Rotated and scaled alamouti coding

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Rotated and Scaled Alamouti Coding

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Abstract—Repetition-based retransmission is used in Alamouti-modulation [1998] for $2 \times 2$ MIMO systems. We propose to use instead of ordinary repetition so-called "scaled repetition" together with rotation. It is shown that the rotated and scaled Alamouti code has a hard-decision performance which is only slightly worse than that of the Golden code [2005], the best known $2 \times 2$ space-time code. Decoding the Golden code requires an exhaustive search over all the codewords (or sphere decoding for higher spectral efficiencies), while our rotated and scaled Alamouti code can be decoded with an acceptable complexity.

I. SCALED-REPLICATION FOR THE SISO CHANNEL

First we consider transmission over a single-input single-output (SISO) additive white Gaussian noise (AWGN) channel (see Fig. 1), and introduce scaled-repetition retransmission. It turns out that scaled-repetition improves upon ordinary-repetition retransmission.

A. Some information theory

![Diagram of AWGN channel](image)

Fig. 1. The AWGN channel.

The real-valued output $y_k$ for transmission $k = 1, 2, \cdots, K$, see Fig. 1, satisfies

$$y_k = x_k + n_k, \quad (1)$$

where $x_k$ is the real-valued channel input for transmission $k$ and $n_k$ is a real-valued Gaussian noise sample with mean $E[N_k] = 0$, variance $E[N_k^2] = \sigma^2$, which is uncorrelated with all other noise samples. The transmitter power is limited, i.e. we require that $E[X_k^2] \leq P$. It is well-known that an $X$ which is Gaussian with mean 0 and variance $P$ achieves capacity. This basic capacity (in bit/transition) equals

$$C = \frac{1}{2} \log_2 (1 + \frac{P}{\sigma^2}). \quad (2)$$

When we retransmit (repeat) codewords, each symbol $x_k$ from such a codeword $(x_1, x_2, \cdots, x_K)$ is actually transmitted and received twice, i.e. $x_{k1} = x_{k2} = x_k$, and

$$y_{k1} = x_k + n_{k1}, \text{ and } y_{k2} = x_k + n_{k2}. \quad (3)$$

An optimal receiver can form $z_k = \frac{y_{k1} + y_{k2}}{2} = x_k + \frac{n_{k1} + n_{k2}}{2}$. Now the variance of the noise sample $(N_{k1} + N_{k2})/2$ is $\sigma^2/2$. Therefore the repetition capacity for a single repetition in bit/transition is

$$C_r = \frac{1}{4} \log_2 (1 + \frac{2P}{\sigma^2}). \quad (4)$$

We can use Benelli’s [3] method to improve upon ordinary-repetition retransmission, i.e. by modulating the retransmitted symbol differently. We could e.g. take

$$x_k = x_k, \text{ and } x_{k2} = M_2(x_k) \text{ for } x_k \in A_{4\text{-PAM}}, \quad (6)$$

where $M_2(\alpha) = 2\alpha - 5$ if $\alpha > 0$ and $M_2(\alpha) = 2\alpha + 5$ for $\alpha < 0$. We call this method scaled repetition since we scale a symbol by a factor (2 here) and then compensate (+5 or -5) in order to obtain a symbol from $A_{4\text{-PAM}}$. This results in the signal points $(x, M_2(x))$ for $x \in A_{4\text{-PAM}}$, see Fig. 2, right part. Also for the scaled-repetition case the maximum transmission rate $I(X; Y_1, Y_2)$ is shown in Figure 3, with red asterisks. Note that this maximum transmission rate is only slightly smaller than the basic capacity $C$. Ordinary repetition is however definitively inferior to the basic transmission if the SNR is not very small.

![Diagram of signal points](image)
C. Demodulation complexity

Scaled repetition outperforms ordinary repetition, but also
has a disadvantage. In an ordinary-repetition system the output
\( y_k = (y_{k1} + y_{k2})/2 \) is simply sliced. In a system that uses
scaled repetition we can only slice after having distinguished
between two cases. More precisely note that \( x_{k2} = M_2(x_k) =
2x_k - D_2(x_k) \), where \( D_2(\alpha) = 5 \) if \( \alpha > 0 \) and \( D_2(\alpha) = -5 \)
if \( \alpha < 0 \). Now we can use a slicer for \( y_{k1} + 2y_{k2} = x_k +
n_{k1} + 2(2x_k - D_2(x_k) + n_{k2}) = 5x_k - 2D_2(x_k) + n_{k1} + 2n_{k2} \).
Assuming that \( x_k \in \{-3, -1\} \) we get that \( D_2(x_k) = -5 \) and
this implies that we should put a threshold at 0 to distinguish
between -3 and -1. Similarly assuming that \( x_k \in \{+1, +3\} \)
we get \( D_2(x_k) = 5 \) and we must slice \( y_{k1} + 2y_{k2} \) again with
a threshold at 0. Then the best overall candidate \( x^*_k \) is found
by minimizing \( (y_{k1} - x^*_k)^2 + (y_{k2} - M_2(x^*_k))^2 \) over the two
candidates.

II. FUNDAMENTAL PROPERTIES FOR THE 2 × 2 MIMO
CHANNEL

A. Model description

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Fig. 4. Model of a 2 × 2 MIMO channel.
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Next consider a 2 × 2 MIMO channel (see Fig. 4). Both
the transmitter and the receiver use two antennas. The output
vector \( (y_{k1}, y_{k2}) \) at transmission \( k \) relates to the corresponding
input vector \( (x_{k1}, x_{k2}) \) as given by

\[
\begin{pmatrix}
    y_{k1} \\
    y_{k2}
\end{pmatrix}
= \begin{pmatrix}
    h_{11} & h_{12} \\
    h_{21} & h_{22}
\end{pmatrix}
\begin{pmatrix}
    x_{k1} \\
    x_{k2}
\end{pmatrix}
+ \begin{pmatrix}
    n_{k1} \\
    n_{k2}
\end{pmatrix}
\tag{7}
\]

where \((N_{1k}, N_{2k})\) is a pair of independent zero-mean cir-
cularly symmetric complex Gaussians, both having variance
\( \sigma^2 \) (per two dimensions). Noise variable pairs in different transmissions are independent.

We assume that the four channel coefficients \( H_{11}, H_{12}, H_{21}, \) and \( H_{22} \) are independent zero-mean circularly symmetric com-
plex Gaussians, each having variance 1 (per two dimensions). The channel coefficients are chosen prior to a block of \( K \)
transmissions and remain constant over that block.

The complex transmitted symbols \((X_{k1}, X_{k2})\) must satisfy a
power constraint, i.e.

\[
E[X_{k1}X_{k1}^* + X_{k2}X_{k2}^*] \leq P.
\tag{8}
\]

B. Telatar capacity

If the channel input variables are independent zero-mean circularly symmetric complex Gaussians both having variance \( P/2 \), then the resulting mutual information (called Telatar capacity here, see [6]) is

\[
C_{\text{Telatar}}(H) = \log_2 \det(I_2 + \frac{P}{\sigma^2} HH^\dagger),
\tag{9}
\]

where \( H = \begin{pmatrix} h_{11} & h_{12} \\ h_{21} & h_{22} \end{pmatrix} \), i.e. the actual channel-coefficient
matrix and \( I_2 \) the 2 × 2 identity matrix. Also in the 2 × 2
MIMO case we define the signal-to-noise ratio as

\[
\text{SNR} \triangleq \frac{P}{\sigma^2}.
\tag{10}
\]

It can be shown (see e.g. Yao ([7], p. 36) that for fixed \( R \)
and SNR large enough \( \Pr\{C_{\text{Telatar}}(H) < R\} \approx \gamma \cdot \text{SNR}^{-4} \), for some constant \( \gamma \).

C. Worst-case error-probabilities

Consider \( M \) (one for each message) \( K \times 2 \) code-matrices
\( \Xi_1, \Xi_2, \ldots, \Xi_M \) resulting in a unit average energy code. Then
Tarokh, Seshadri and Calderbank [5] showed that for large
SNR

\[
\Pr\{\xi \to \xi'\} \approx \gamma' (\det((\xi' - \xi)(\xi' - \xi)^*)^4)^{-2} \text{SNR}^{-4}.
\tag{11}
\]

for some \( \gamma' \) if the rank of the difference matrices \( \xi - \xi' \) is 2, and we transmit \( \xi = \sqrt{P} \xi \). This holds for all difference matrices
we say that the diversity order is 4. Therefore it makes sense
to maximize the minimum modulus of the determinant over all
code-matrix differences.

III. ALAMOUTI: ORDINARY REPEITION

Alamouti [1] proposed a modulation scheme (space-time
code) for the 2 × 2 MIMO channel which allows for a very
simple detector. Two complex symbols \( s_1 \) and \( s_2 \) are trans-
mitted in the first transmission (an odd transmission) and in
the second transmission (the next even transmission) these
symbols are more or less repeated. More precisely

\[
\begin{pmatrix}
    x_{11} \\
    x_{12} \\
    x_{21} \\
    x_{22}
\end{pmatrix}
= \begin{pmatrix}
    s_1 \\
    s_2 \\
    -s_1^* \\
    s_1^*
\end{pmatrix}.
\tag{12}
\]

The received signal is now

\[
\begin{pmatrix}
    y_{11} \\
    y_{12} \\
    y_{21} \\
    y_{22}
\end{pmatrix}
= \begin{pmatrix}
    h_{11} & h_{12} & h_{11}^* & h_{12}^* \\
    h_{21} & h_{22} & h_{21}^* & h_{22}^*
\end{pmatrix}
\begin{pmatrix}
    s_1 \\
    s_2 \\
    -s_1^* \\
    s_1^*
\end{pmatrix}
+ \begin{pmatrix}
    n_{11} \\
    n_{21} \\
    n_{12} \\
    n_{22}
\end{pmatrix}.
\tag{13}
\]
or more compactly
\[ y = s_1\hat{a} + s_2\hat{b} + \hat{n}, \]
with
\[ \hat{y} = (\hat{y}_{11}, \hat{y}_{21}, \hat{y}_{12}, \hat{y}_{22})^T, \]
\[ \hat{a} = (h_{11}, h_{21}, h_{12}^*, h_{22}^*)^T, \]
\[ \hat{b} = (h_{12}, h_{22}, -h_{11}, -h_{21})^T, \]
and
\[ \hat{n} = (n_{11}, n_{21}, n_{12}^*, n_{22}^*)^T. \] (14)

Since \( \hat{a} \) and \( \hat{b} \) are orthogonal the symbol estimates \( \hat{s}_1 \) and \( \hat{s}_2 \) can be determined by simply slicing \((\hat{a}^\dagger \hat{y})/(\hat{a}^\dagger \hat{a})\) and \((\hat{b}^\dagger \hat{y})/(\hat{b}^\dagger \hat{b})\) respectively.

Another advantage of the Alamouti method is that the densities of \( \hat{a}^\dagger \hat{a} \) and \( \hat{b}^\dagger \hat{b} \) are (identical and) chi-square with 8 degrees of freedom. This results in a diversity order 4, i.e.
\[ \Pr\{ (\hat{S}_1, \hat{S}_2) \neq (S_1, S_2) \} \approx \gamma'' \cdot \text{SNR}^{-4}, \] (15)
for fixed rate and large enough SNR.

A disadvantage of the Alamouti method is that only two complex sign of \( X \) are transmitted every two transmissions, but more importantly that the symbols transmitted in the second transmission are more or less repetitions of the symbols in the first transmission. Section I however suggests that we can improve upon ordinary repetition.

A. Method description

Having seen in section I that scaled-repetition improves upon ordinary repetition in the SISO case, we use this concept to improve upon the standard Alamouti scheme for MIMO transmission. Instead of just repeating the symbols in the second transmission we scale them. More precisely, when \( s_1 \) and \( s_2 \) are elements of \( A_{16-QAM} \), we could transmit for some value of \( \theta \) the signals
\[ \begin{pmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{pmatrix} = \begin{pmatrix} s_1 \cdot \exp(j\theta) & -s_2^* \\ s_2 \cdot \exp(j\theta) & s_1^* \end{pmatrix} \begin{pmatrix} M_{2}(s_2) & M_{2}(s_1^*) \\ 2s_2 & 2s_1^* \end{pmatrix} = \begin{pmatrix} 0 & 0 \\ D_2(s_2) & D_2(s_1^*) \end{pmatrix}, \] (16)
where \( M_2(\alpha) = 2\alpha - D_2(\alpha) \) with \( D_2(\alpha) = 5\beta \) when \( \beta \) is the complex sign of \( \alpha \).

A first question is to determine a good value for \( \theta \). Therefore we determine for \( 0 \leq \theta \leq \pi/2 \) the minimum modulus of the determinant \( \min_{(s_1, s_2), (s'_1, s'_2)} |\text{det}(X(s_1, s_2, \theta) - X(s'_1, s'_2, \theta))| \),
\[ X = \begin{pmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{pmatrix} \] (17)
where \( X \) is the code matrix. The minimum modulus of the determinant as a function of \( \theta \) can be found in Fig. 5. The maximum value of the determinant minimum determinant (i.e. 7,613) occurs for
\[ \theta_{opt} = 1.028. \] (18)

We will use this value for \( \theta \) in what follows.

B. Hard-decision Performance

We have compared the message-error-rate for several \( R = 4 \) space-time codes in Fig. 6. By message-error-rate we mean the probability \( \Pr\{ \hat{X} \neq X \} \). Note that for each “test” we generate a new message (8-bit) and a new channel matrix. The decoder is optimal for all codes, it performs ML-decoding (exhaustive search). The methods that we have considered are:

1) Uncoded, in green. We transmit
\[ X = \begin{pmatrix} x_{11} & x_{12} \\ x_{21} & x_{22} \end{pmatrix}, \] (19)
where \( x_{11}, x_{12}, x_{21}, \) and \( x_{22} \) are symbols from \( A_{4-QAM} \).

2) Alamouti, in blue, see (12), where \( s_1 \) and \( s_2 \) are symbols from \( A_{16-QAM} \).

3) Tilted QAM, in cyan. Proposed by Yao and Wornell [8]. Let \( s_a, s_b, s_c, \) and \( s_d \) symbols from \( A_{4-QAM} \). Then we transmit
\[ \begin{pmatrix} x_{11} \\ x_{22} \end{pmatrix} = \begin{pmatrix} \cos(\theta_1) & -\sin(\theta_1) \\ \sin(\theta_1) & \cos(\theta_1) \end{pmatrix} \begin{pmatrix} s_a \\ s_b \end{pmatrix}, \]
\[ \begin{pmatrix} x_{21} \\ x_{12} \end{pmatrix} = \begin{pmatrix} \cos(\theta_2) & -\sin(\theta_2) \\ \sin(\theta_2) & \cos(\theta_2) \end{pmatrix} \begin{pmatrix} s_c \\ s_d \end{pmatrix} \] (20)
for \( \theta_1 = \frac{1}{2} \arctan(\frac{1}{2}) \) and \( \theta_2 = \frac{1}{2} \arctan(2) \).

4) Rotated and scaled Alamouti, in red, see (16) for \( \theta = 1.028 \), and with \( s_1 \) and \( s_2 \) from \( A_{16-QAM} \).

5) Golden code, in magenta. Proposed by Belfiore et al. [2]. Now
\[ X = \frac{1}{\sqrt{\overline{2}}} \begin{pmatrix} \alpha(z_1 + z_2 \overline{\alpha}) \\ j \cdot \frac{\alpha(z_3 + z_4 \overline{\alpha})}{\overline{\alpha}(z_1 + z_2 \overline{\alpha})} \end{pmatrix}, \] (21)
with \( \overline{\alpha} = \frac{1+j}{\sqrt{2}}, \overline{\theta} = \frac{1-j}{\sqrt{2}}, \alpha = 1 + j - j\theta, \) and \( \overline{\alpha} = 1 + j - j\overline{\theta} \) and where \( z_1, z_2, z_3, \) and \( z_4 \) are \( A_{4-QAM} \) symbols.

6) Telatar, in black. This is the probability that the Telatar capacity of the channel is smaller than 4.

Clearly it follows from Fig. 6 that the winner is the Golden code. However rotated and scaled Alamouti is only slightly worse, roughly 0.2 dB. Important is that Alamouti coding is roughly 2 dB worse than the Golden code.
we can write
\[
\begin{pmatrix}
x_{11} \\
x_{21} \\
x_{12} \\
x_{22}
\end{pmatrix}
= \begin{pmatrix} -M_2(t_1) & M_2(t_2) \\ -2t_1 & 2t_2 \\ 0 & 0 \end{pmatrix}
\begin{pmatrix} t_1 \\ t_2 \\ t_1^* \\ t_2^* \end{pmatrix},
\]
(24)
since \( t = M_2(s) \) implies that \( s = -M_2(t) \). Now
\[
\begin{pmatrix}
y_{11} \\
y_{21} \\
y_{12} \\
y_{22}
\end{pmatrix}
= \begin{pmatrix} -2h_{11} & h_{12} \\ -2h_{21} & h_{22} \\ h_{12}^* & 2h_{11}^* \\ h_{22}^* & 2h_{21}^* \end{pmatrix}
\begin{pmatrix} t_1 \\ t_2 \end{pmatrix}
- \begin{pmatrix} -h_{11} & 0 \\ 0 & h_{11}^* \end{pmatrix}
\begin{pmatrix} D_2(t_1) \\ D_2(t_2) \end{pmatrix}
+ \begin{pmatrix} n_{11} \\ n_{21} \\ n_{12} \\ n_{22} \end{pmatrix},
\]
(25)
We can write this as
\[
y = t_1 g + t_2 b - D_2(t_1) c - D_2(t_2) d + n,
\]
\[
g' = (2h_{11}, -2h_{21}, h_{12}^*, h_{22}^*)^T,
\]
\[
b' = (h_{12}, h_{22}, 2h_{11}^*, 2h_{21}^*)^T,
\]
\[
c' = (2h_{11}^*, -2h_{21}^*, 0, 0)^T,
\]
\[
d' = (0, 0, h_{11}^*, h_{22}^*)^T,
\]
and for the "\(\cos(\phi')\)" of the angle between \( g' \) and \( b' \) we can write
\[
\cos(\phi') = \frac{2(\Theta - 1)(h_{11}h_{12}^* + h_{21}h_{22}^*)}{|2(\Theta - 1)h_{11}|^2 + |2h_{12}|^2 + |h_{21}|^2 + |h_{22}|^2}.
\]
(26)
C. It now follows from the inequality \( 2r_1r_2 \leq r_1^2 + r_2^2 \) (where \( r_1 \) and \( r_2 \) are reals), that
\[
\cos(\phi) \leq |\Theta - 1| \cdot \frac{|h_{11}|^2 + |h_{12}|^2 + |h_{21}|^2 + |h_{22}|^2}{|2|}.
\]
\[
\cos(\phi') \leq |\Theta - 1| \cdot \frac{|h_{11}|^2 + |h_{12}|^2 + |h_{21}|^2 + |h_{22}|^2}{4|}.
\]
(27)
If
\[
|h_{12}|^2 + |h_{22}|^2 \geq |h_{11}|^2 + |h_{21}|^2,
\]
then \( \cos(\phi) \leq \frac{2|\Theta - 1|}{3} = 0.393 \), else \( \cos(\phi') \leq \frac{2|\Theta - 1|}{3} = 0.393 \). Therefore it makes sense to decode \((s_1, s_2)\) when (28) holds and \((t_1, t_2)\) when (28) does not hold. Using zero-forcing to decode, the noise enhancement is then at most \(1/1 - 0.393^2 = 1.183\) which is 0.729 dB. We shall see later that noise enhancement is unnoticeable in practise.

D. The decoding procedure is straightforward. Focus on the case where we decode \((s_1, s_2)\) for a moment. For all 16 alternatives of \((D_2(s_1), D_2(s_2))\) the vector
\[
\tilde{z} = y + D_2(s_1)c + D_2(s_2)d = s_1a + s_2b + n
\]
(29)
and is determined. Then compute the sufficient statistic
\[
\left( \begin{array}{c}
\tilde{a} \\
\tilde{b}
\end{array} \right) = \left( \begin{array}{cc}
a^T & b^T \\
a^T & b^T
\end{array} \right) \left( \begin{array}{c}
s_1 \\
s_2
\end{array} \right) + \left( \begin{array}{cc}
a^T & b^T \\
a^T & b^T
\end{array} \right) \left( \begin{array}{c}
n_1 \\
n_2
\end{array} \right).
\]
(30)
Use inverted matrix $M = \left( \begin{array}{ccc} b_1^r b_1 & -a_1^r b_1 \\ b_1^r a_1 & a_1^r a_1 \end{array} \right) / D$ where $D = (a_1^r a_1)(b_1^r b_1) - (b_1^r a_1)(a_1^r b_1)$ to obtain $\left( \begin{array}{c} s_1 \\ s_2 \end{array} \right) = M \left( \begin{array}{c} a^r z \\ b^r z \end{array} \right)$. Next both $s_1$ and $s_2$ are sliced under the restriction that only alternatives that match the assumed values $D_2(s_1)$ and $D_2(s_2)$ are possible outcomes. This is done for all 16 alternatives $(D_2(s_1), D_2(s_2))$. The best result is in terms of Euclidean distance is now chosen.

In considering all alternatives $(D_2(s_1), D_2(s_2))$ we only need to slice when the length of $z - s_1 a - s_2 b$ is smaller than the closest distance we have observed so far. This reduces the number of slicing steps. We call this approach METHOD 1.

E. The number of slicing steps can even be further decreased if we start slicing with the most promising alternative $(D_2(s_1), D_2(s_2))$. This approach is called METHOD 2. Therefore we note that the "direct" $s_1$-signal-component in $X$ is $\left( \begin{array}{cc} s_1 \Theta & 0 \\ 0 & -s_1^2 / 2 \end{array} \right)$. Therefore we can slice $(c_1^r y) / (c_1^r c_1)$ in order to find a good guess for $D_2(s_1)$. Similarly we slice $(c_2^r y) / (c_2^r c_2)$ to find a good first guess for $D_2(s_2)$. Here

$$\begin{align*}
\xi_1 &= \left( h_{11} \Theta, h_{21} \Theta, -h_{12}^r / 2, -h_{22}^r / 2 \right)^T, \\
\xi_2 &= \left( -h_{12}^r / 2, -h_{22}^r / 2, -h_{11}^r, -h_{21}^r \right)^T.
\end{align*}$$

(31)

Then we consider the other 15 alternatives and only slice if necessary. Similar methods apply if we want to decode $(t_1, t_2)$.

F. We have carried out simulations, first to find out what the degradation of the suboptimal decoders according to method 1 and method 2 is relative to ML-decoding. The result is shown in Fig. 7. Conclusion is that the suboptimal decoders do not demonstrate a performance degradation. We have also considered the number of slicings for both method 1 and method 2. This is shown in Fig. 8. It can be observed that method 1 leads to roughly 7 slicings (as opposed to 16). Method 1 further decreases the number of slicing to roughly 3.5.

VI. CONCLUSION

Rotated and scaled Alamouti has a hard-decision performance which is only slightly worse than that of the Golden code, but can be decoded with an acceptable complexity. We have obtained similar results for codes based on mapping $M_3(\cdot)$ for 9-PAM, see Fig. 9. Recently also Sezginer and Sari [4] investigated complexity reducing methods for alternatives to the Golden code.

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