log n\) minimization algorithm, as presented in [Wat93b, Algorithm 4.8].

Implementation: The member function uses some encoding tricks to effectively implement the abstract algorithm. The combination of the out-transitions of all of the States is stored in a CRSet \(C\). Set \(L\) from the abstract algorithm is implemented as a mapping from States to int (an array of int is used). Array \(L\) should be interpreted as follows: if State \(q\) a representative, then the following pairs still require processing (are still in abstract set \(L\)):

\[
([q], C_0), \ldots, ([q], C_{L(q)-1})
\]

The remaining pairs do not require processing:

\[
([q], C_{L(q)}), \ldots, ([q], C_{|C|-1})
\]

This implementation facilitates quick scanning of \(L\) for the next valid State-CharRange pair.

```c
/* (c) Copyright 1994 by Bruce W. Watson */
// $Revision: 1.2 $
// $Date: 1994/08/15 14:00:05 $
#include "crset.h"
#include "state.h"
#include "stateset.h"
#include "st-eqrel.h"
#include "dfa.h"

// Implement Algorithm 4.8 (Hopcroft's \(O(n \log n)\) algorithm).
DFA& DFA::min_Hopcroft() {
    assert( class_invariant() );

    // This algorithm requires that the DFA not have any final unreachable
    // States.
    assert( Useful() );

    auto State q;

    // Keep track of the combination of all of the out labels of State's.
    auto CRSet C;
    for( q = 0; q < Q.size(); q++ ) {
        C.combine( T.out_labels( q ) );
    }

    // Encode set \(L\) as a mapping from State to \([0,|C|]\) where:
    // if \(q\) is a representative of a class in the partition \(P\), then
    // \(L\) (the abstract list) contains
    // \([q], C_0), ([q], C_1), \ldots, ([q], C_{L(q)-1})
    // but not
    // ([q], C_{L(q)}), \ldots, ([q], C_{|C|-1})
    auto int *const L( new int [Q.size()] );
    for( q = 0; q < Q.size(); q++ ) {
        L[q] = 0;
    }

    // Initialize \(P\) to be the total equivalence relation.
    StateEqRel P( Q.size() );
    // Now set \(P\) to be \(E\).
```
P.split( F );

// Now, build the set of representatives and initialize L.
auto StateSet repr( P.representatives() );
if( F.size() <= (Q.size() - F.size()) ) {
    repr.intersection( F );
} else {
    repr.remove( F );
}

// Do the final setup of L
for( repr.iter_start( q ); !repr.iter_end( q ); repr.iter_next( q ) ) {
    L[q] = C.size();
}

// Use a break to get of this loop.
while( 1 ) {
    // Find the first pair in L that still needs processing.
    for( q = 0; q < Q.size() && !L[q]; q++ );

    // It may be that we're at the end of the processing.
    if( q == Q.size() ) {
        break;
    } else {
        // Mark this element of L as processed.
        L[q]--; } // Iterate over all eq. classes, and try to split them.
auto State p;
    repr = P.representatives();
    for( repr.iter_start( p ); !repr.iter_end( p );
        repr.iter_next( p ) ) {
        // Now split [p] w.r.t. (q,C(L[q]))
        auto State r( split( p, q, C.iterator( L[q] ), P ) );
        // r is the representative of the new split of the
        // eq. class that was represented by p.
        if( r != Invalid ) {
            // p and r are the new representatives.
            // Now update L with the smallest of
            // [p] and [r].
            if( P.equiv_class( p ).size() <= P.equiv_class( r ).size() ) {
                L[r] = L[p];
                L[p] = C.size();
            } else {
                L[r] = C.size();
        } // if
        } // /\ for
    } // while

// L is no longer needed.
delete L;

// We can now use P to compress the DFA.
compress( P );
assert( class_invariant() );
return *this ;
}
21 A new minimization algorithm

User function: DFA::min_Watson

Files: min-bww.cpp

Uses: CRSet, DFA, State, StateEqRel, SymRel

Description: Member function min_Watson implements the new minimization algorithm appearing in [Wat93b, Sections 4.6-4.7]. The algorithm computes the equivalence relation \( E \) (on states) from below. The importance of this is explained in [Wat93b, p. 16].

Implementation: Function min_Watson is an implementation of Algorithm 4.10 (of [Wat93b]). The helper function are_eq is an implementation of the second algorithm appearing in Section 4.6 of [Wat93b]. Function are_eq takes two parameters more than the version appearing in [Wat93b]: \( H \) (a StateEqRel) and \( Z \) (a SymRel). These parameters are used to discover equivalence (or distinguishability) of States earlier than the abstract algorithm would.

Performance: Memoization in function are_eq would provide a great speedup.

```c
/* (c) Copyright 1994 by Bruce W. Watson */
/ * Revision: 1.2 */
/ * $Date: 1994/08/15 14:00:04 $ */
#include "state.h"
#include "erset.h"
#include "symrel.h"
#include "st-eqrel.h"
#include "dfa.h"
#include <assert.h>

// The following is function equiv from Section 4.6 of the min. Taxonomy.
int DFA::are_eq(State p, State q, SymRel& S, const StateEqRel& H, const SymRel& Z) const {
    if (S.contains_pair(p, q) || H.equivalent(p, q)) {
        // p and q are definitely equivalent.
        return(1);
    } else if (!Z.contains_pair(p, q)) {
        assert(!S.contains_pair(p, q));
        assert(!H.equivalent(p, q));
        // p and q were already compared (in the past) and found to be
        // inequivalent.
        return(0);
    } else {
        assert(!H.equivalent(p, q));
        assert(Z.contains_pair(p, q));

        // Now figure out all of the valid out-transitions from p and q.
        auto CRSet C(T.out_labels(p));
        C.combine(T.out_labels(q));

        // Just make sure that p and q have the same out-transitions.
        auto int it;
        for (it = 0; !C.iter_end(it); it++) {
            if (!T.transition_on_range(p, C.iterator(it)) || T.transition_on_range(q, C.iterator(it)) == Invalid)
                // There's something that one transitions on, but
                // the other doesn't
                return(0);
        }
    } // if
}
```

```
// p and q have out-transitions on equivalent labels.
// Keep track of whether checking needs to continue.
S.add_pair(p, q);
for (it = 0; !C.iter_end(it); it++) {
    // if
```

auto State pdest( T.transition_on_range( p, C.iterator( it ) ) );
auto State qdest( T.transition_on_range( q, C.iterator( it ) ) );
if( !are_eq( pdest, qdest, S, H, Z ) ) {
    // p and q have been found distinguished.
    S.remove_pair( p, q );
    return( 0 );
}
} // for

// p and q have been found equivalent.
S.remove_pair( p, q );
return( 1 );
} // if

DFA& DFA::min_Watson() {
    assert( class_invariant() );
    assert( Useful() );

    // (Symmetrical) State relation S is from p.14 of the min. Taxonomy.
    auto SymRel S;
    S.set_domain( Q.size() );

    // H is used to accumulate equivalence relation E.
    auto StateEqRel H( Q.size() );
    // Start with the identity since this is approximation from below
    // w.r.t. refinement.
    H.identity();

    // Z is a SymRel containing pairs of States still to be considered.
    auto SymRel Z;
    Z.set_domain( Q.size() );
    Z.identity();
    // We will need the set of non-final States to initialize Z.
    auto StateSet nonfinal( F );
    nonfinal.complement();
    Z.add_pairs( F, nonfinal );
    // Z now contains those pairs that definitely do not need comparison.
    Z.complement();
    // Z initialized properly now.

    auto State p;
    // Consider each p.
    for( p = 0; p < Q.size(); p++ ) {
        auto State q;
        // Consider each q that p still needs to be compared to.
        for( Z.image( p ).iter_start( q ); Z.image( p ).iter_end( q );
             Z.image( p ).iter_next( q ) ) {
            // Now compare p and q.
            if( are_eq( p, q, S, H, Z ) ) {
                // p and q are equivalent.
                H.equivalize( p, q );
            } else {
                // Don't do anything since we aren't computing
                // distinguishability explicitly.
            }
        } // for
    } // for

    compress( H );
assert( class_invariant() );
return( *this );
}
References


### Computing Science Notes

**In this series appeared:**

<table>
<thead>
<tr>
<th>Volume</th>
<th>Authors</th>
<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>91/01</td>
<td>D. Alstein</td>
<td>Dynamic Reconfiguration in Distributed Hard Real-Time Systems, p. 14</td>
</tr>
<tr>
<td>91/02</td>
<td>R.P. Nederpelt, H.C.M. de Swart</td>
<td>Implication. A survey of the different logical analyses &quot;if...,then...&quot;, p. 26</td>
</tr>
<tr>
<td>91/03</td>
<td>J.P. Katoen, L.A.M. Schoenmakers</td>
<td>Parallel Programs for the Recognition of $P$-invariant Segments, p. 16</td>
</tr>
<tr>
<td>91/05</td>
<td>D. de Reus</td>
<td>An Implementation Model for GOOD, p. 18</td>
</tr>
<tr>
<td>91/06</td>
<td>K.M. van Hee</td>
<td>SPECIFICATIEMETHODEN, een overzicht, p. 20</td>
</tr>
<tr>
<td>91/07</td>
<td>E. Poll</td>
<td>CPO-models for second order lambda calculus with recursive types and subtyping, p. 49</td>
</tr>
<tr>
<td>91/08</td>
<td>H. Schepers</td>
<td>Terminology and Paradigms for Fault Tolerance, p. 25</td>
</tr>
<tr>
<td>91/09</td>
<td>W.M.P.v.d.Aalst</td>
<td>Interval Timed Petri Nets and their analysis, p.53</td>
</tr>
<tr>
<td>91/11</td>
<td>R.C. Backhouse, P.J. de Bruin, G.Malcolm, E.Voermans, J. vander Woude</td>
<td>Relational Catamorphism, p. 31</td>
</tr>
<tr>
<td>91/12</td>
<td>E. van der Sluis</td>
<td>A parallel local search algorithm for the travelling salesman problem, p. 12</td>
</tr>
<tr>
<td>91/13</td>
<td>F. Rietman</td>
<td>A note on Extensionality, p. 21</td>
</tr>
<tr>
<td>91/14</td>
<td>P. Lemmens</td>
<td>The PDB Hypermedia Package. Why and how it was built, p. 63</td>
</tr>
<tr>
<td>91/16</td>
<td>A.J.J.M. Marcelis</td>
<td>An example of proving attribute grammars correct; the representation of arithmetical expressions by DAGs, p. 25</td>
</tr>
<tr>
<td>Year</td>
<td>Authors</td>
<td>Title</td>
</tr>
<tr>
<td>------</td>
<td>---------</td>
<td>-------</td>
</tr>
<tr>
<td>91/18</td>
<td>Rik van Geldrop</td>
<td>Transformational Query Solving, p. 35.</td>
</tr>
<tr>
<td>91/19</td>
<td>Erik Poll</td>
<td>Some categorical properties for a model for second order lambda calculus with subtyping, p. 21.</td>
</tr>
<tr>
<td>91/23</td>
<td>K.M. van Hee, L.J. Somers, M. Voorhoeve</td>
<td>Z and high level Petri nets, p. 16.</td>
</tr>
<tr>
<td>91/24</td>
<td>A.T.M. Aerts, D. de Reus</td>
<td>Formal semantics for BRM with examples, p. 25.</td>
</tr>
<tr>
<td>91/25</td>
<td>P. Zhou, J. Hooman, R. Kuiper</td>
<td>A compositional proof system for real-time systems based on explicit clock temporal logic: soundness and completeness, p. 52.</td>
</tr>
<tr>
<td>91/27</td>
<td>F. de Boer, C. Palamidessi</td>
<td>Embedding as a tool for language comparison: On the CSP hierarchy, p. 17.</td>
</tr>
<tr>
<td>91/28</td>
<td>F. de Boer</td>
<td>A compositional proof system for dynamic process creation, p. 24.</td>
</tr>
<tr>
<td>91/30</td>
<td>J.C.M. Baeten, F.W. Vaandrager</td>
<td>An Algebra for Process Creation, p. 29.</td>
</tr>
<tr>
<td>91/31</td>
<td>H. ten Eikelder</td>
<td>Some algorithms to decide the equivalence of recursive types, p. 26.</td>
</tr>
<tr>
<td>91/33</td>
<td>W. v.d. Aalst</td>
<td>The modelling and analysis of queueing systems with QNM-ExSpect, p. 23.</td>
</tr>
<tr>
<td>91/34</td>
<td>J. Coenen</td>
<td>Specifying fault tolerant programs in deontic logic, p. 15.</td>
</tr>
<tr>
<td>Year</td>
<td>Authors</td>
<td>Title</td>
</tr>
<tr>
<td>------</td>
<td>---------</td>
<td>-------</td>
</tr>
<tr>
<td>91/35</td>
<td>F.S. de Boer J.W. Klop C. Palamidessi</td>
<td>Asynchronous communication in process algebra</td>
</tr>
<tr>
<td>92/01</td>
<td>J. Coenen J. Zwiers W.-P. de Roever</td>
<td>A note on compositional refinement</td>
</tr>
<tr>
<td>92/02</td>
<td>J. Coenen J. Hooman</td>
<td>A compositional semantics for fault tolerant real-time systems</td>
</tr>
<tr>
<td>92/03</td>
<td>J.C.M. Baeten J.A. Bergstra</td>
<td>Real space process algebra</td>
</tr>
<tr>
<td>92/04</td>
<td>J.P.H.W.v.d.Eijnde</td>
<td>Program derivation in acyclic graphs and related problems</td>
</tr>
<tr>
<td>92/05</td>
<td>J.P.H.W.v.d.Eijnde</td>
<td>Conservative fixpoint functions on a graph</td>
</tr>
<tr>
<td>92/06</td>
<td>J.C.M. Baeten J.A. Bergstra</td>
<td>Discrete time process algebra</td>
</tr>
<tr>
<td>92/07</td>
<td>R.P. Nederpelt</td>
<td>The fine-structure of lambda calculus</td>
</tr>
<tr>
<td>92/08</td>
<td>R.P. Nederpelt F. Kamareddine</td>
<td>On stepwise explicit substitution</td>
</tr>
<tr>
<td>92/09</td>
<td>R.C. Backhouse</td>
<td>Calculating the Warshall/Floyd path algorithm</td>
</tr>
<tr>
<td>92/10</td>
<td>P.M.P. Rambags</td>
<td>Composition and decomposition in a CPN model</td>
</tr>
<tr>
<td>92/11</td>
<td>R.C. Backhouse J.S.C.P.v.d.Woude</td>
<td>Demonic operators and monotype factors</td>
</tr>
<tr>
<td>92/12</td>
<td>F. Kamareddine</td>
<td>Set theory and nominalisation, Part I</td>
</tr>
<tr>
<td>92/13</td>
<td>F. Kamareddine</td>
<td>Set theory and nominalisation, Part II</td>
</tr>
<tr>
<td>92/14</td>
<td>J.C.M. Baeten</td>
<td>The total order assumption</td>
</tr>
<tr>
<td>92/15</td>
<td>F. Kamareddine</td>
<td>A system at the cross-roads of functional and logic programming</td>
</tr>
<tr>
<td>92/16</td>
<td>R.R. Seljée</td>
<td>Integrity checking in deductive databases; an exposition</td>
</tr>
<tr>
<td>92/17</td>
<td>W.M.P. van der Aalst</td>
<td>Interval timed coloured Petri nets and their analysis</td>
</tr>
<tr>
<td>92/18</td>
<td>R.Nederpelt F. Kamareddine</td>
<td>A unified approach to Type Theory through a refined lambda-calculus</td>
</tr>
<tr>
<td>92/19</td>
<td>J.C.M Baeten J.A.Bergstra S.A.Smolka</td>
<td>Axiomatizing Probabilistic Processes: ACP with Generative Probabilities</td>
</tr>
<tr>
<td>92/20</td>
<td>F.Kamareddine</td>
<td>Are Types for Natural Language?</td>
</tr>
</tbody>
</table>
92/21 F.Kamareddine
Non well-foundedness and type freeness can unify the interpretation of functional application, p. 16.

92/22 R. Nederpclt
A useful lambda notation, p. 17.
F.Kamareddine

92/23 F.Kamareddine
Nominalization, Predication and Type Containment, p. 40.
E.Klein

92/24 M.Codish
Bottom-up Abstract Interpretation of Logic Programs, p. 33.
D.Dams
Eyal Yardeni

92/25 E.Poll
A Programming Logic for Fo, p. 15.

92/26 T.H.W.Beelen
A modelling method using MOVIE and SimCon/ExSpect, p. 15.
W.J.J.Sut
P.A.C.Verkoulen

92/27 B. Watson
A taxonomy of keyword pattern matching algorithms, p. 50.
G. Zwaan

93/01 R. van Geldrop
Deriving the Aho-Corasieck algorithms: a case study into the synergy of programming methods, p. 36.

93/02 T. Verhoeff
A continuous version of the Prisoner's Dilemma, p. 17

93/03 T. Verhoeff
Quicksort for linked lists, p. 8.

93/04 E.H.L. Aans
Deterministic and randomized local search, p. 78.
J.H.M. Korst
P.J. Zwietering

93/05 J.C.M. Baeten
A congruence theorem for structured operational semantics with predicates, p. 18.
C. Verhoef

93/06 J.P. Veltkamp
On the unavoidability of metastable behaviour, p. 29

93/07 P.D. Moerland
Exercises in Multiprogramming, p. 97

93/08 J. Verhoosel
A Formal Deterministic Scheduling Model for Hard Real-Time Executions in DEDOS, p. 32.

93/09 K.M. van Hee

93/10 K.M. van Hee
Systems Engineering: a Formal Approach Part II: Frameworks, p. 44.

93/11 K.M. van Hee

93/12 K.M. van Hee

93/13 K.M. van Hee
Systems Engineering: a Formal Approach
<table>
<thead>
<tr>
<th>Book ID</th>
<th>Authors</th>
<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>93/16</td>
<td>H. Schepers, J. Hooman</td>
<td>A Trace-Based Compositional Proof Theory for Fault Tolerant Distributed Systems, p. 27</td>
</tr>
<tr>
<td>93/17</td>
<td>D. Alstein, P. van der Stok</td>
<td>Hard Real-Time Reliable Multicast in the DEDOS system, p. 19.</td>
</tr>
<tr>
<td>93/18</td>
<td>C. Verhoef</td>
<td>A congruence theorem for structured operational semantics with predicates and negative premises, p. 22.</td>
</tr>
<tr>
<td>93/19</td>
<td>G.J. Houben</td>
<td>The Design of an Online Help Facility for ExSpect, p. 21.</td>
</tr>
<tr>
<td>93/22</td>
<td>E. Poll</td>
<td>A Typechecker for Bijective Pure Type Systems, p. 28.</td>
</tr>
<tr>
<td>93/23</td>
<td>E. de Kogel</td>
<td>Relational Algebra and Equational Proofs, p. 23.</td>
</tr>
<tr>
<td>93/24</td>
<td>E. Poll and Paula Severi</td>
<td>Pure Type Systems with Definitions, p. 38.</td>
</tr>
<tr>
<td>93/26</td>
<td>W.M.P. van der Aalst</td>
<td>Multi-dimensional Petri nets, p. 25.</td>
</tr>
<tr>
<td>93/27</td>
<td>T. Kloks and D. Kratsch</td>
<td>Finding all minimal separators of a graph, p. 11.</td>
</tr>
<tr>
<td>93/28</td>
<td>F. Kamareddine and R. Nederpelt</td>
<td>A Semantics for a fine λ-calculus with de Bruijn indices, p. 49.</td>
</tr>
<tr>
<td>93/29</td>
<td>R. Post and P. De Bra</td>
<td>GOLD, a Graph Oriented Language for Databases, p. 42.</td>
</tr>
<tr>
<td>93/30</td>
<td>J. Deogun, T. Kloks, D. Kratsch, H. Müller</td>
<td>On Vertex Ranking for Permutation and Other Graphs, p. 11.</td>
</tr>
<tr>
<td>93/31</td>
<td>W. Körver</td>
<td>Derivation of delay insensitive and speed independent CMOS circuits, using directed commands and production rule sets, p. 40.</td>
</tr>
</tbody>
</table>
93/33 L. Loyens and J. Moonen  
ILIAS, a sequential language for parallel matrix computations, p. 20.

93/34 J.C.M. Baeten and J.A. Bergstra  
Real Time Process Algebra with Infinitesimals, p. 39.

93/35 W. Ferrer and P. Severi  
Abstract Reduction and Topology, p. 28.

93/36 J.C.M. Baeten and J.A. Bergstra  
Non Interleaving Process Algebra, p. 17.

93/37 J. Brunekreef J-P. Katoen R. Koymans S. Mauw  
Design and Analysis of Dynamic Leader Election Protocols in Broadcast Networks, p. 73.

93/38 C. Verhoef  
A general conservative extension theorem in process algebra, p. 17.

93/39 W.P.M. Nuijten E.H.L. Aarts D.A.A. van Erp Taalman Kip K.M. van Hee  
Job Shop Scheduling by Constraint Satisfaction, p. 22.

93/40 P.D.V. van der Stok M.M.M.P.J. Claessen D. Alstein  

93/41 A. Bijlsma  
Temporal operators viewed as predicate transformers, p. 11.

93/42 P.M.P. Rambags  
Automatic Verification of Regular Protocols in P/T Nets, p. 23.

93/43 B.W. Watson  
A taxonomy of finite automata construction algorithms, p. 87.

93/44 B.W. Watson  
A taxonomy of finite automata minimization algorithms, p. 23.

93/45 E.J. Luit J.M.M. Martin  
A precise clock synchronization protocol, p.

93/46 T. Kloks D. Kratsch J. Spinrad  

93/47 W. v.d. Aalst P. De Bra G.J. Houben Y. Komatzky  

93/48 R. Gerth  
Verifying Sequentially Consistently Memory using Interface Refinement, p. 20.
<table>
<thead>
<tr>
<th>Paper Code</th>
<th>Authors</th>
<th>Title</th>
</tr>
</thead>
<tbody>
<tr>
<td>94/01</td>
<td>P. America, M. van der Kammen, R.P. Nederpelt, O.S. van Roosmalen, H.C.M. de Swart</td>
<td>The object-oriented paradigm, p. 28.</td>
</tr>
<tr>
<td>94/02</td>
<td>F. Kamareddinc, R.P. Nederpelt</td>
<td>Canonical typing and Π-conversion, p. 51.</td>
</tr>
<tr>
<td>94/04</td>
<td>J.C.M. Baeten, J.A. Bergstra</td>
<td>Graph isomorphism Models for Non Interleaving Process Algebra, p. 18.</td>
</tr>
<tr>
<td>94/06</td>
<td>T. Basten, T. Kunz, J. Black, M. Coffin, D. Taylor</td>
<td>Time and the Order of Abstract Events in Distributed Computations, p. 29.</td>
</tr>
<tr>
<td>94/08</td>
<td>O.S. van Roosmalen</td>
<td>A Hierarchical Diagrammatic Representation of Class Structure, p. 22.</td>
</tr>
<tr>
<td>94/09</td>
<td>J.C.M. Baeten, J.A. Bergstra</td>
<td>Process Algebra with Partial Choice, p. 16.</td>
</tr>
<tr>
<td>94/10</td>
<td>T. verhoeff</td>
<td>The testing Paradigm Applied to Network Structure, p. 31.</td>
</tr>
<tr>
<td>94/13</td>
<td>R. Seljée</td>
<td>A New Method for Integrity Constraint checking in Deductive Databases, p. 34.</td>
</tr>
<tr>
<td>94/14</td>
<td>W. Peremans</td>
<td>Ups and Downs of Type Theory, p. 9.</td>
</tr>
<tr>
<td>94/16</td>
<td>R.C. Backhouse, H. Doombos</td>
<td>Mathematical Induction Made Calculational, p. 36.</td>
</tr>
</tbody>
</table>
94/18  F. Kamareddine
         R. Nederpelt  Refining Reduction in the Lambda Calculus, p. 15.

94/19  B.W. Watson  The performance of single-keyword and multiple-keyword pattern matching algorithms, p. 46.

94/20  R. Bloo
         F. Kamareddine
         R. Nederpelt  Beyond β-Reduction in Church’s λ→, p. 22.